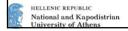
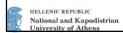
Decision-Making for Uncoordinated User Access to Limited (Distributed) Resources
Or
Resource Competition
in a Highly Networked World of
Humans and Things

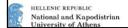


Decision-Making for <u>Uncoordinated</u> User Access to <u>Limited</u> (Distributed) Resources



Decision-Making for Uncoordinated User Access to (Distributed) Limited Resources

*** A classical and very old problem! ***



Accessing a Single (common) Acoustical Channel by uncoordinated Users

(... an old problem for humans, e.g., at a cocktail party ...)

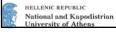


(shared air, acoustical)

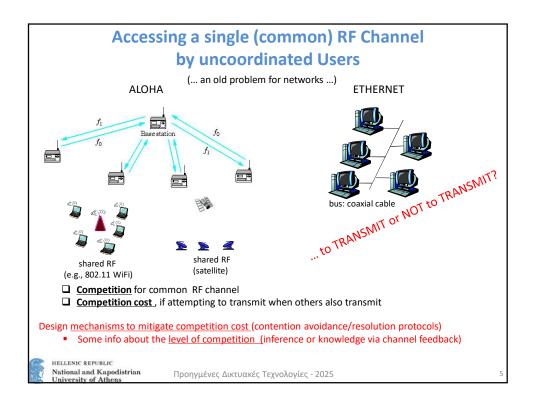
- ☐ <u>Competition</u> for common acoustical channel
- ☐ Competition cost, if attempting to talk when others also talk

Humans adopt behaviors to mitigate competition costs (conflict avoidance / resolution)

Some info about the <u>level of competition</u> (sense, inference or knowledge)



Προηγμένες Δικτυακές Τεχνολογίες - 2025

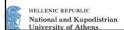


Uncoordinated access of a common resource by distributed users can bring benefits

(simplicity, efficiency, feasibility, privacy,...)

.... but it costs

(collisions / their resolution and avoidance)



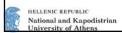
Προηγμένες Δικτυακές Τεχνολογίες - 2025

"price of Anarchy" (PoA) * E. Koutsoupias, C. H. Papadimitriou, "Worst-case equilibria", Computer Science Review, 2009. PoA= { Cost of the uncoordinated approach } / { Cost under optimal (coordinated) approach } = {max benefit (coordinated)}/{ max benefit of the (uncoordinated) approach} max throughput of ALOHA \neq 0.36) (of 0.18) without even time coordination) max throughput (coordinated) € 1.00 PoA = $\{1\}/\{0.36\}$ (2.77) or $\{1\}/\{0.16\}$ (5.55) without even time coordination

Accessing a Common/Limited Resource by uncoordinated Users

Προηγμένες Δικτυακές Τεχνολογίες - 2025

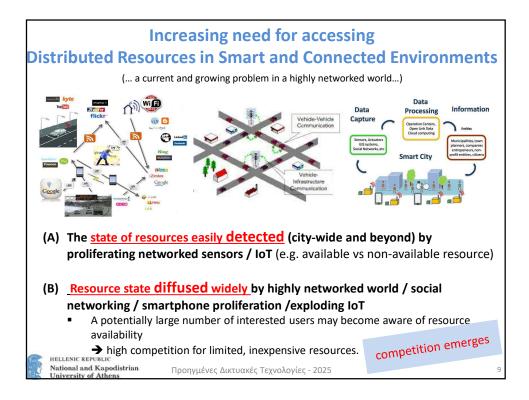
*An increasingly relevant problem appearing in today's smart and connected world! *



For ALOHA

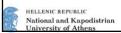
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Προηγμένες Δικτυακές Τεχνολογίες - 2025



Is wide resource state diffusion always good?

(more) information → easier/better decisions



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Some earlier work on benefits of wide information diffusion

Enhance <u>NAPS</u> (<u>Non-Assisted Parking Search</u>) exploiting ICT-enabled information diffusion

OAPS: opportunistically-assisted parking search

Parking spots equipped with sensors

Vehicles with wireless interfaces / storage / processing capability

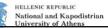
Inter-vehicle & vehicleparking spot communication







E. Kokolaki, M. Karaliopoulos, I. Stavrakakis, "Opportunistically-assisted parking service discovery: now it helps, now it does not", Pervasive and Mobile Computing (PMC), Elsevier, 2012.



Προηγμένες Δικτυακές Τεχνολογίες - 2025

11

Observation

under concentrated destinations within a particular (hotspot) road:

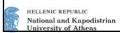
- ➤ information dissemination → synchronizes movement patterns
 - → intensifies competition / increases # of competitors

> OAPS worse than NAPS

(More) Information may yield worse performance!

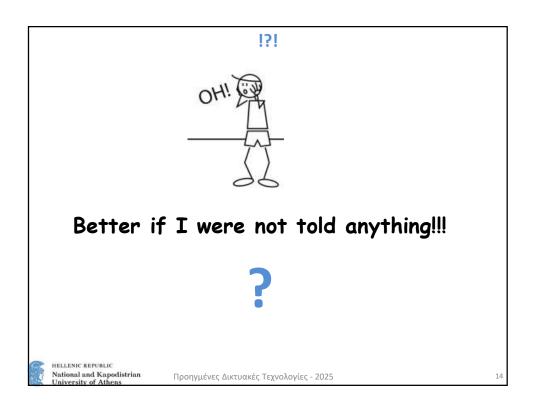
(" more is less")

Increased Competition is to blame!



Προηγμένες Δικτυακές Τεχνολογίες - 2025







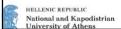
The broad environment and problem formulation

Some info about the state (availability) of a desirable low-cost resource provided to potential users.

- ☐ If a user attempts to use such a resource (COMPETE):
 - > if it is successful (i.e. not taken by others), it will incur a low cost.
 - ➤ If it fails, it will incur a failure cost. Then, the user will have to resort to an always available (unlimited) alternative resource and pay also the high cost of that resource.
- ☐ If a user ignores the low-cost resource availability info (NOT COMPETE) and goes for the high-cost (unlimited) resource, it will pay the high cost.

The challenge:

Users need to decide how to access the limited resources effectively – Humans are frequently driving such decisions



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Major Question posed

TO COMPETE OR NOT TO COMPETE?

that is the question



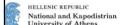
to compete for a limited and inexpensive resource or not and go for the unlimited, expensive alternative?

(cost of failing in the competition: use of the expensive resource PLUS pay a **failing penalty**)

Uncoordinated Access to Limited Resources

Competing for congestible goods

Resource Selection problem



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17

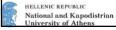
Example: Access point association

Cheap but best-effort (public) WLAN

vs. More reliable, high-speed fee-based wireless access point



Competition reduces quality of service

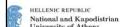


Προηγμένες Δικτυακές Τεχνολογίες - 2025





Competition reduces quality of service



Vs Faster toll road

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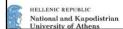
Example: Parking spot selection

Cheap but scarce (public) on-street parking

vs. Expensive but abundant parking lots



Competition reduces quality of service



Προηγμένες Δικτυακές Τεχνολογίες - 2025

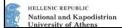
Problem Formulation

Uncoordinated Access to Limited Resources -Competing for congestible goods - Resource Selection problem

(parking spot resource example)

- o K distributed non-communicating users (competitors)
 - o aware of the availability of R resources
- o **Ternary** cost model for the users
 - Low Cost of Competing with success: Cosp.s
 - \circ <u>Medium</u> Cost for not Competing: $c_{pl} = β c_{osp,s}$ (β>1)
 - \circ <u>High</u> Cost of Competing and failing: $c_{osp,f} = \gamma c_{osp,s}$ (γ>β)

 $\delta = \gamma - \beta > 0$ (penalty factor of anarchy)



Προηγμένες Δικτυακές Τεχνολογίες - 2025

2

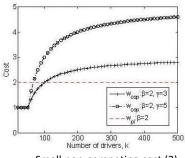
Congestion curve/dynamics: cost of a user when K compete

- \square Cost of a non-competing (pl) user is fixed (independent of k): $w_{pl}(k) = c_{pl}$
- ☐ Cost of a competing (osp) user when *K>R*:

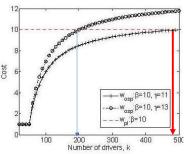
$$w_{osp}(k) = \min(1, R/k)c_{osp,s} + (1 - \min(1, R/k))\gamma c_{osp,s}$$

non-decreasing function of demand K

Up to 480 (δ =1) or 200 (δ =3) users is beneficiary to compete for 50 resources

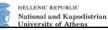


Small non-competing cost (2)



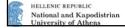
Large non-competing cost (10)

R=50 resources – competing with success cost, $c_{osp,s} = 1$



Προηγμένες Δικτυακές Τεχνολογίες - 2025

How are decisions to compete or not to compete taken?



Προηγμένες Δικτυακές Τεχνολογίες - 2025

2

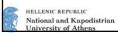
The full rationality vs. bounded rationality assumptions

Assumptions affecting how selections / decisions are made:

- Perfect vs. imperfect information/knowledge on number of competitors
- Full or limited computational capacity of decision-maker to assess the impact of choices
- Cognitive biases of decision-maker

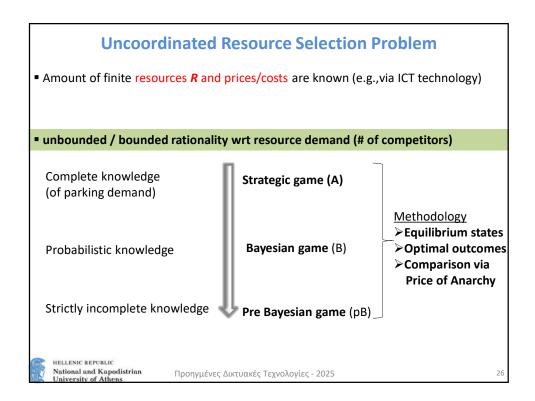
Uncoordinated Resource Selection problem formulation:

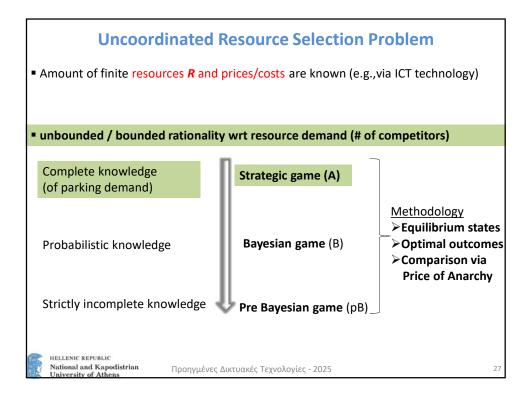
- ☐ Full Rationality: perfect knowledge, unlimited computational capacity
- ☐ Bounded Rationality: Limited / imperfect information / computational limits
- Bounded Rationality: (human-related) computational limits cognitive biases
 - inputs from behavioral economics, cognitive psychology, etc
 - models for (human-driven) bounded rationality



Προηγμένες Δικτυακές Τεχνολογίες - 2025

CASE A | Full Rationality: perfect knowledge, unlimited computational capacity | Bounded Rationality: Limited / imperfect information / computational limits | Study based on (classical) Expected Utility Maximization Framework | E. Kokolaki, M. Karaliopoulos, I. Stavrakakis, "Leveraging information in parking assistance systems", IEEE Transactions on Vehicular Technology, 2013. | | HELLENG REPUBLIC | RATIONAL AND KARPORTISTIAN | Προηγμένες Δικτυακές Τεχνολογίες - 2025 | 25





The (strategic) parking spot selection game $\Gamma(N)$

- N drivers / players
- R on-street-parking osp (public) + Infinite parking lot pl (private) spots
- Action set: {osp (public), pl (private) }
- Action of player i : α_i
- Actions of all players except player i: α_{-i}
- Action profile $\alpha = (\alpha_i, \alpha_{-i}) \rightarrow 2^{**}N$ of them
- Action meta-profile α(m): any profile with m players competing
 N+1 meta-profiles



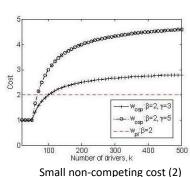
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RECALL: Cost for choosing action PL or OSP

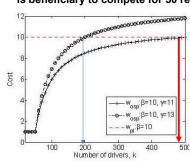
- \square Cost of a non-competing (pl) user is fixed (independent of k): $w_{nl}(k) = c_{nl}$
- \square Cost of a competing (osp) user when k>R:

$$w_{osp}(k) = \min(1, R/k)c_{osp,s} + (1 - \min(1, R/k))\gamma c_{osp,s}$$

non-decreasing function of demand k



Up to 480 (δ =1) or 200 (δ =3) users is beneficiary to compete for 50 resources



II non-competing cost (2) Large non-competing cost (10) R=50 resources – competing with success cost, $c_{osp,s} = 1$

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Προηγμένες Δικτυακές Τεχνολογίες - 2025

The (strategic) parking spot selection game

Expected **cost for player** *i* under action profile $\alpha = (\alpha_i, \alpha_{-i})$

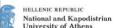
$$c_i^N(a_i,a_{-i}) = \left\{ \begin{array}{ll} w_{osp}(N_{osp}(a)), & \text{for } a_i = osp \\ \\ w_{pl}(N-N_{osp}(a)), & \text{for } a_i = pl \end{array} \right. \quad \text{($=c_{pl}$, fixed)}$$

■ Nash Equilibrium state:

a situation in which no player can decrease its cost (increase its utility) by changing his strategy unilaterally

Every symmetric game with two strategies has an equilibrium in pure strategies

 Cheng, S.G. et al.: Notes on the equilibria in symmetric games, Proc. 6th Workshop On Game Theoretic And Decision Theoretic Agents (collocated with IEEE AAMAS). New York, USA (2004)



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Derivation of pure equilibrium states/strategies

• The action profile $\alpha = (\alpha_i, \alpha_{-i})$ is a pure Nash equilibrium if for all $i \in N$

$$a_i \in \arg\min_{a_i' \in A_i} (c_i^N(a_i', a_{-i}))$$

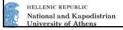
■ How to find EQ:

Identify the conditions on the number of competing agents that break the equilibrium definition and reverse them.

A driver is motivated to change his action in the following circumstances

when
$$a_i = pl$$
 and $w_{osp}(N_{osp}(a) + 1) < c_{pl}$ (*)

when
$$a_i = osp$$
 and $w_{osp}(N_{osp}(a)) > c_{pl}$ (**)



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Derivation of pure equilibrium states/strategies

<u>Lemma</u>: a player is motivated to change his action α_i as follows

•
$$a_i = \operatorname{pl}$$
 and $(a) N_{\operatorname{osp}}(a) < R \le N$ or

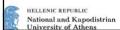
(b)
$$R \le N_{\text{osp}}(a) < N_0 - 1 \le N$$
 or

(c)
$$N_{\rm osp}(a) < N \le R$$

•
$$a_i = \operatorname{osp}$$
 and $R < N_0 < N_{\operatorname{osp}}(a) \le N$

where
$$N_0 = R(\gamma - 1)/\delta \in \mathbb{R}$$
.

<u>Hint</u>: for (b), require that (*) holds and for the second bullet require that (**) holds and get the condition on $N_{osp}(a)$

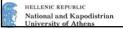


Προηγμένες Δικτυακές Τεχνολογίες - 2025

Pure equilibrium strategies

The strategic parking spot selection game has the following EQ profiles

- a) for $N \leq N_0$, a unique NE profile a^* with $N_{\rm osp}(a^*) = N_{\rm osp}^{\rm NE,\,1} = N;$
- b.1) for $N > N_0$ and $N_0 \in (R, N) \setminus \mathbb{N}^*$, $\binom{N}{\lfloor N_0 \rfloor}$ NE profiles a' with $N_{\text{osp}}(a') = N_{\text{osp}}^{\text{NE}, 2} = \lfloor N_0 \rfloor$;
- b.2) for $N > N_0$ and $N_0 \in [R+1,N] \cap \mathbb{N}^*$, $\binom{N}{N_0}$ NE profiles a' with $N_{\mathrm{osp}}(a') = N_{\mathrm{osp}}^{\mathrm{NE},\,2} = N_0$ and $\binom{N}{N_0-1}$ NE profiles a^* with $N_{\mathrm{osp}}(a^*) = N_{\mathrm{osp}}^{\mathrm{NE},\,3} = N_0 1$.



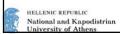
Προηγμένες Δικτυακές Τεχνολογίες - 2025

33

Pure Nash EQ strategies

The strategic resource selection game has the following EQ profiles

# drivers, N	#competing drivers for OSP (public) parking under NE
$\leq \frac{R(\gamma-1)}{\delta}$	N
$> \frac{R(\gamma-1)}{\delta}$	$\frac{R(\gamma-1)}{\delta}$



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Efficiency of pure equilibrium states/strategies

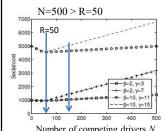
Efficiency Metric**: The Price of Anarchy (PoA \geq 1)

= [social cost under EQ state] / [optimal social cost]

Social cost under action profile α :

$$C(N_{osp}(a)) = \sum_{i=1}^{N} c_i^N(a) = \begin{cases} c_{osp,s}(N\beta - N_{osp}(a)(\beta - 1)), & \text{if } N_{osp}(a) \le R \\ c_{osp,s}(N_{osp}(a)\delta - R(\gamma - 1) + \beta N), & \text{if } R < N_{osp}(a) \le N \end{cases}$$

Use under a* : if N $\leq N_{\theta}$ use $N_{osp}(\alpha^*)$ = N ; if N $\geq N_{\theta}$ use $N_{osp}(\alpha^*)$ = $N_{osp}^{\rm NE,\,2} = \lfloor N_0 \rfloor$



cost minimized at $N_{osp} = R = \underline{\bf 50}$

Social cost increases as N_{osp} moves away from R (lack of coordination penalty)

Cost at EQ: $N_{osp} = N_{osp}^{NE, 2} = \lfloor N_0 \rfloor = \rfloor R(\gamma - 1)/\delta \rfloor$ = $R*(15-1)/(15-10) = R*2.8 = \underline{140}$ (<N=500)

Number of competing drivers N_{osp} lack of coordination penalty of 500

4000

** E. Koutsoupias, C. H. Papadimitriou, "Worst-case equilibria", Computer Science Review, 2009.

National and Kapodistrian University of Athens

Προηγμένες Δικτυακές Τεχνολογίες - 2025

Efficiency of pure equilibrium states/strategies

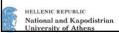
Optimal social cost: (exactly min{R,N} players compete – no fail and all osp spots used)

$$C_{opt} = \sum_{i=1}^{N} c_i^N(a_{opt}) = c_{osp,s}[min(N,R) + \beta \cdot max(0,N-R)]$$

Price of Anarchy

$$\text{PoA} = \left\{ \begin{array}{l} \frac{\gamma N - (\gamma - 1) \min(N, R)}{\min(N, R) + \beta \max(0, N - R)}, & if \ N_0 \geq N \\ \frac{\lfloor N_0 \rfloor \delta - R(\gamma - 1) + \beta N}{R + \beta (N - R)} =, & if \ N_0 < N \end{array} \right.$$

 $PoA \le 1 / [1-R/N]$, for N>R



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Efficiency of pure equilibrium states/strategies

$c = [\min(N, R) + \beta \max(0, N - R)]$	Optimal social cost, Copt
cosp,s[mm(1,1,1) + p max(0,1,1 1)]	$c_{osp,s}[\min(N,R) + \beta \max(0,N-R)]$

# drivers, N	Social cost in EQ, Ceq
$\leq \frac{R(\gamma-1)}{\delta}$	$c_{osp,s}[N\gamma - \min(N,R)(\gamma - 1)]$
$> \frac{R(\gamma-1)}{\delta}$	$c_{osp,s}oldsymbol{eta}N$

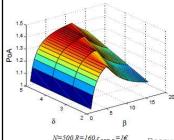
N compete

 $R(\gamma\text{-}1)/\delta$ compete

R competing drivers

PoA>1: due to competition and paying the (lack-of-coordination) cruising cost

Pricing (β) and failure cost/ overhead (δ) shaping guidelines



β	δ	PoA→ 1
$\geq \frac{\delta(N-R)+R}{R}$	> 0	Υβ
$<\frac{\mathcal{S}(N-R)+R}{R}$	> 0	√β
>1	$\leq \frac{R(\beta-1)}{N-R}$	√δ

Προηγμένες Δικτυακές Τεχνολογίες - 2025

The (strategic) mixed-action selection game

Practical strategies for real systems / a mixed action $p = (p_{osp}, p_{pl})$

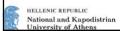
Expected cost for player choosing osp or pl, respectively, when all N-1 play according to the mixed-action p

$$c_i^N(osp, p) = \sum_{N_{osp}=0}^{N-1} w_{osp}(N_{osp} + 1)B(N_{osp}; N - 1, p_{osp})$$

$$c_i^N(pl,p) = c_{pl}$$

Expected cost for symmetric profile (all play according to the mixed-action p)

$$c_i^N(p, p) = p_{osp} \cdot c_i^N(osp, p) + p_{pl} \cdot c_i^N(pl, p)$$



Προηγμένες Δικτυακές Τεχνολογίες - 2025

The (strategic) mixed-action selection game

• Existence of Mixed-action Nash Equilibrium:

Every symmetric game with more than two players and increasing cost functions (of the number of players) has a unique mixed-action equilibrium

- Ashlagi, I., Monderer, D., Tennenholtz, M.: Resource selection games with unknown number of players. In: Proc. AAMAS '06. Hakodate, Japan (2006)
- The Mixed-action Nash Equilibrium state:

The strategic parking spot selection game has a unique mixed-action NE

$$p^{\text{NE}} = \left(p_{\text{osp}}^{\text{NE}}, \, p_{\text{pl}}^{\text{NE}}\right), \quad \text{where } p_{\text{osp}}^{\text{NE}} = 1 \text{ if } N \leq N_0 \text{ and } p_{\text{osp}}^{\text{NE}} = N_0/N \text{ if } N > N_0,$$
 with $p_{\text{osp}}^{\text{NE}} + p_{\text{pl}}^{\text{NE}} = 1 \text{ and } N_0 \in \mathbb{R}.$

Sketch of proof: Set the requirement to be fulfilled by the profiles at EQ:

$$c_i^N(osp,p^{NE}) = c_i^N(pl,p^{NE})$$



Προηγμένες Δικτυακές Τεχνολογίες - 2025

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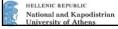
Equilibrium states/strategies (pure / mixed-action strategies

Pure NE

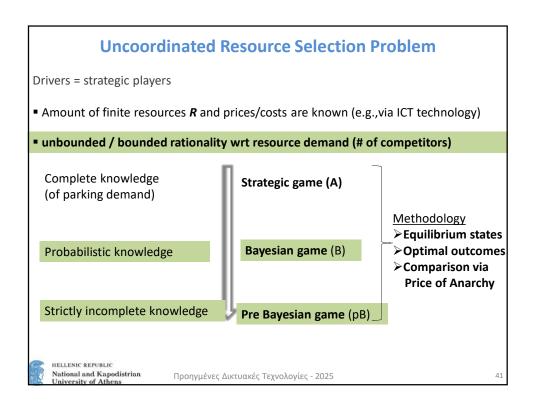
# drivers, N	#competing drivers for public parking space
$\leq \frac{R(\gamma-1)}{\delta}$	N
$> \frac{R(\gamma-1)}{\delta}$	$\frac{R(\gamma-1)}{\delta}$

Symmetric mixed-action NE

# drivers, N	Probability of competing for public parking space
$\leq \frac{R(\gamma - 1)}{\delta}$	1
$> \frac{R(\gamma-1)}{\delta}$	$\frac{R(\gamma-1)}{\delta N}$

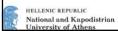


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Decision-making under demand knowledge constraints (a type of bounded rationality)

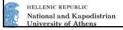
- The Bayesian model: probabilistic (demand) information
 - Players know
 - a) probability for a player to be active (interested in parking)
 - b) total number of players
- The pre Bayesian model: strictly incomplete (demand) information
 - Players know their total number (*upper bound on competitors*)



Προηγμένες Δικτυακές Τεχνολογίες - 2025

A Bayesian parking spot selection game

- $A_i = \{osp, pl, \emptyset\}$ is the set of potential actions for each driver $i \in \mathcal{N}$;
- $\Theta_i = \{0,1\}$ is the set of types for each driver $i \in \mathcal{N}$, where 1 (0) stands for active (inactive) drivers;
- $S_i: \Theta_i \to A_i$ is the set of possible strategies for each driver $i \in \mathcal{N}$;
- $c_i^{N_B}(s(\vartheta),\vartheta)$ is the cost functions for each driver $i\in\mathcal{N}$, for every type profile $\vartheta\in\times_{k=1}^N\Theta_k$ and strategy profile $s(\vartheta)\in\times_{k=1}^NS_k$, that are functions of $w_{osp}(\cdot)$ and $w_{pl}(\cdot)$, as defined for $\Gamma(N)$, and also written as $c_i^{N_B}(s(\vartheta),\vartheta)=c_i^{N_B}(s_i(\vartheta_i),s_{-i}(\vartheta_{-i}),\vartheta_i,\vartheta_{-i})$;
- p_{act} is the probability for a driver to be active.



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4

A Bayesian parking spot selection game

Equilibria:

For the game $\Gamma_B(N)$, the strategy profile $s' \in \times_{k=1}^N S_k(\vartheta_k = 1)$ is a Bayesian NE if, for all $i \in \mathcal{N}$ with $\vartheta_i = 1$ $s_i(\vartheta_i) \in \arg\min_{s' \in S_i} (c_i^{N_B}(s_i(\vartheta_i), s_{-i}(\vartheta_{-i}), \vartheta_i, \vartheta_{-i})) \qquad \text{or,}$

$$s_i(\vartheta_i) \in \arg\min_{s_i' \in S_i} \sum_{\vartheta_{-i}} f_{\Theta}(\vartheta_{-i}/\vartheta_i) c_i^{\sum_k \vartheta_k}(s_i', s_{-i}(\vartheta_{-i}), \vartheta_i, \vartheta_{-i})$$

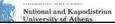
Derivation approach

$$c_{i}^{N_{B}}(pl,p) = c_{pl} \qquad c_{i}^{N_{B}}(osp,p) = \sum_{n_{act}=0}^{N-1} c_{i}^{n_{act}+1}(osp,p) B(n_{act};N-1,p_{act})$$

Equilibria:

The Bayesian parking spot selection game $\Gamma_B(N)$ has unique symmetric equilibrium profiles $p^{NE_B}=(p_{osp}^{NE_B},p_{pl}^{NE_B})$, with $p_{osp}^{NE_B}+p_{pl}^{NE_B}=1$. More specifically:

- a unique pure (Bayesian Nash) equilibrium with $p_{osp}^{NE_B}=1$, if $p_{act}<\frac{N_0}{N}$,
- a unique symmetric mixed-action Bayesian Nash equilibrium with $p_{osp}^{NE_B} = \frac{N_0}{Np_{act}}$, if $p_{act} \ge \min(\frac{N_0}{N}, 1)$, where $N_0 \in \mathbb{R}$.



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Equilibrium states under probabilistic knowledge(Bayesian Game)

Symmetric mixed-action Bayesian Nash equilibria

Activation probability, pact	(symmetric) Probability of competing
$<\frac{R(\gamma-1)}{\delta N}$	1
$\geq \min\left(\frac{R(\gamma-1)}{\delta N},1\right)$	$\frac{R(\gamma-1)}{\delta Np_{act}}$

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45

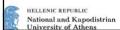
Equilibrium states under strict uncertainty (pre-Bayesian model)

Knowledge of an upper bound on demand, N

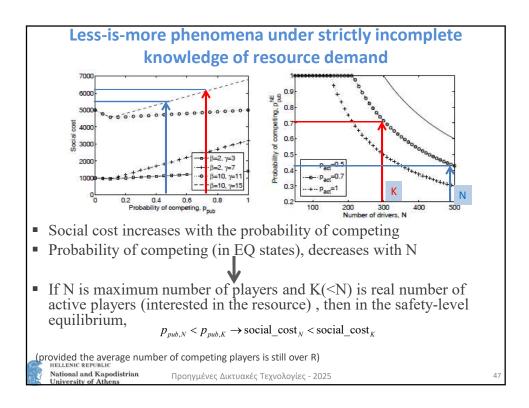
For the pre-Bayesian resource selection game it holds that

Symmetric mixed-action safety-level equilibrium (playing to min worst cost) \uparrow

Symmetric mixed-action equilibrium of the strategic game $\Gamma(N)$ with N players



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EQ strategies for Strategic, Bayesian and pre-Bayesian Resource Selection Game

 ${\bf TABLE~I}\\ {\bf EQUILIBRIUM~STRATEGIES~FOR~THE~STRATEGIC,~BAYESIAN~AND~PRE-BAYESIAN~PARKING~SPOT~SELECTION~GAME}$

Condition	Equilibrium type	Equilibrium expression
$N \leq N_0,N_0 \in \mathbb{R}$	pure Nash Eq	$N_{osp}^{NE} = N$
$N > N_0, N_0 \in (R, N) \backslash \mathbb{N}^*$	pure Nash Eq	$N_{osp}^{NE} = \lfloor N_0 floor$
$N > N_0, N_0 \in [R+1, N] \cap \mathbb{N}^*$	pure Nash Eq	$N_{osp}^{NE} = N_0, \ N_{osp}^{NE} = N_0 - 1$
$N>N_0,\ N_0\in\mathbb{R}$	mixed-action Nash Eq	$p_{osp}^{NE}=rac{N_0}{N}$
Bayesia	n Parking Spot Selection Game, Γ	$_{3}(N)$
Condition	Equilibrium type	Equilibrium expression
$p_{act} < rac{N_0}{N}, N_0 \in \mathbb{R}$	pure Bayesian Nash Eq	$p_{osp}^{NE_B} = 1$
$p_{act} \geq \min(\frac{N_0}{N}, 1), \ N_0 \in \mathbb{R}$	mixed-action Bayesian Nash Eq	$p_{osp}^{NE_B} = rac{N_0}{Np_{act}}$
pre-Bayes	ian Parking Spot Selection Game, l	$\Gamma_{pB}(N)$
Condition	Equilibrium type	Equilibrium expression
$N \leq N_0,N_0 \in \mathbb{R}$	pure safety-level Eq	$p_{osp}^{NE_{pB}} = 1$
$N>N_0,N_0\in\mathbb{R}$	mixed-action safety-level Eq	$p_{osp}^{NE_{pB}} = \frac{N_0}{N}$

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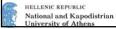
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CASE B:

- Bounded Rationality: (human-related) computational limits cognitive biases
 - inputs from behavioral economics, cognitive psychology, etc
 - models for bounded rationality and assessment of its impact

Classical Expected Utility Maximization Framework not adequate

E. Kokolaki, M. Karaliopoulos, I Stavrakakis, "On the human-driven decision-making process in competitive environments", Internet Science Conference, 2013



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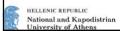
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RECALL: Decisions under Full Rationality Expected Utility Theory framework

- Strategic agents with perfect information, without behavioral biases, aiming at maximizing own welfare
 - -quantified by the expected gain/cost of their actions through EUT framework
- Expected Utility Theory (EUT) framework
 - -Expected utility of a lottery equals the sum of the utilities of the lottery outcomes, $U(x_i)$, times their probabilities of the outcomes, $p(x_i)$

$$EU = \sum_{x_i} p(x_i)U(x_i), x_i$$
: choice / alternativ e.

- Nash equilibrium
 - -captures best response in terms of expected utility maximization



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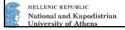
Deviations from Full Rationality

☐ (Cumulative) Prospect Theory

- maintains most of the concepts/assumptions of EUT
- manipulates both utility measures and prob. to account for biases against risk

☐ Alternative decision-making models & Equilibrium (EQ) concepts (Quantal Response, Rosenthal)

- Use probabilistic choice models to capture any unobserved and omitted elements, estimation/computational errors, individual's mood, perceptual variations or cognitive biases
- In line with the fact that individuals are more likely to make better choices than worse choices, but do not necessarily make the very best choice



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51

Prospect Theory motivation (1)

In several choice problems, individuals' preferences systematically violate EUT !!!

Allais' paradox (indication that people assess utilities and probabilities of outcomes differently from what full rationality/ EUT predicts => contradictions under EUT formulation)

```
PROBLEM 1: Choose between

A: 2,500 with probability .33, B: 2,400 with certainty
2,400 with probability .66,
0 with probability .01;

N = 72 [18]

Percentage of responses
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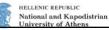
 $EUT_{prospectB} \geq EUT_{prospectA} \Leftrightarrow \text{u(2400)} > .33\text{u(2500)} + .66\text{u(2400)} \Leftrightarrow \textbf{.34u(2400)} > \textbf{.33u(2500)} + .66\text{u(2400)} \Leftrightarrow \textbf{.34u(2400)} > \textbf{.33u(2500)} + .66\text{u(2400)} \Leftrightarrow \textbf{.34u(2400)} > \textbf{.34u(2400$

PROBLEM 2: Choose between

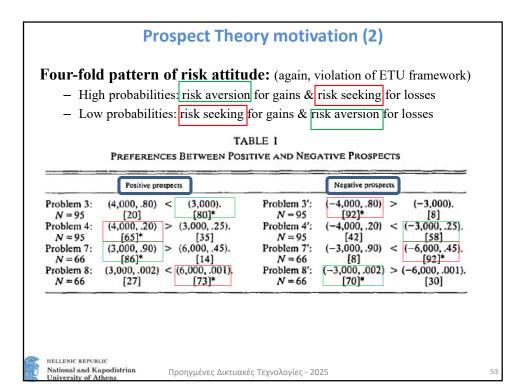
C: 2,500 with probability .33, D: 2,400 with probability .34,

0 with probability .67; 0 with probability .66. N = 72 [83]* [17]

 $EUT_{prospectC} > EUT_{prospectD} \Leftrightarrow .33u(2500) > .34u(2400)$



Προηγμένες Δικτυακές Τεχνολογίες - 2025



Prospect Theory formulation (Kahneman & Tversky, 1979)*

Defines prospects

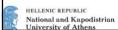
Prospect : $(x_1, p_1; x_2, p_2; ...; x_n, p_n)$

Desirability of a prospect is quantified through generalization of the utility functions and their weighting through weighting functions $\pi(p_i)$

$$U = \sum_{i=1}^{n} \prod(p_i) v(x_i)$$

· Decision maker is still a utility maximizer

* Daniel Kahneman and Amos Tversky, "Prospect Theory: An Analysis of Decision under Risk", *Econometrica*, *47*(2), pp. 263-291, March 1979



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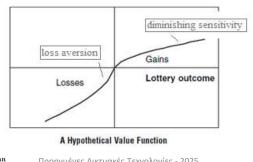
The Prospect Theory (PT) model

Diminishing sensitivity:

Impact of a chance diminishes with distance from Reference Point (a gain from reference 50 to 100 is less valuable than from 0 to 50) (people are more sensitive to extreme outcomes and less to intermediate ones)

Loss aversion:

Curve is steeper for losses than for gains (a high loss hurts more than a high pleasure)



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The Cumulative PT (CPT) model (Tversky & Kahneman, 1992)

Prospect: $(x_1, p_1; x_2, p_2; ...; x_n, p_n)$

CPT fixes some (experimentally observed) inconsistencies of PT

Modifies the probability weighting functions

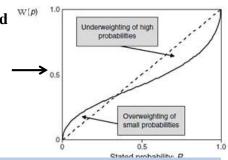
- PT: transforms probabilities of separate outcomes
- CPT: transforms probabilities of {an outcome or anything better (or worse) than that}

Desirability of a prospect quantified through generalizing:

A. Outcome probabilities {p_i} via decision weights $\{w(p_i)\}$

> (diminishing sensitivity: more sensitive around 0 and 1, and less in the middle)

B. Outcome values $\{x_i\}$ via utility functions $\{U(x_i)\}$



A. Tversky and D. Kahneman, "Advances in prospect theory: cumulative representation of uncertainty", Journal of Risk and Uncertainty, 5, 1992.

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The Cumulative PT (CPT) model (Tversky & Kahneman, 1992)

Desirability of a prospect quantified through generalizing:

- A. Outcome probabilities $\{p_i\}$ via decision weights $\{w(p_i)\}$
- B. Outcome values $\{x_i\}$ via utility functions $\{U(x_i)\}$

CPT value of the prospect $(x_1, p_1; ...; x_n, p_n)$:

$$U = \sum_{i=1}^{k} \pi_{i}^{-} v(x_{i}) + \sum_{i=k+1}^{n} \pi_{i}^{+} v(x_{i}) , \qquad x_{1} \leq \ldots \leq x_{k} \leq 0 \leq x_{k+1} \leq \ldots \leq x_{n}.$$

where the decision weights (i.e. the numbers π_i^-, π_i^+) are defined by:

$$\pi_1^- = w^-(p_1), \qquad \pi_i^- = w^-(p_1 + \dots + p_i) - w^-(p_1 + \dots + p_{i-1}) \qquad 2 \le i \le k$$

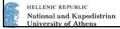
$$\pi_n^+ = w^+(p_n), \quad \pi_i^+ = w^+(p_i + \dots + p_n) - w^+(p_{i+1} + \dots + p_n) \quad k+1 \le i \le n-1$$

$$\pi_n^+ = w^+(p_n), \qquad \pi_i^+ = w^+(p_i + \dots + p_n) - w^+(p_{i+1} + \dots + p_n) \qquad k+1 \leqslant i \leqslant n-1$$

$$v(x) = \begin{cases} x^{\alpha} & \text{if } x \geq 0 & w^+(p) = p^{\gamma}/[p^{\gamma} + (1-p)^{\gamma}]^{1/\gamma} & w^+(0) = w^-(0) = 0 \\ -\lambda(-x)^{\beta} & \text{if } x < 0, \quad w^-(p) = p^{\delta}/[p^{\delta} + (1-p)^{\delta}]^{1/\delta}, \quad w^+(1) = w^-(1) = 1. \end{cases}$$
The values that had fit the constructed results of Kaharana 2. Targing and
$$\pi_n^+ = w^+(p_n), \qquad \pi_i^+ = w^+(p_i + \dots + p_n) - w^+(p_{i+1} + \dots + p_n) \qquad k+1 \leqslant i \leqslant n-1$$

$$-\lambda(-x)^{\beta} & \text{if } x < 0, \quad w^-(p) = p^{\delta}/[p^{\delta} + (1-p)^{\delta}]^{1/\delta}, \quad w^+(1) = w^-(1) = 1.$$

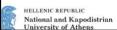
The values that best fit the experimental results of Kahneman & Tversky are: $\left\{ \begin{array}{l} \alpha=\beta=0.88;\ \lambda=2.25;\\ \gamma=0.61;\ \delta=0.69 \end{array} \right.$



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Equilibrium concepts over Cumulative Prospect Thy

- Nash Equilibrium state: no player can increase his/her utility (expected utility value) by changing his/her strategy unilaterally
 - Mixed-action profiles:
 - EU value of strategy I = EU value of strategy 2
- CPT Equilibrium state: no player can increase his/her utility (cumulative prospect value) by changing his/her strategy unilaterally
 - -Mixed-action profiles:
 - •CPT value of prospect 1 = CPT value of prospect 2



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Applying CPT to Resource Selection Problem

Two alternatives/prospects: \boldsymbol{l} (limited resource, osp) and \boldsymbol{u} (unlimited resource, pl)

CPT value for prospects l and, u respectively:

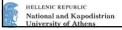
$$CPT_l = \sum_{n=1}^{N} \pi_n^- u(g_l(n))$$
 $CPT_u = u(c_u)$

With the cost of selecting prospect l when k others do the same given by

$$g_l(k) = min(1, R/k)c_{l,s} + (1 - min(1, R/k))c_{l,f}$$

and the probabilities p_k for this are Binomial (N, p_l^{CPT})

Both prospects consist of negative outcomes / costs

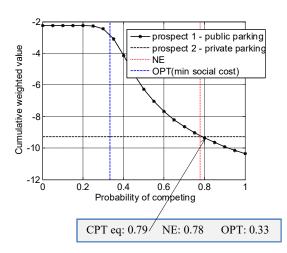


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59

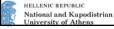
Applying CPT to Resource Selection Problem

- N= 150, R=50
- $c_{osp,s}=1$
- $\beta = 5 \ (c_{pl} = 5)$
- γ=8
- $\delta = 3$

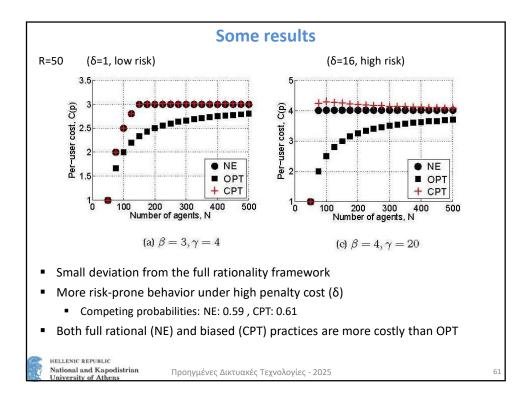


CPT equilibrium condition :

CPT value of prospect 1 = CPT value of prospect 2

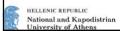


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Deviations from Full Rationality

- ☐ Alternative decision-making models & Equilibrium (EQ) concepts (Quantal Response, Rosenthal)
 - Use probabilistic choice models to capture any unobserved and omitted elements, estimation/computational errors, individual's mood, perceptual variations or cognitive biases
 - In line with the fact that individuals are more likely to make better choices than worse choices, but do not necessarily make the very best choice
 - > Due to noise/disturbances in their anticipation of exact choices' payoffs



Προηγμένες Δικτυακές Τεχνολογίες - 2025

Quantal response equilibrium (McKelvey & Palfrey, 1995)

Introduce some randomness in the decision-making process to capture people's inability to play always the strategy that maximizes the expected utility

- "Choices are made with probabilities that are monotone in their expected payoffs"
- Logit QRE => disturbances/errors follow extreme value distribution (smaller mistakes are more likely to occur that more serious ones)

$$p(r_1) = \frac{e^{-\lambda EU(r_1)}}{e^{-\lambda EU(r_1)} + e^{-\lambda EU(r_2)}}$$

$$p(r_1) = 1 - p(r_2)$$

 $\lambda \in [0, \infty]$: rationality control parameter

 $\lambda \rightarrow 0$: random decision

 $\lambda \rightarrow \infty$: full rationality (Nash EQ) [cost differences (i.e., EU(.)) are emphasized more through a more responsive distribution to cost changes, in line with the more emphasis in differences expected by a more rational decision-

R. McKelvey and T. Palfrey, "Quantal response equilibria for normal form games", Games and Economic Behavior, 1995.



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Quantal response equilibrium for Resource Selection Problen

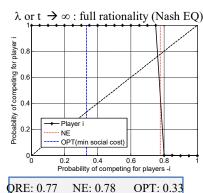
$$p_l^{QRE} = \frac{e^{-tc(l,p^{QRE})}}{e^{-tc(l,p^{QRE})} + e^{-tc(u,p^{QRE})}} \qquad c(l,p) = \sum_{n=0}^{N-1} g_l(n+1)B(n;N-1,p_l)$$

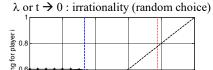
$$c(l,p) = \sum_{n=0}^{N-1} g_l(n+1)B(n;N-1,p_l)$$

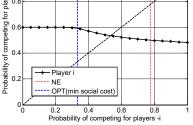
$$p^{QRE} = (p_l^{QRE}, p_u^{QRE}), p_u^{QRE} = 1 - p_l^{QRE}$$

$$c(u, p) = c_u$$

At EQ, prob of competing of player i equals the belief for the way (i.e. prob) the others play which is used in calculating c(l,p)







QRE: 0.55 NE: 0.78

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Rosenthal equilibrium (Rosenthal, 1989)

"The difference in probabilities with which two actions are played equals a parameter *t* multiplied by the difference of the corresponding expected costs"

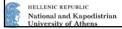
$$p(r_1) - p(r_2) = t(EU(r_1) - EU(r_2))$$

$$p(r_1) = 1 - p(r_2)$$

 $t \in [0,\infty]$: rationality control parameter

 $t \rightarrow \infty$: full rationality

R. Rosenthal, "A bounded-rationality approach to the study of noncooperative games", Int. J. Game Theory, 1989.



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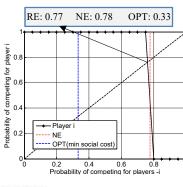
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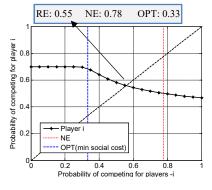
Rosenthal equilibrium for Resource Selection Problem

$$egin{aligned} p_{l}^{RE} - p_{u}^{RE} &= -t(c(l, p^{RE}) - c(u, p^{RE})) \ p_{u}^{RE} &= (p_{l}^{RE}, p_{u}^{RE}), \ p_{u}^{RE} &= 1 - p_{l}^{RE} \end{aligned}$$

 $t \rightarrow \infty$: full rationality (NE)

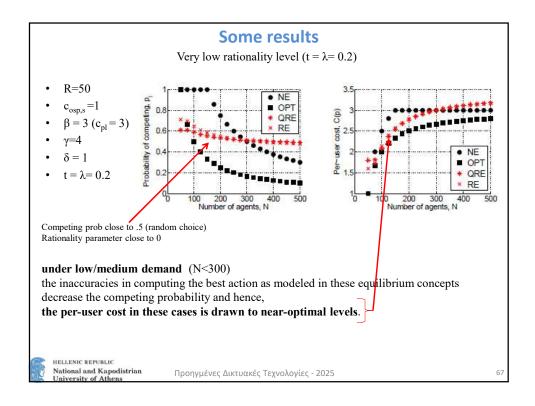
 $t \rightarrow 0$: no rationality (random choice)

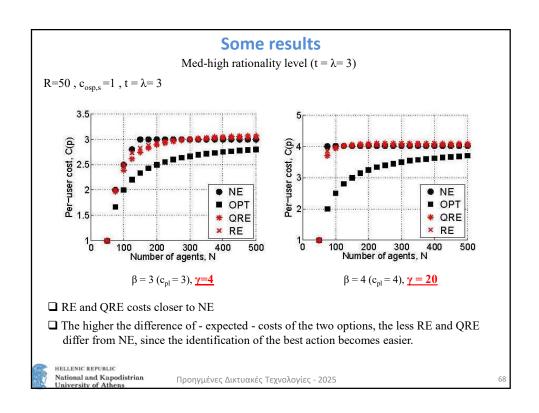


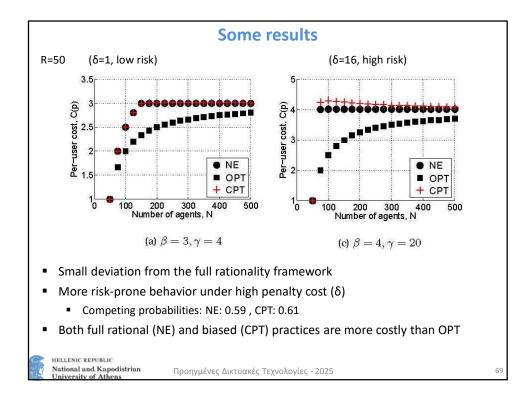


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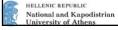






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- ☐ Alternative decision-making models & Equilibrium (EQ) concepts (Quantal Response, Rosenthal)
 - Use probabilistic choice models to capture any unobserved and omitted elements, estimation/computational errors, individual's mood, perceptual variations or cognitive biases
 - In line with the fact that individuals are more likely to make better choices than worse choices, but do not necessarily make the very best choice
- **□** Heuristics
 - ☐ Fast and frugal reasoning solutions / decisions
 - ☐ Emphasis on cognitive processes underlying decisions
 - ☐ Satisfying instead of maximization of expected utilities (Simon 1955)



Προηγμένες Δικτυακές Τεχνολογίες - 2025

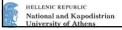
Heuristic strategy for the resource selection problem

Satisficing: instead of computing/comparing expected costs, it estimates the probability to hit an empty public spot and plays according to this.

Confidence heuristic rule: "risk competing for resource r₁ according to the probability of winning one of the R resources"

- Under the belief that other players think the same way
- Fixed-point equation

$$p(r_1) = \Pr(\#competitors < R) = \sum_{j=0}^{R-1} Bin(j, N-1; p(r_1))$$



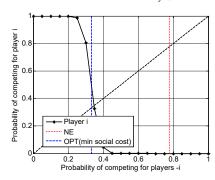
Προηγμένες Δικτυακές Τεχνολογίες - 2025

71

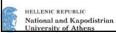
Heuristic strategy for the resource selection problem

• *Confidence* heuristic: "risk for public parking space according to the probability of winning public parking space"

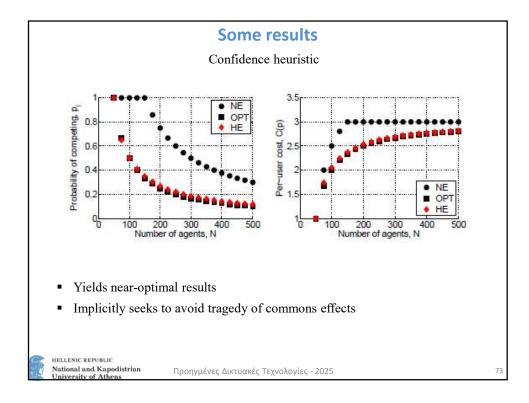
$$p_{osp} = \Pr(\#competitors < R) = \sum_{i=0}^{R-1} Bin(j; N-1, p_{osp})$$



HE = 0.34, NE = 0.78, OPT = 0.33



Προηγμένες Δικτυακές Τεχνολογίες - 2025



Conclusions for the models of bounded rationality

Formulation of the parking spot selection application drawing on models from behavioral economics and cognitive psychology

• (Cumulative) Prospect Theory:

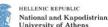
- The decision maker is still a utility maximizer
- The desirability of outcomes is expressed through the transformed probabilities of them
- Small deviation from the full rationality framework

Alternative equilibrium concepts (Quantal Response Equilibrium, Rosenthal equilibrium)

- The decision maker is a satisfizer
- Symmetric mixed-action equilibria as fixed-point solutions
- A degree of freedom quantifies the rationality in the model (convergence to the Nash equilibrium as the free parameter goes to infinity)

Heuristics

- The decision maker is a satisfizer
- Confidence heuristic: fast and frugal reasoning solution that is shown to yield nearoptimal results within the particular application concept

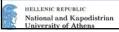


Προηγμένες Δικτυακές Τεχνολογίες - 2025

Real Players vs Nash EQ vs Rosenthal EQ vs Random vs Optimal

An experimental comparative study

María Pereda, Juan Ozaita, Ioannis Stavrakakis, and Angel Sanchez, "Competing for congestible goods: experimental evidence on parking choice", Scientific Reports (a Nature Research journal), 2020.



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75

Nash EQ vs Rosenthal EQ vs optimal

An experimental study

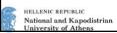
N=20 players with N known to all.

Choices: a yellow lot with only S {5, 10} slots, or a blue lot with unlimited capacity.

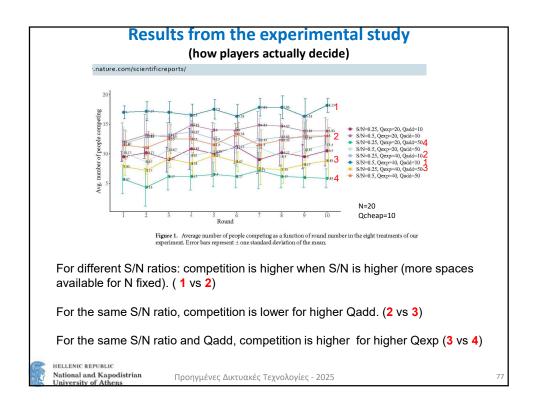
At the end of each round, participants had an associated payoff, resulting from subtracting their decision costs from their initial endowed 100 points.

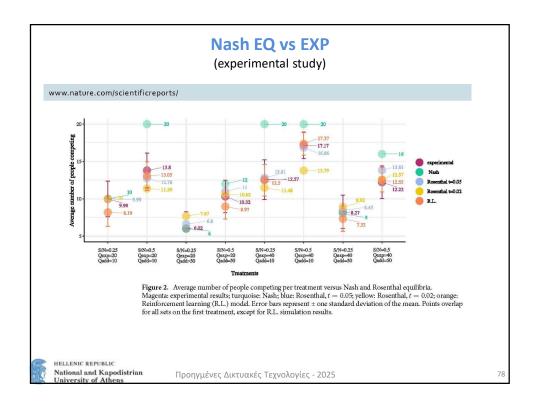
After each round, participants were reminded of:

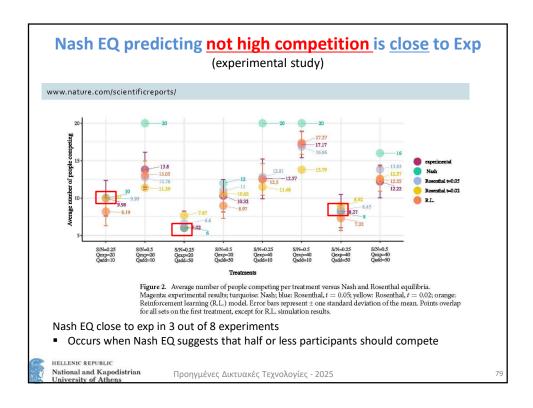
- their previous decision
- the type of slot they ended up using,
- their payoff for that round,
- the number of yellow slots occupied.

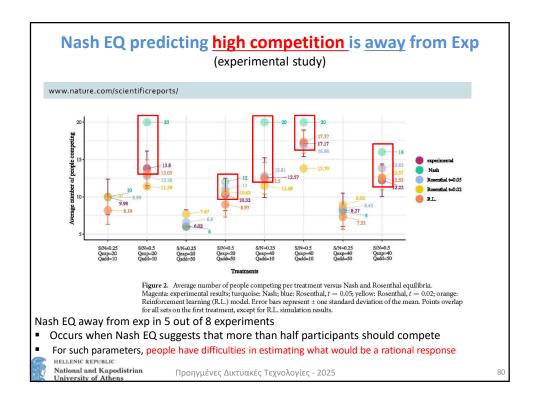


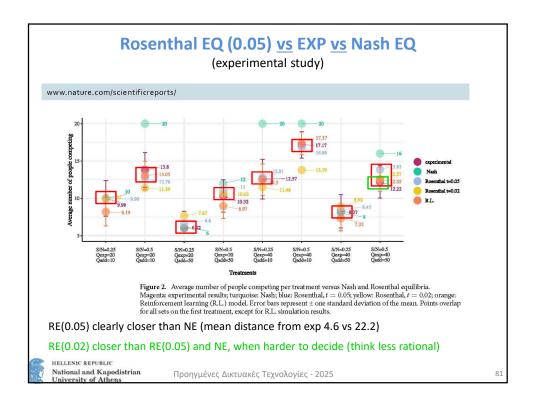
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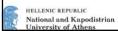




Nash EQ / Rosenthal EQ / EXP / vs RANDOM / OPTIMAL

An experimental study

- ☐ RANDOM decision (0.5*N compete):
 - > EXP away from random choice (mean distance 21.5)
- ✓ Real users do not decide RANDOM
- ☐ OPTIMAL (coordinated): min {N, S} = S
 - > EXP and RE(0.05) are about equally close to the OPTIMAL
 - > Nash EQ furthest away from OPTIMAL
- ✓ Real users decide more effectively than the fully rational approach
- √ The bounded rationality of real users can be captured well by the Rosenthal model



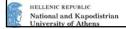
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Decision-Making for <u>Uncoordinated</u> User Access to (Distributed) <u>Limited Resources</u>

* A classical, old but ever modern and challenging problem *

Presented and discussed various decision models under <u>full rationality and bounded rationality (human-driven)</u>

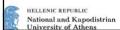
Tried to shed some light into the relevance and effectiveness of <u>alternative</u> <u>decision models and equilibrium concepts</u>, that could be considered when humans drive the decisions.



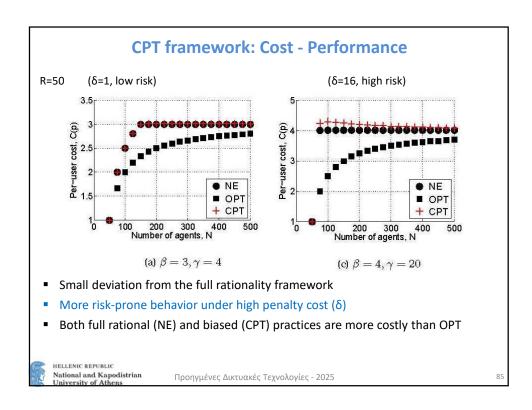
Προηγμένες Δικτυακές Τεχνολογίες - 2025

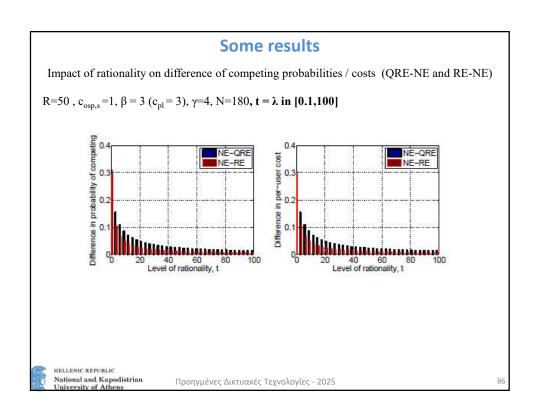
83

Extra slides start



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Congestion-cost-cutting Approaches In Resource-limited Competitive Environments

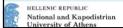


Coordinated Resource Allocation through Social Applications

Resolve competition through ICT and IoT



E. Kokolaki, M. Karaliopoulos, I. Stavrakakis, "Parking assisting applications: effectiveness and side-issues in managing public goods", 3rd AWARE workshop on Challenges for Achieving Self-Awareness in Autonomic Systems of the Self-Adaptive and Self-Organizing systems conference (SASO 2013), Sept. 9-13, 2013, Philadelphia



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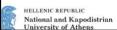
87

Decentralized Resource Allocation through Social Applications

User would subscribe to be able to more effectively use public resources in a competitive environment

Social apps for parking resources: (Sfpark, **Parking Defenders**, Parkomotivo)





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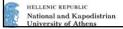
Some important questions

Are these apps effective and "fair" to their users?

- Do they require high user subscription to be effective?
- Do they treat equally similar users?
- Do they create a wealth (through coordination and congestion cost cutting) that rightfully distribute to their users?
 - Or simply benefit by eliminating competition by non-users (exclusion)?

What is the impact on non-users?

- Are non-users of these apps suffering substantially
 - (or almost excluded from) accessing public goods?

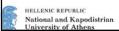


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89

3 Driver profiles

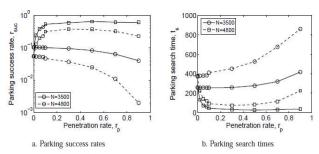
- <u>Traditional</u> user/driver (non-users of the application)
- App user/driver:
 - **Defender** (sharer, fully cooperative)
 - · Announces upcoming freed-up spot
 - Waits for selected app user to come
 - · Rates other defender upon parking
 - Earns credit and improves its ranking
 - <u>Seeker</u> (free rider not fully cooperative)
 - Rates defender upon parking



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Effectiveness of the social parking application: Very high parking demand (105%, 145%)

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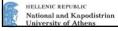
(the effective competition that non-users experience is mitigated)

Performance indices for Non-users (circles) and Sharers (squares) at high parking demand:

At high penetration rates:

HELLENIC REPUBLIC
National and Kapodistrian

- Non-user exclusion trends
- Sharer performance deterioration due to own competition



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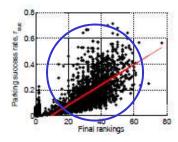
Effectiveness of the incentive mechanism and some concerns on its fairness

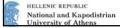
It is *effective*...

• higher success rates are coupled with higher rankings

But not fair ...

- it discriminates against identical users (=users with similar interests, needs and attitude towards cooperation)
- It induces rich-get-richer phenomena: winners in the initial competition round earn credits that offer them a competitive edge in the following rounds





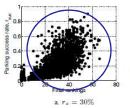
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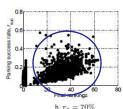
9:

Impact of Seekers on the incentive mechanism's fairness

Seekers tend to restore fairness

- \uparrow Seekers' portion \rightarrow \downarrow parking spot handovers \rightarrow
- ↓ opportunities for credit-building/emergence of high rankings





HELLENIC REPUBLIC
National and Kapodistrian
University of Athens

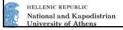
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Conclusions

Investigated effectiveness / appropriateness of distributed, social public resource (parking) management apps

Effective and mostly non-exclusive to non-users

- ➤ Users' improved performance is mostly due to the increased efficiency they generate in the parking process, rather than excluding traditional users from competing for the resources.
- ➤ Non-Users also benefit from the reduced anarchy and coordination that the App brings
- ➤ Incentive mechanism is effective
 - > But it induces rich-club phenomena and difficulties to newcomers
- Seekers (free-riders) seem to alleviate those problems



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95

SUMMARY of Resource Competition in a Highly Networked World of Humans and Things

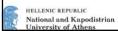
Motivation (environment – early study)

Decision-making in uncoordinated competitive environment - formulation

- ■rational case Price of Anarchy
- limited info case
- ■human driven case
 - Prospect Thy
 - alternative models
 - heuristics

Alternative, partially coordinated approaches

■ICT-supported distributed apps



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