

PART C

**DIGITAL MODULATION - DEMODULATION
(MODEM) TECHNIQUES**

P. Takis Mathiopoulos, Ph.D.

Professor (elected)

Department of Informatics and Telecommunications

University of Athens

Greece

e-mail: mathio@di.uoa.gr

- *Non-constant and Constant Envelope Signals*
- *Performance (noise, fading, interference, non-linearities)*
- *Techniques for Performance Improvement (Coherent and Non-coherent)*
- *Modem Implementation and Testing*

- PERFORMANCE

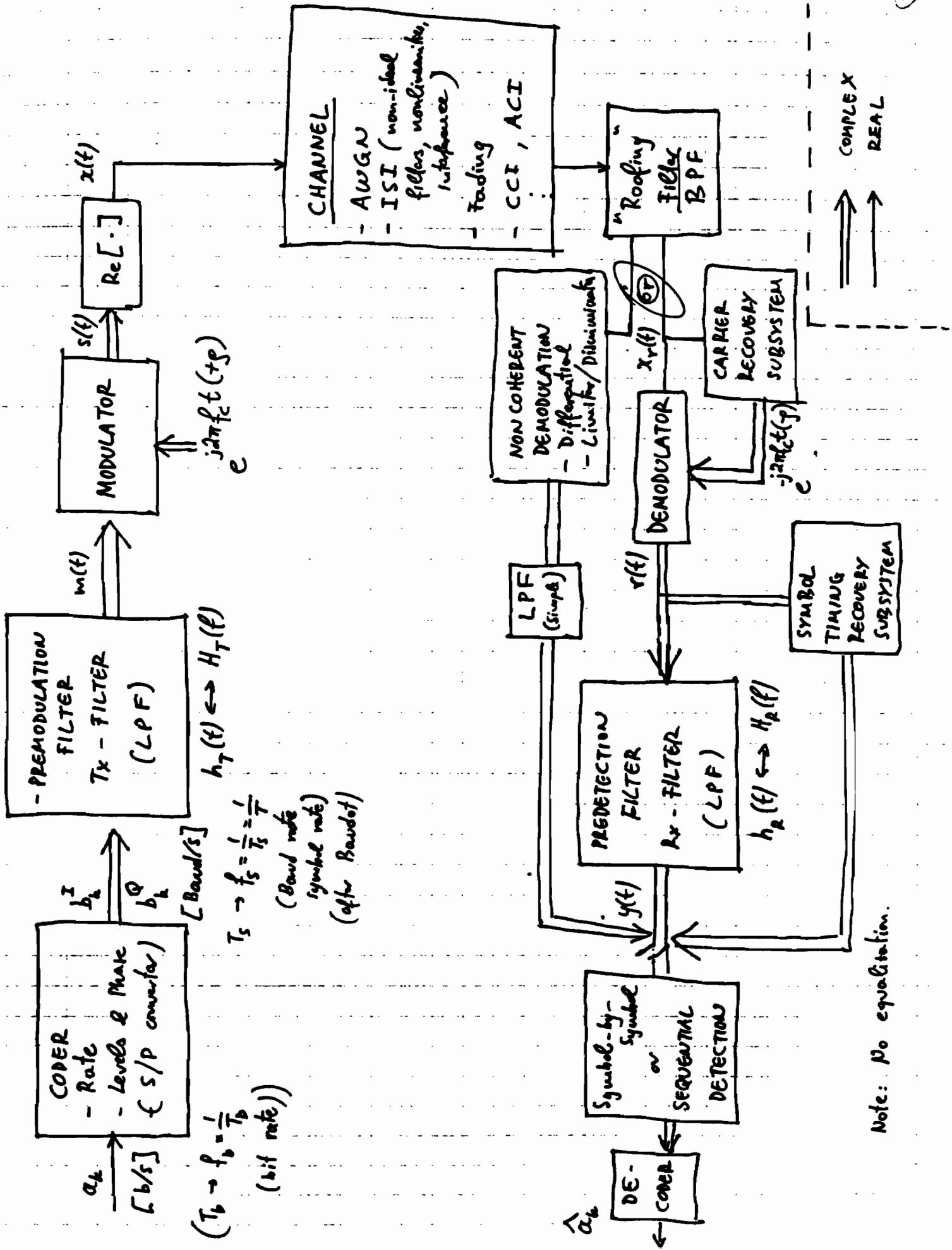
- noise
- fading
- co-channel interference
- nonlinearities

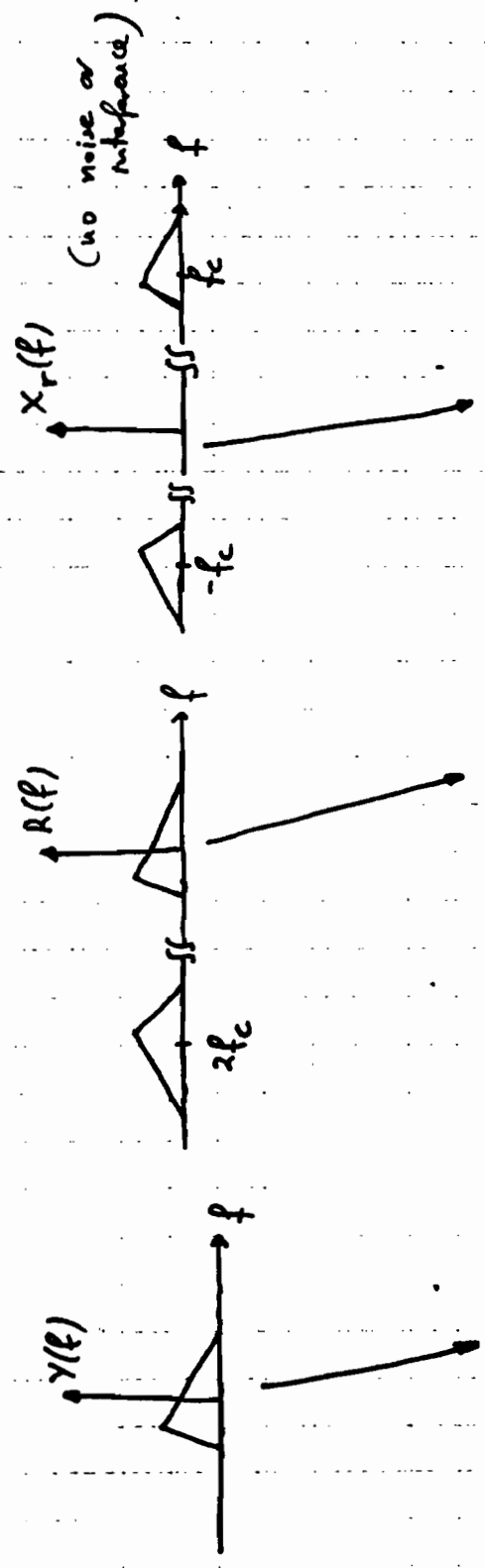
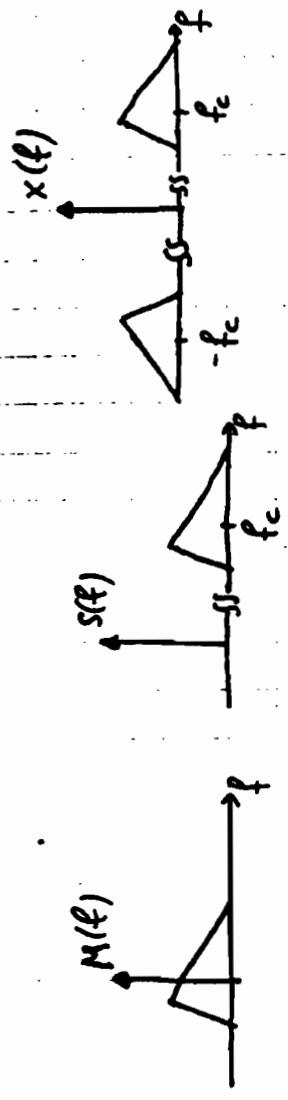
- TECHNIQUES TO IMPROVE THE PERFORMANCE

- Combat interference

- MODEM IMPLEMENTATION

- hardware
- software
- DSP





$$s(t) = e^{j2\pi f_c t} \sum_{k=-\infty}^{\infty} c_k h_T(t - kT)$$

$$c_k = b_k^I + j b_k^Q = a_k e^{j\theta_k} \rightarrow \begin{cases} b_k^I = a_k \cos \theta_k \\ b_k^Q = a_k \sin \theta_k \end{cases}$$

$$s(t) = [\cos(2\pi f_c t) + j \sin(2\pi f_c t)] \sum_k [b_k^I + j b_k^Q] h_T(t - kT)$$

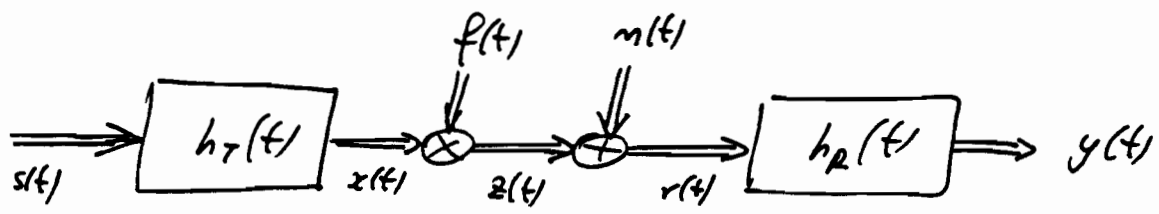
$$= \cos(2\pi f_c t) \sum_k b_k^I h_T(t - kT) - \sin(2\pi f_c t) \sum_k b_k^Q h_T(t - kT)$$

$$+ j [\cos(2\pi f_c t) \sum_k b_k^Q h_T(t - kT) + \sin(2\pi f_c t) \sum_k b_k^I h_T(t - kT)]$$

$$\therefore \boxed{x(t) = \text{Re}[s(t)] = \cos(2\pi f_c t) \sum_k b_k^I h_T(t - kT) - \sin(2\pi f_c t) \sum_k b_k^Q h_T(t - kT)}$$

Example for Computer Simulations [AWGN $n(t)$ fading $f(t)$]

1) Mathematical Model (usually convolution)

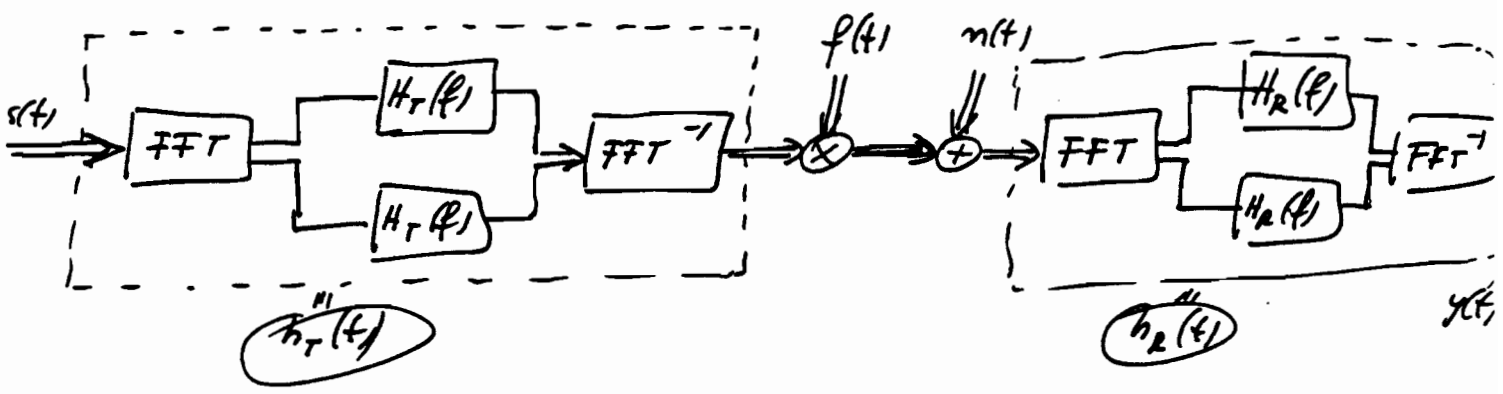


$$s(t) = s_I(t) + j s_Q(t) ; f(t) = f_I(t) + j f_Q(t) ; n(t) = n_I(t) + j n_Q(t)$$

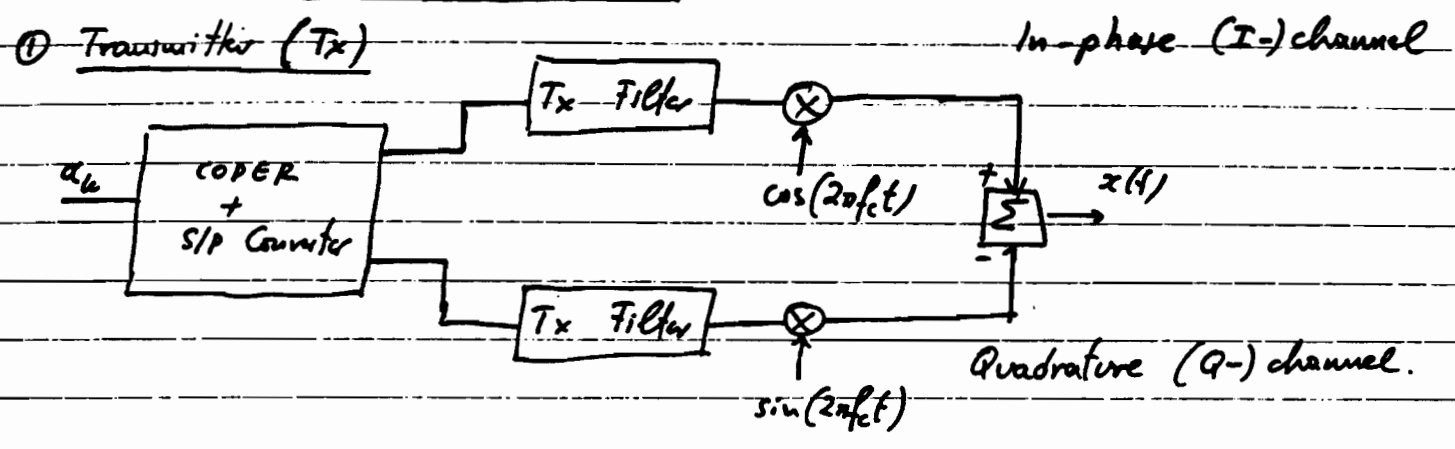
(Although theoretically $h_T(t)$, $h_R(t)$ can be complex, in computer simulation they are real.)

$$h_T(t) \leftrightarrow H_T(f) ; h_R(t) \leftrightarrow H_R(f)$$

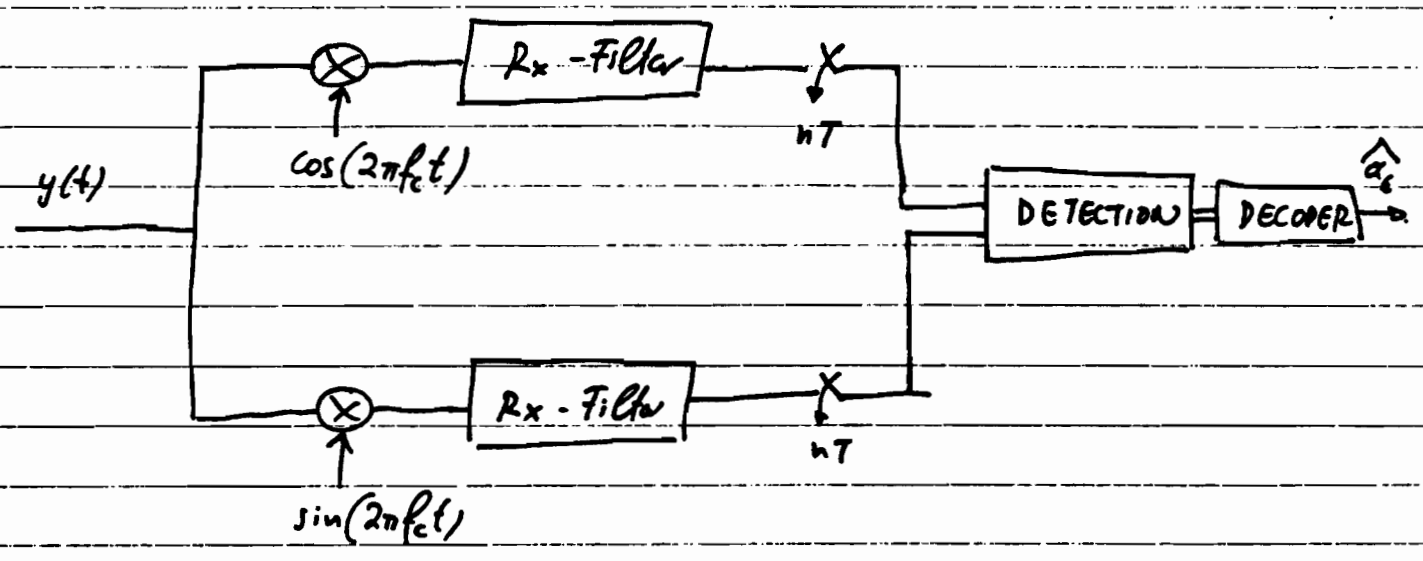
2) Computer Simulation Implementation (usually only multiplication)



IMPLEMENTATION.



② Receiver (Rx)



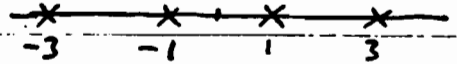
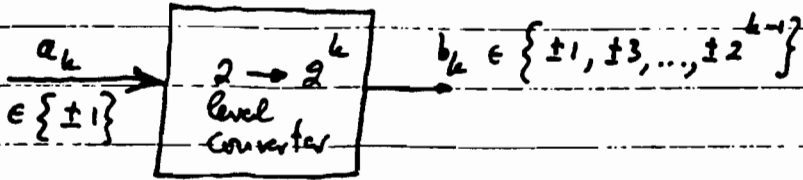
[QAM - Quadrature Amplitude Modulation]

MOST POPULAR "UNCODED" MODULATION SCHEMES

[Baseband] - unfiltered

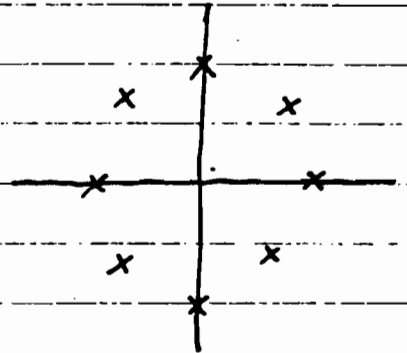
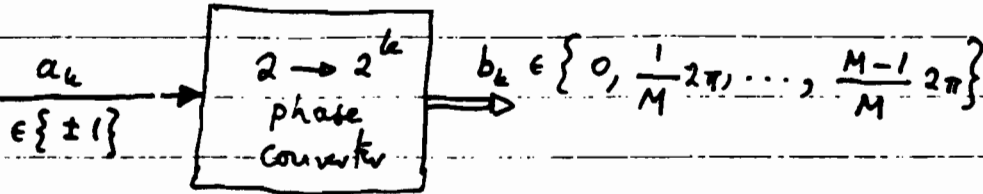
① M-ary PAM (Pulse Amplitude Modulation)

$M=4$

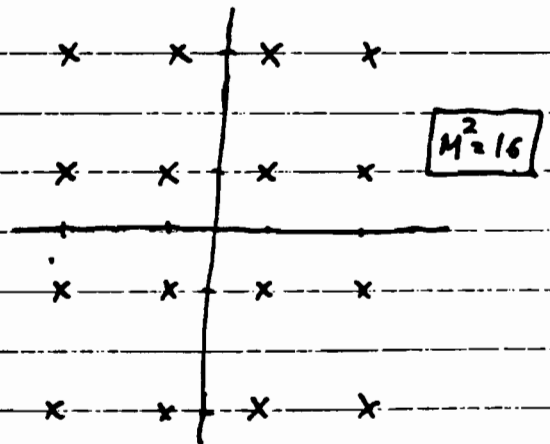
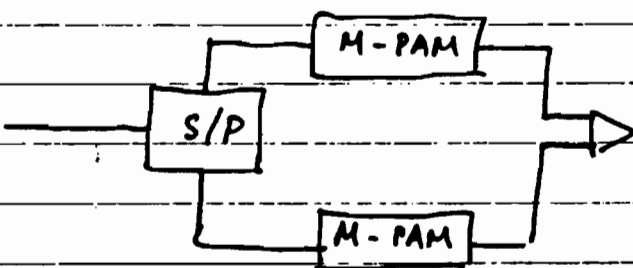


② M-ary PSK (Phase Shift Keying)

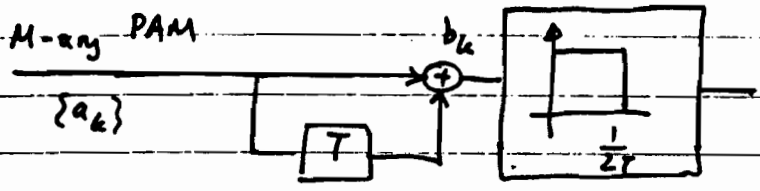
$M=8$



③ M²-QAM [belongs to the APK (Amplitude & Phase Shift Keying) family]

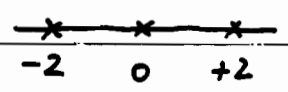


④ (2M-1) - PR (Partial Response)



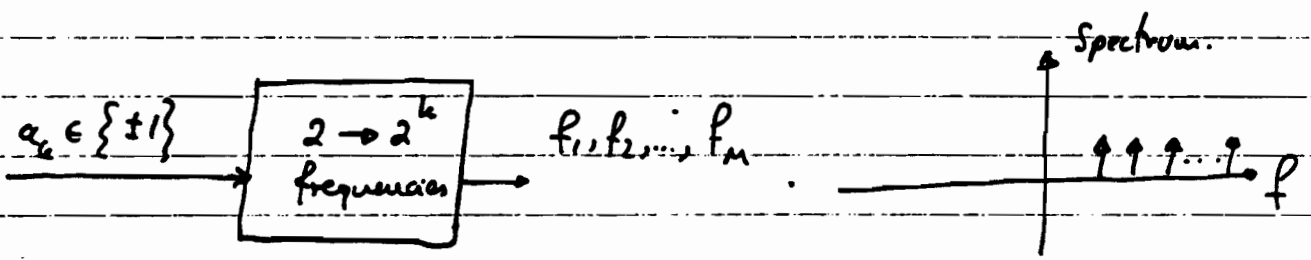
Example: $a_k \in \{\pm 1\} \rightarrow b_k \in \{+2, 0, -2\}$
 Prob: $\frac{1}{4}, \frac{1}{2}, \frac{1}{4}$

$2M-1 = 3$



- class -I PR
- Adding +1 or -1 \Rightarrow controlled amount of interference
- Comparison with M-PAM
 - Worst Performance
 - Better Spectrum
- QPR

⑤ M-FSK [Frequency Shift Keying]



- Continuous Phase FSK (e.g., MSK $f_2 - f_1 = \frac{1}{2T}$)
-

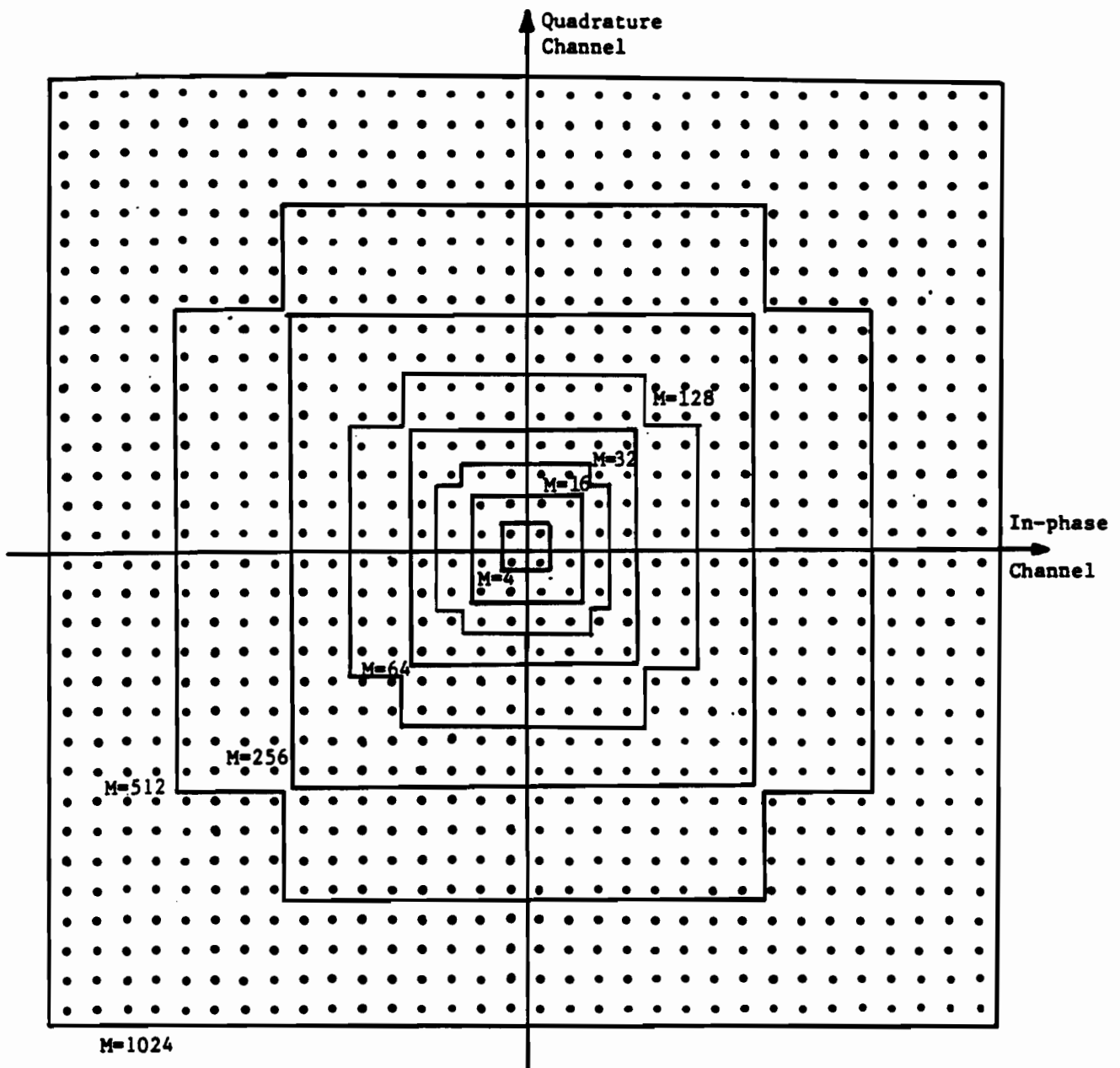


Figure 1.1: Constellations of M -ary QAM schemes with $M = 4, 16, 32, 64, \dots, 1024$.

I: in-phase carrier (channel) ; Q: quadrature-phase carrier (channel).

Ref. P. Mathiopoulos, "On bandwidth efficient QAM transmission systems," Ph.D. Thesis, University of Ottawa, Dept. of Electrical Engineering, 1989.

R: Information Rate in $\frac{\text{bits}}{\text{sec}}$

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$$P_M = k \cdot P_b$$

with $M = 2^k$

Also: $\gamma_M = \gamma_b \cdot k$

$$P_M = 10^{-5}$$

FSK with $\Delta f = \frac{1}{2T}$

Orthogonal signals coherent detection

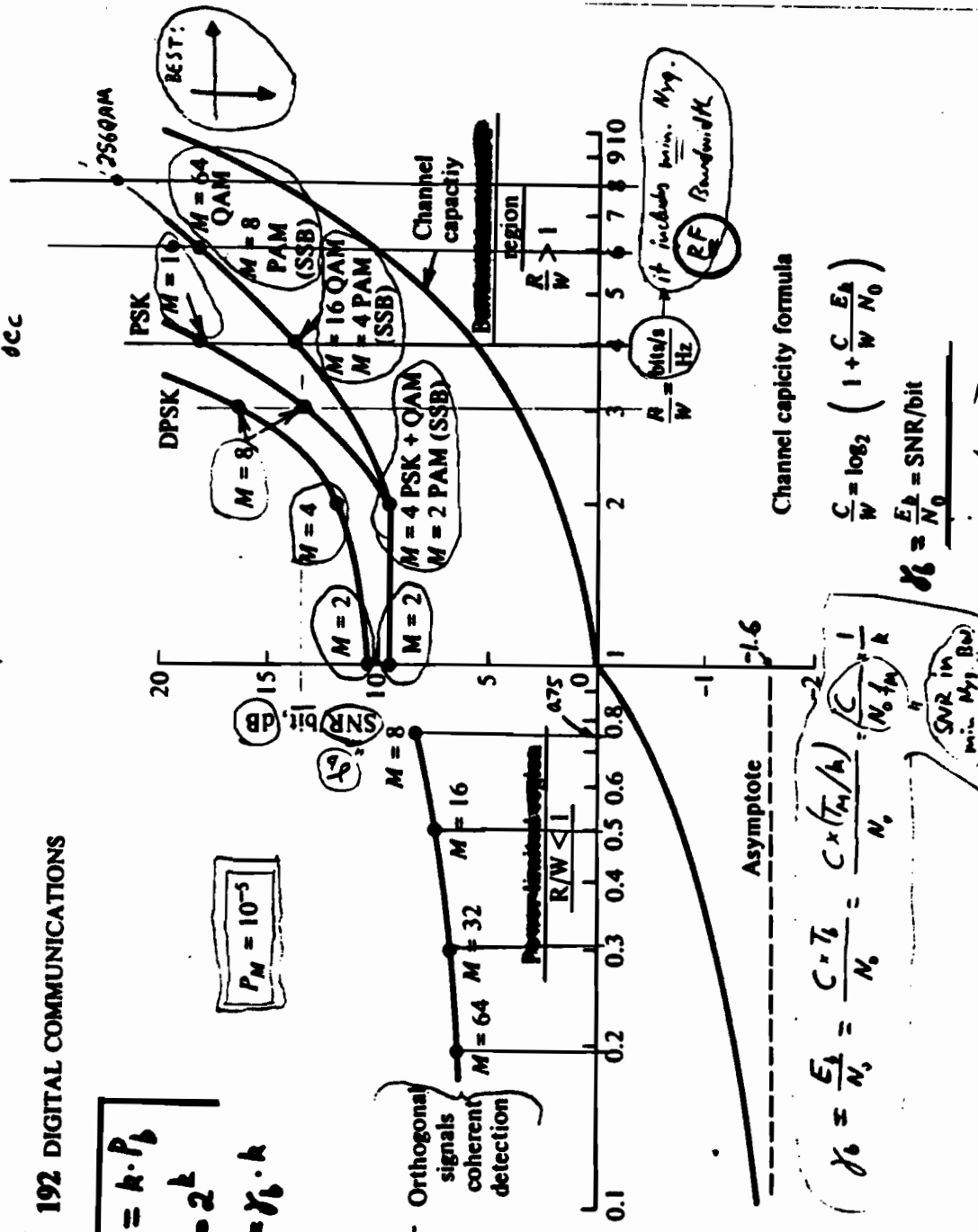


Figure 4.2.31 Comparison of several modulation methods at 10^{-5} symbol error probability.

Ref:

Proakis

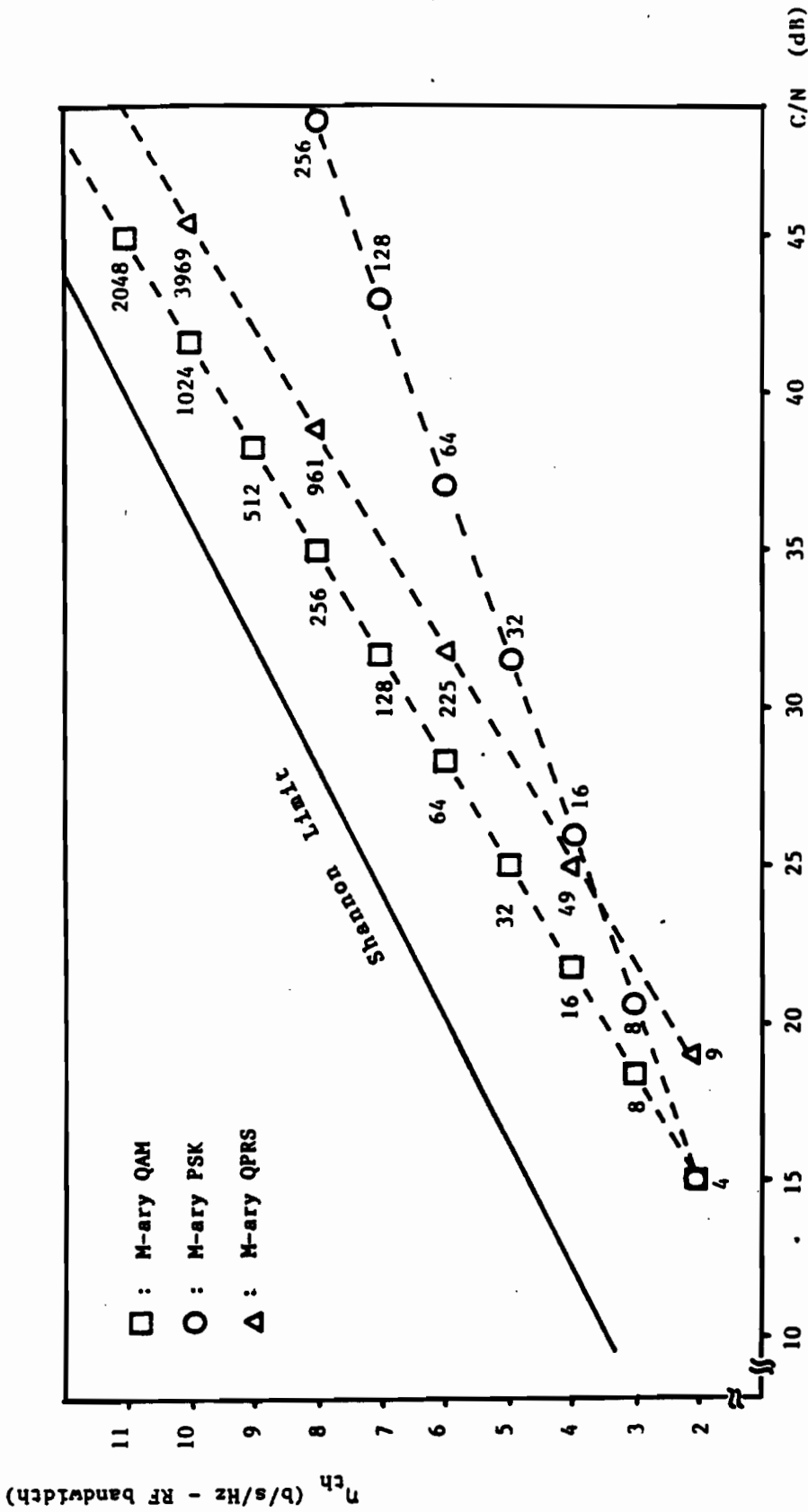
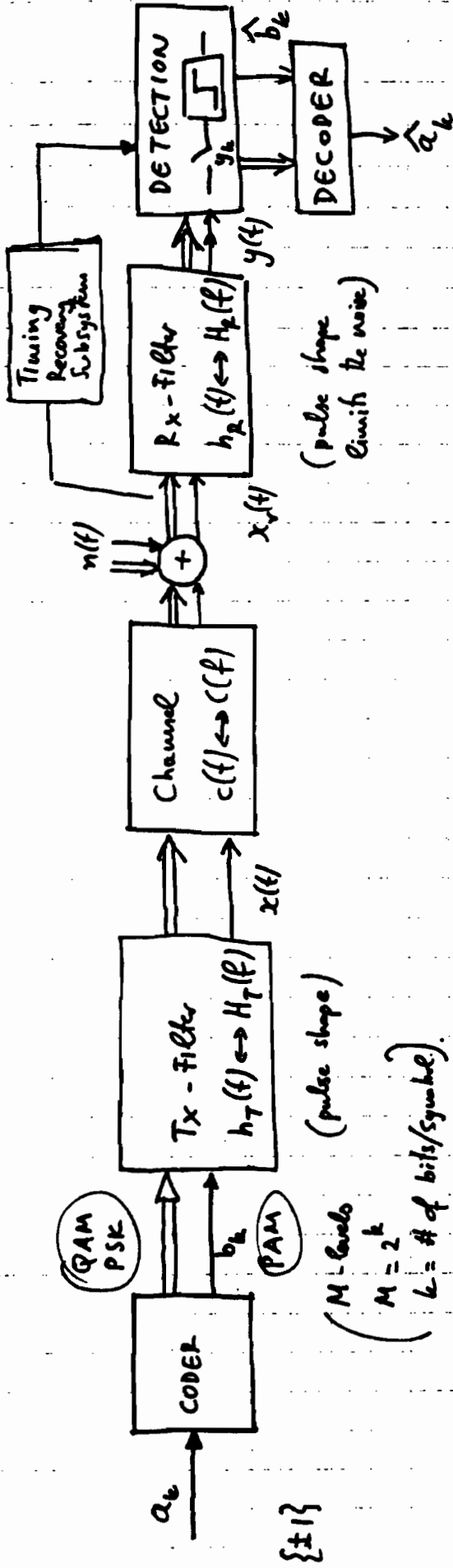


Figure 1.2: Theoretical spectral efficiency η_{th} in b/s/Hz of RF bandwidth and C/N requirement at $P_s = 10^{-6}$ of various spectrally efficient modulation systems. The average C/N is specified in the double-sided Nyquist bandwidth which equals the symbol rate. Since we refer to η_{th} ideal $\alpha = 0.0$ raised-cosine filters have been assumed.

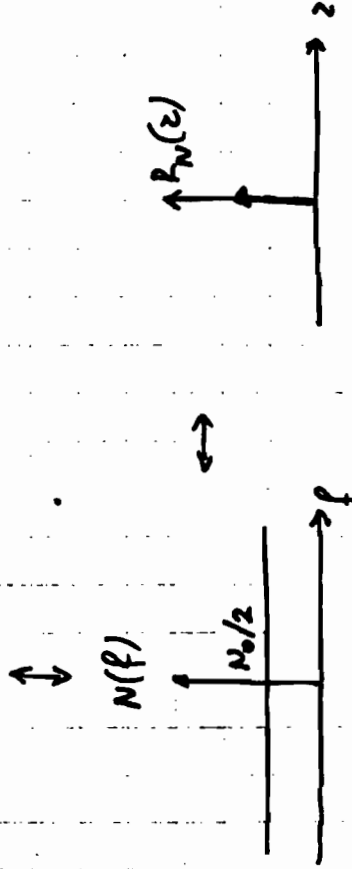
Ref: P. Mathiopoulos, "On bandwidth efficient QAM transmission systems," Ph.D. Thesis, University of Ottawa, Dept. of Electrical Engineering, 1987.

PERFORMANCE IN THE AWGN CHANNEL

• Baseband Systems.



$n(t)$: white Gaussian noise (additive)



white \Leftrightarrow uncorrelated
 Gaussian
 zero mean, σ^2 .

$$x(t) = \sum_k b_k h_T(t - kT) ; (b_k \text{ takes equiprobable values from an } M\text{-ary alphabet, e.g., } \pm d, \pm 3d, \dots, \pm(M-1)d)$$

$$x_r(t) = \sum_k b_k g(t - kT) + n(t) ; g(t) = h_T(t) \otimes c(t) = \int_{-\infty}^{\infty} h_T(z) c(t-z) dz$$

(Filtering: time dispersion reduction of bandwidth) $G(f) = H_T(f) C(f)$

$$y(t) = \sum_k b_k h(t - kT) + \tilde{n}(t) ; h(t) = h_T(t) \otimes c(t) \otimes h_R(t) ; H(f) = H_T(f) C(f) H_R(f)$$

$$\tilde{n}(t) = n(t) \otimes h_R(t)$$

Sample at $t = nT + t_0$ results:

$$y(nT + t_0) = \sum_k b_k h(nT - kT + t_0) + \tilde{n}(nT + t_0)$$

Assume ideal timing ($\Rightarrow t_0 = 0$)

$$y_n = y(nT) = \sum_k b_k h_{n-k} + \tilde{n}_n = h_0 \left(\overset{\substack{\uparrow \\ \text{useful} \\ \text{signal}}}{b_n} + \underbrace{\frac{1}{h_0} \sum_{k \neq n} b_k h_{n-k}}_{\text{Inter-Symbol-Interference (ISI)}} + \frac{\tilde{n}_n}{h_0} \right) \overset{\substack{\uparrow \\ \text{noise}}}{}$$

$h_0 =$ gain or attenuation through the system (easily compensated: $h_0 = 1$)

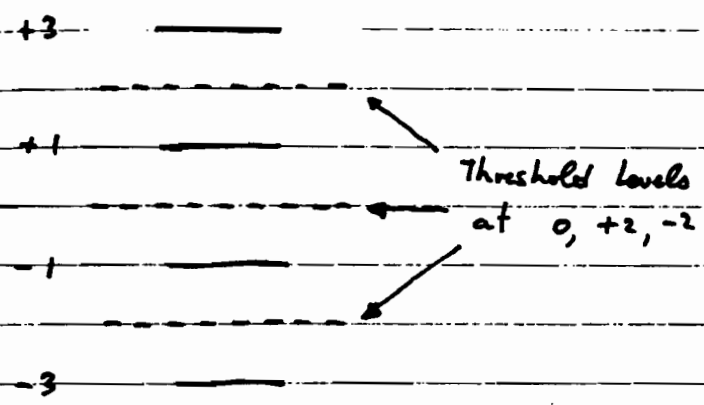
An error occurs when $\left| \frac{1}{h_0} \left[\sum_{k \neq n} b_k h_{n-k} + \tilde{n}_n \right] \right| > d$

$2d$ is the geometrical distance between the points

Example

Assume $d=1, b_0=1$

4-PAM



Problems: (ISI & Noise)
(or "useful" phenomena!?)

NYQUIST CRITERION FOR ZERO ISI

There are 3 Nyquist criteria.

The first deals with pulse shapes (essentially filters) which result in zero ISI at the sampling instant.

$$\text{If } \{h(t) = 0 \text{ for } t = nT \text{ (} n \neq 0)\} \iff \{h_n = 0 \text{ for } n \neq 0\}$$

$$\iff \{h_n = h(nT) = \delta_n\} \implies y_n = b_n + \tilde{n}_n \text{ (i.e., no ISI)}$$

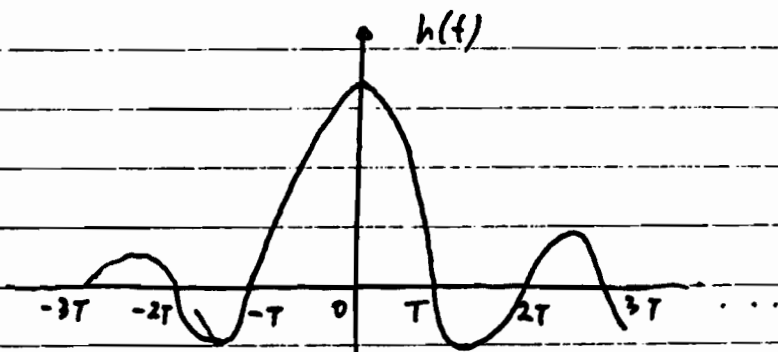
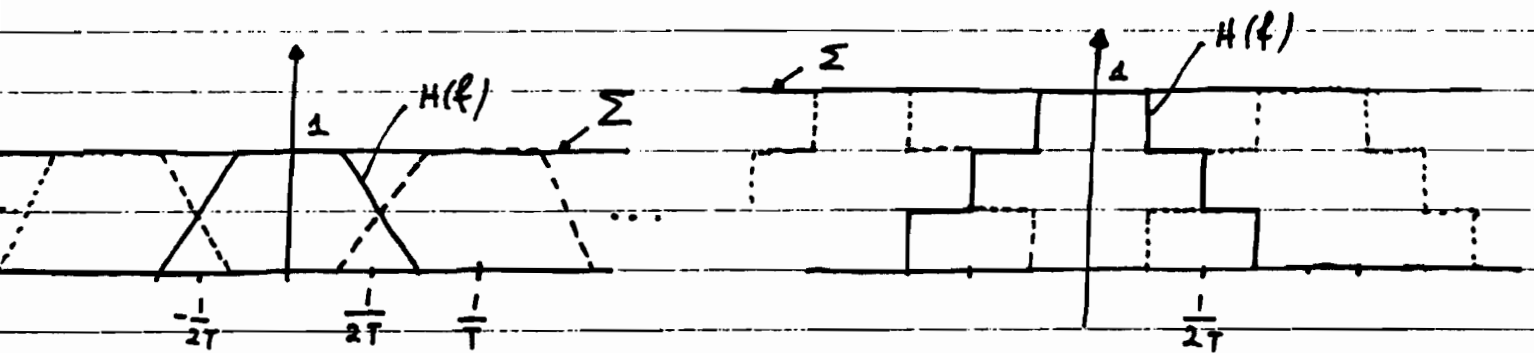
$$\circ\circ \text{ For no ISI } \implies \sum_k h(kT) \delta(t - kT) = \delta(t)$$

Nyquist
First
Criterion
(SOS)²

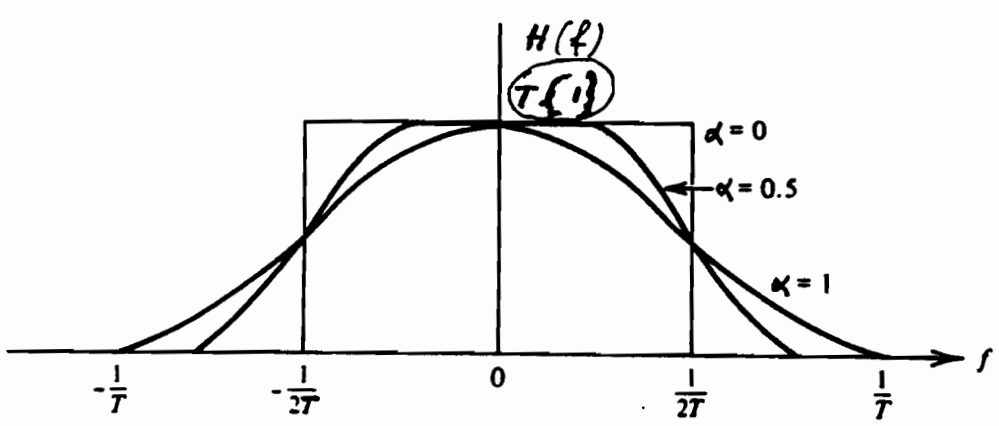
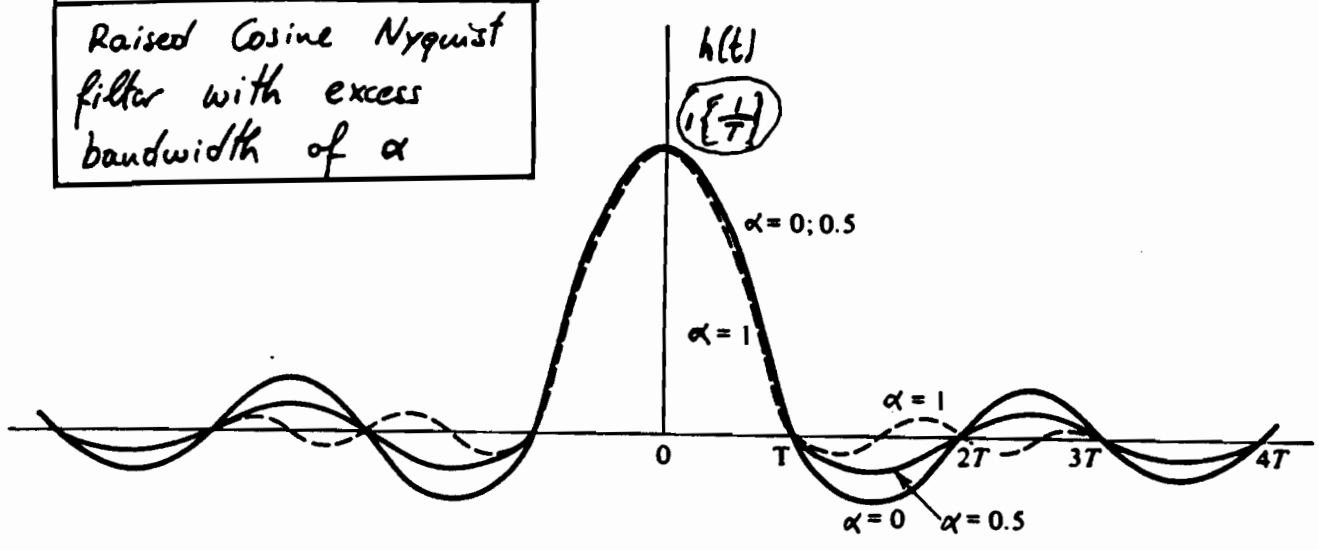
$h(t) \sum_k \delta(t - kT) = \delta(t)$ <p style="text-align: center;">Time Domain</p>	$\iff \frac{1}{T} \sum_m H\left(f - \frac{m}{T}\right) = 1$ <p style="text-align: center;">Frequency Domain</p>
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Comments

- ① Impulse response includes Tx-filter, channel and Rx-filter.
(For the Nyquist it does not matter how they are "split")
- ② In the time domain, we force the overall impulse response $h(t)$ to be zero at all sampling instances (except of course $t=0$).
- ③ In the frequency domain, this is equivalent of having an even symmetry around the Nyquist frequency $f_N = \frac{1}{2T}$



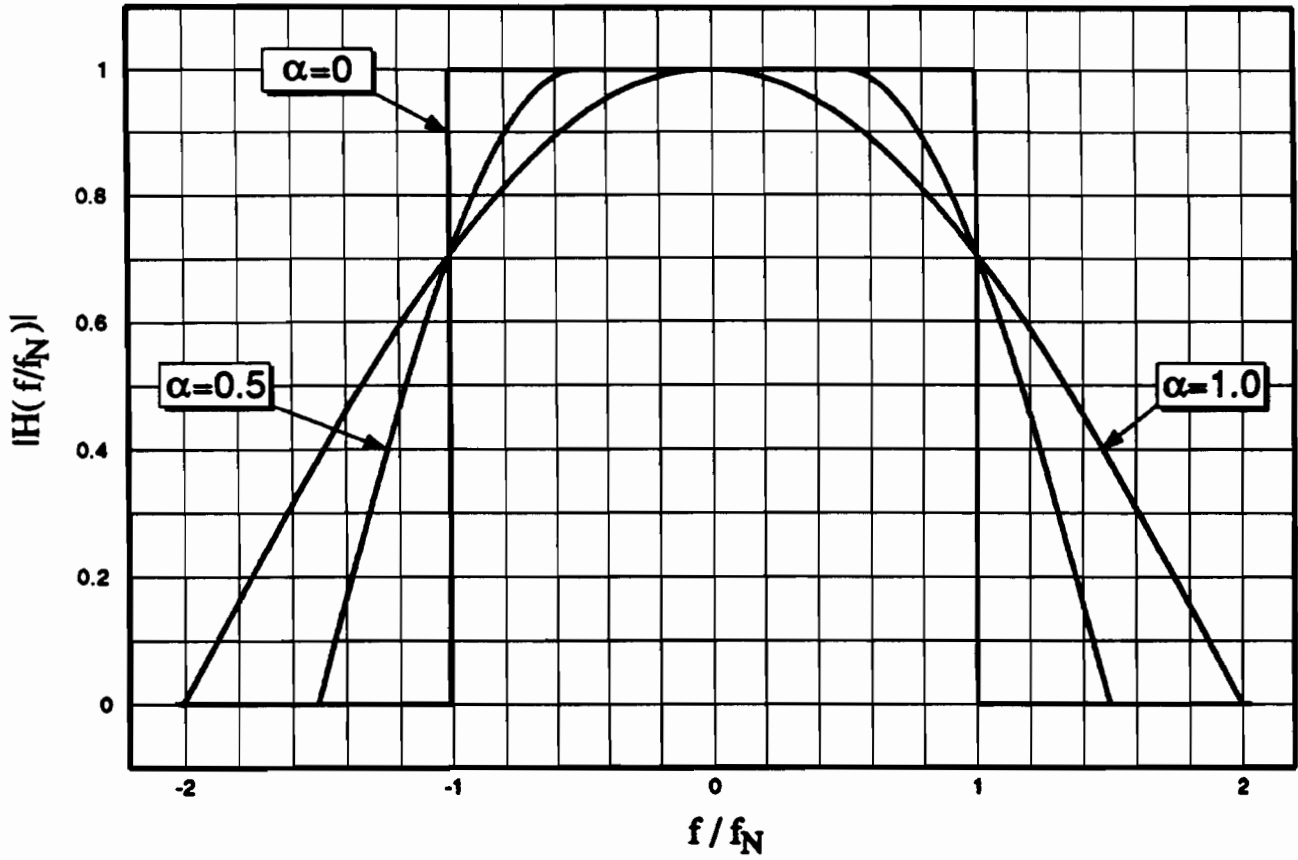
Raised Cosine Nyquist filter with excess bandwidth of α



$$H(f) = \begin{cases} 1 & \text{for } 0 \leq f < \frac{1-\alpha}{2T} \\ \cos^2 \left[\frac{\pi T}{2\alpha} \left(f - \frac{1-\alpha}{2T} \right) \right] & \text{for } \frac{1-\alpha}{2T} \leq f < \frac{1+\alpha}{2T} \\ 0 & \text{for } f \geq \frac{1+\alpha}{T} \end{cases}$$

$$h(t) = \frac{\sin\left(\frac{\pi t}{T}\right)}{\frac{\pi t}{T}} \frac{\cos\left(\frac{\alpha \pi t}{T}\right)}{1 - \frac{4\alpha^2 t^2}{T^2}}$$

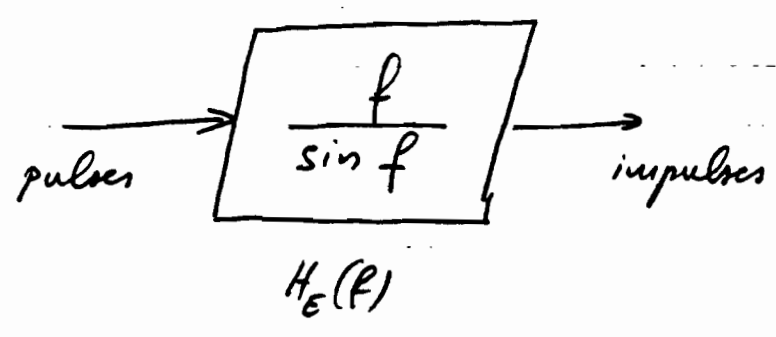
Root of raised-cosine transfer function



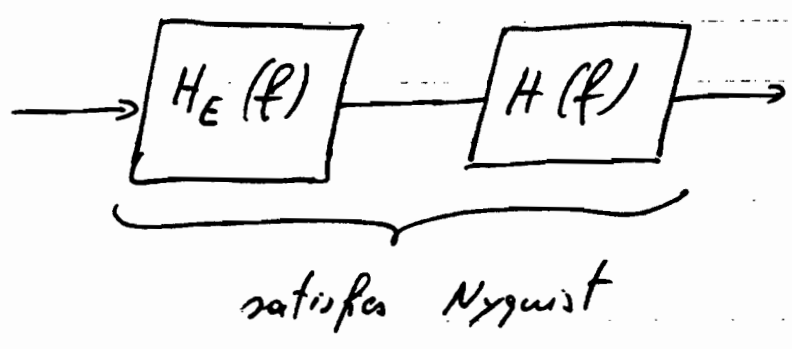
Important Note

The Nyquist filters (or equivalently impulse responses) have been derived assuming impulse transmission.

If we have pulse, which is the case for most practical applications, then we need a filter which transforms the pulses \rightarrow impulses

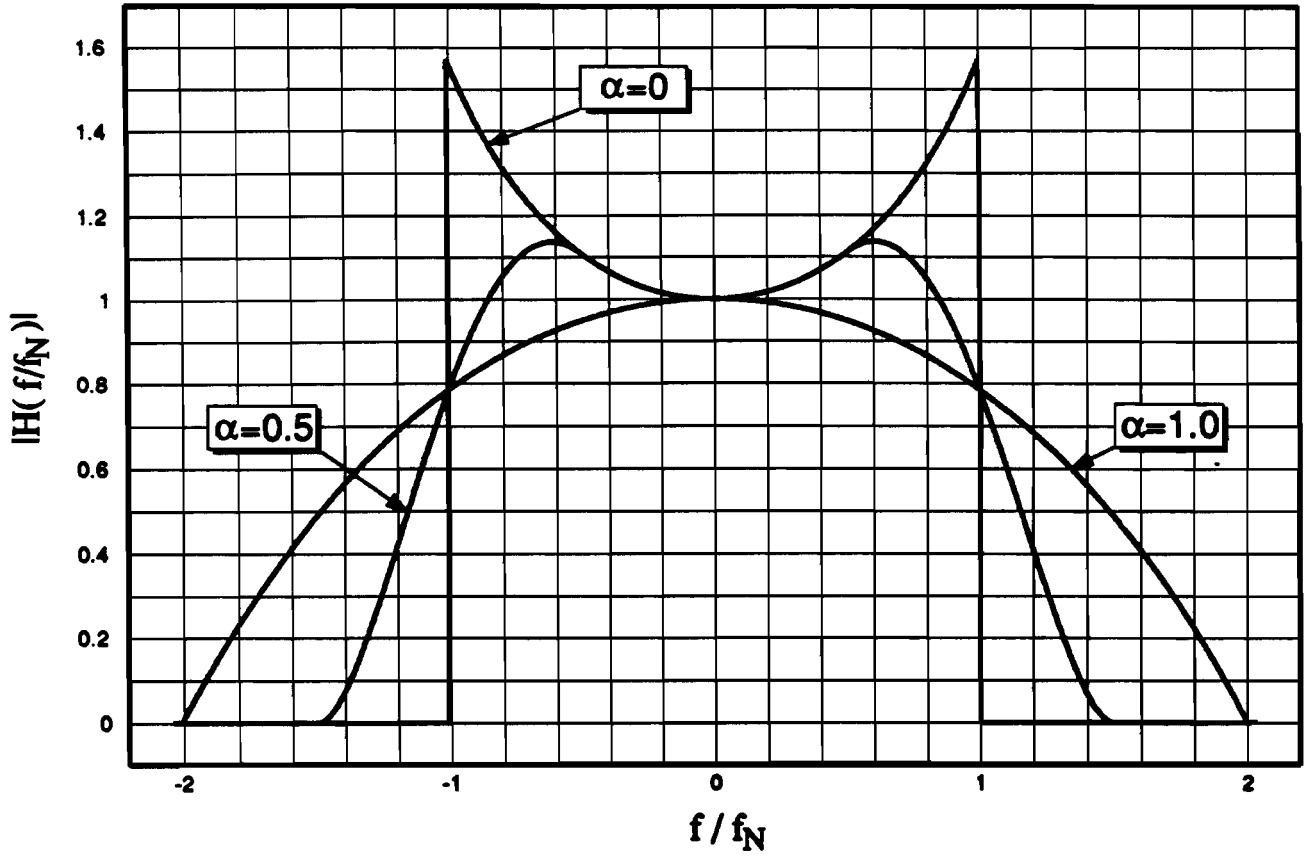


∴ For pulse transmission:

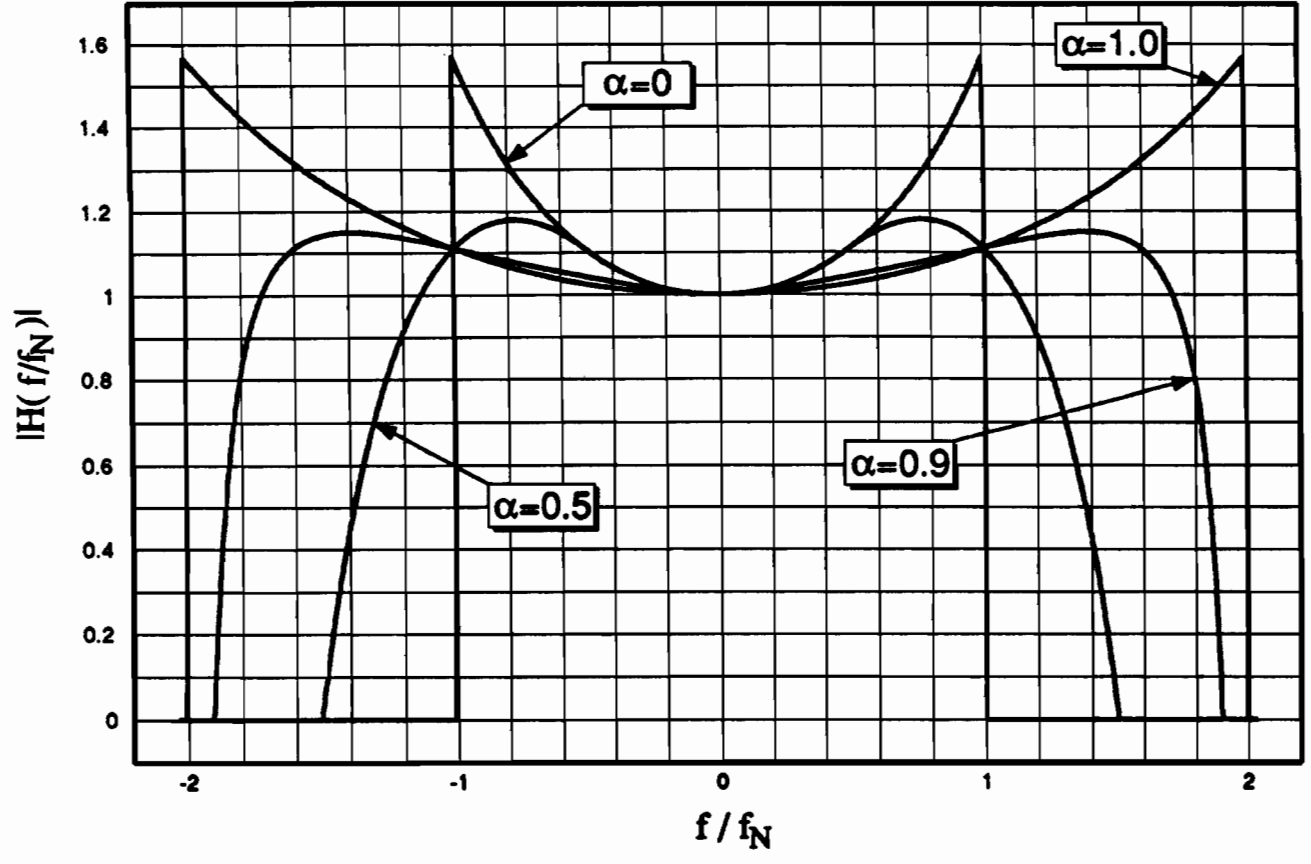


Also linear phase \Rightarrow constant group delay.

$x/\sin(x)$ equalized, raised-cosine transfer function



$x/\sin(x)$ equalized, root of raised-cosine transfer function



PARTIAL RESPONSE SIGNALS

- Bandlimited signals with controlled amount of intersymbol interference (ISI).
- For a raised cosine (α) pulse we have:

$$h(nT) = \begin{cases} 1 & \text{for } n=0 \\ 0 & \text{for } n \in \{\pm 1, \pm 2, \pm 3, \dots\} \end{cases}$$

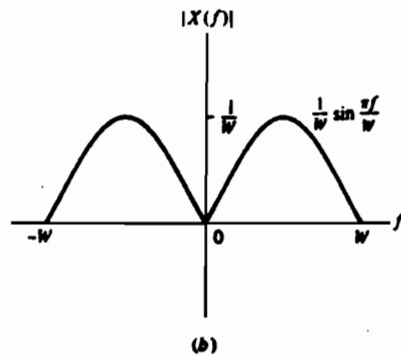
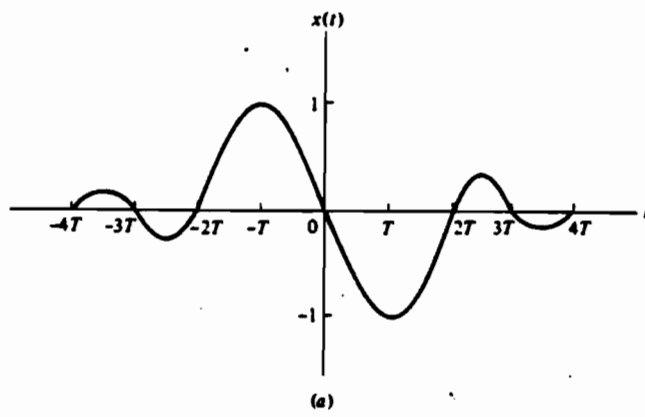
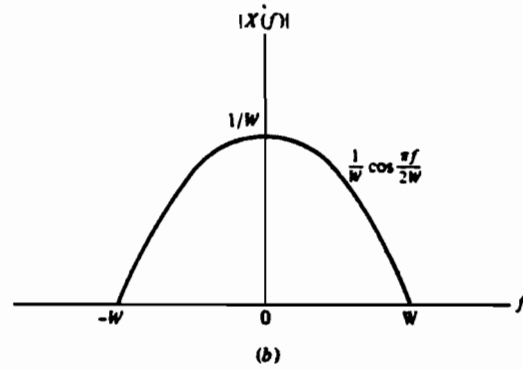
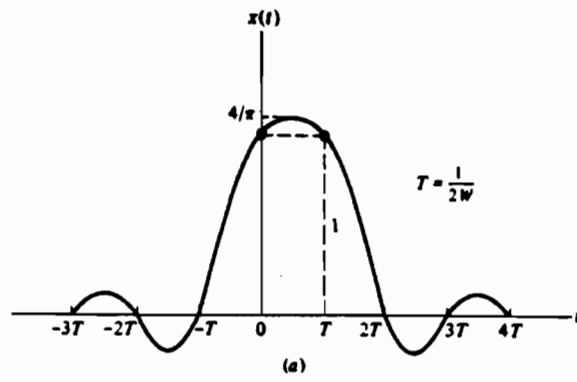
- Consider now that:

1)
$$h(nT) = \begin{cases} 1 & \text{for } n=0, 1 \\ 0 & \text{otherwise.} \end{cases}$$

[duobinary signalling]

2)
$$h(nT) = \begin{cases} 1 & \text{for } n=-1 \\ 1 & \text{for } n=1 \\ 0 & \text{otherwise} \end{cases}$$

[modified duobinary signalling]



- Problem with Error Propagation

Duobinary: $B_n = A_n + A_{n-1}$

$$A_n \in \{\pm 1\} \Rightarrow B_n \in \{\pm 2, 0\}$$

To decode A_n use $B_n - A_{n-1}$ (Note: Subtraction of ISI)

If A_{n-1} in error $\Rightarrow A_n$ in error too ...

\therefore Error propagation.

- Solution: Precoding

Essentially, it is a differential encoding in which we are transmitting the "difference in symbols" rather than the symbols themselves

Binary signaling with duobinary pulses

Data																
sequence D_n :	1	1	1	0	1	0	0	1	0	0	0	1	1	0	1	
Precoded																
sequence P_n :	0	1	0	1	1	0	0	0	1	1	1	1	0	1	1	0
Transmitted																
sequence A_n :	-1	1	-1	1	1	-1	-1	-1	1	1	1	1	-1	1	1	-1
Received																
sequence B_n :	0	0	0	2	0	-2	-2	0	2	2	2	0	0	2	0	
Decoded																
sequence D_n :	1	1	1	0	1	0	0	1	0	0	0	1	1	0	1	

$$D_n \in \{0, 1\}$$

$$P_n = D_n \oplus P_{n-1}$$

$$A_n = 2P_n - 1$$

$$B_n = A_n + A_{n-1} = 2(P_n + P_{n-1} - 1)$$

$$D_n = \frac{B_n}{2} \oplus 1 \quad [D_n = P_n \oplus P_{n-1}]$$

(0 denotes mod -2)

Four-level signal transmission with duobinary pulses

18E

Data													
sequence D_n :	0	0	1	3	1	2	0	3	3	2	0	1	0
Precoded													
sequence P_n :	0	0	0	1	2	3	3	1	2	1	1	3	2
Transmitted													
sequence A_n :	-3	-3	-3	-1	1	3	3	-1	1	-1	-1	3	1
Received													
sequence B_n :	-6	-6	-4	0	4	6	2	0	0	-2	2	4	2
Decoded													
sequence D_n :	0	0	1	3	1	2	0	3	3	2	0	1	0

$$D_n \in \{0, 1, 2, 3\}$$

Four-level signal transmission with modified duobinary pulses

Data													
sequence D_n :	0	1	3	1	2	0	3	3	2	0	1	0	0
Precoded													
sequence P_n :	0	0	0	1	3	2	1	2	0	1	2	1	3
Transmitted													
sequence A_n :	-3	-3	-3	-1	3	1	-1	1	-3	-1	1	-1	3
Received													
sequence B_n :	0	2	6	2	-4	0	-2	-2	4	0	2	0	0
Decoded													
sequence D_n :	0	1	3	1	2	0	3	3	2	0	1	0	0

In general: $D_n \in \{0, 1, 2, \dots, M-1\}$; $0 \equiv \text{mod } -M$

$$P_n = D_n \ominus P_{n-1} \quad \boxed{\text{Duobinary}} \quad \left| \sum_{k=0}^{\infty} \right| \quad P_n = D_n \oplus P_{n-2} \quad \boxed{\text{Modified Duobinary}}$$

$$A_n = 2P_n - (M-1)$$

$$D_n = \frac{B_n}{2} + (M-1) \quad \boxed{\text{Duobinary}} \quad \left| \sum_{k=0}^{\infty} \right| \quad D_n = \frac{B_n}{2} \quad \boxed{\text{Modified Duobinary}}$$

Generalization of a Technique for Binary Data Communication

E. R. KRETZMER, SENIOR MEMBER, IEEE

Abstract—A technique for binary data transmission is described, in which each binary symbol is chosen to be a prescribed superposition of n impulses of form $(\sin 2\pi Ft)/2\pi Ft$, spaced at intervals $1/2F$. Such superposition leads to more than two received levels with binary input, but in return permits realization of the Nyquist rate ($2F$ symbols/s in bandwidth F). Appropriate choice of the superposition coefficients provides a variety of spectral distributions to suit individual applications. The most interesting classes of system functions are defined, the associated coding procedures are described, and the performance characteristics are summarized.

INTRODUCTION

A useful generalization of the so-called biternary, duobinary, or polybinary [1]–[6] technique for binary data transmission is reported herewith. It is based on recognizing that one may choose an end-to-end system function $H(f)$ whose transform $h(t)$, the overall response per binary symbol, is a prescribed weighted linear superposition of n impulses of form $(\sin 2\pi Ft)/2\pi Ft$, spaced at intervals $1/2F$. Use of such a superposition rather than a single impulse leads to more than two received levels with binary input, but in return it renders practical the attainment of the Nyquist rate [7] ($2F$ symbols/s in bandwidth F), since $H(f)$ is zero for $f \geq F$ and is also continuous at $f = F$. Resulting constraints on level transitions in the received signal make possible a limited amount of error detection. Furthermore, appropriate choice of the superposition coefficients makes available a variety of spectral distributions to suit individual applications. This communication will define the most interesting classes of system functions (channels), describe the associated coding procedures, and summarize the performance characteristics.

DEFINITION

The transmission channel will be characterized by the superposition relationship mentioned: A received sample value c_n depends on n successive transmitted sample values $b_1 \dots b_n$ as follows:

$$c_n = k_1 b_n + k_2 b_{n-1} + \dots + k_n b_1, \quad (1)$$

where the k 's are integer weighting coefficients, with the smallest $|k|$ equal to unity.

The corresponding channel frequency function is

$$H(f) = \int_{-\infty}^{\infty} \left[k_1 \delta(t) + k_2 \delta\left(t - \frac{1}{2F}\right) + \dots + k_n \delta\left(t - \frac{n-1}{2F}\right) \right] e^{-j2\pi ft} dt \quad (2)$$

$[0 < |f| < F]$; the band limitation permits the use of the delta function instead of $(\sin 2\pi Ft)/2\pi Ft$.

With these two characterizations in mind we can now tabulate those channels which have been found of greatest interest (see Table I).

Each row of this table defines just one member of a class, namely the one having the smallest possible number n of k 's and hence of received levels. Thus, Class 1 is the polybinary class of systems [4] having all unity-weight coefficients, while Class 2 has only linearly tapered distributions of coefficient values. As n is increased within each class, the number of received levels also rises and $H(f)$ becomes increasingly "concentrated," yielding a more drastic spectral tapering. However, the full band F is still required to convey $2F$ bps, and performance margins are substantially reduced. To ensure the essential condition of continuity in $H(f)$ at $f = F$, certain superpositions must be ruled out, e.g., odd values of n in Class 1.

DECODING OR PRECODING

The received multilevel signal can be interpreted with a digital version of McColl's apparatus [8], in which the contributions of $n - 1$ preceding (transmitted) sample values—all binary in our case—are subtracted from the present received sample value. A binary decision is then made on the difference, which is either 0 or k_i in the absence of noise. A shift register stores the resulting binary digit as well as the $n - 1$ preceding ones (see Fig. 1). An error in any received sample tends to propagate until c_n reaches the top or bottom level.

Alternatively, a precoding operation can be undertaken at the transmitter, such that the n th binary digit a_n of the original sequence is the modulo-2 summation of the n th binary digit b_n actually transmitted and the $n - 1$ preceding digits transmitted, all weighted by their assigned coefficients, i.e.,

TABLE I

Class	k_1	k_2	k_3	k_4	k_5	$h(t)$	$H(f)$ [0 < f < F]	No. of Rec. Levels
Binary (ideal)	1						1	2
1 ($n = 2$)	1	1					$2 \cos \frac{\pi f}{2F}$	3
2 ($n = 3$)	1	2	1				$4 \cos^2 \frac{\pi f}{2F}$	5
3 ($n = 3$)	2	1	-1				$2 + \cos \frac{\pi f}{F} - \cos \frac{2\pi f}{F}$ $+ j \left[\sin \frac{\pi f}{F} - \sin \frac{2\pi f}{F} \right]$	5
4 ($n = 3$)	1	0	-1				$2 \sin \frac{\pi f}{F}$	3
5 ($n = 5$)	-1	0	2	0	-1		$4 \sin^2 \frac{\pi f}{F}$	5

* Entry in figure is area under curve.

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The author is with Bell Telephone Labs., Inc., Holmdel, N. J.

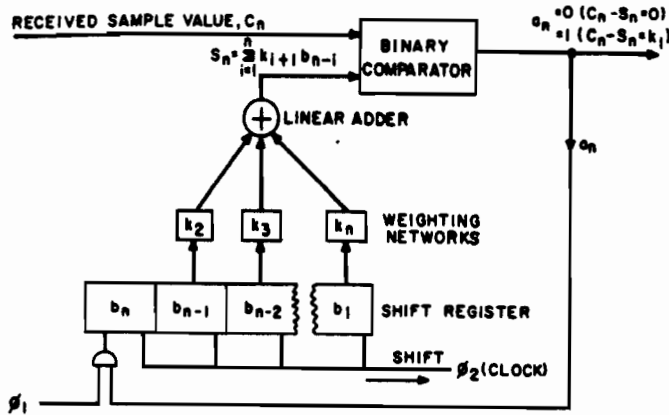
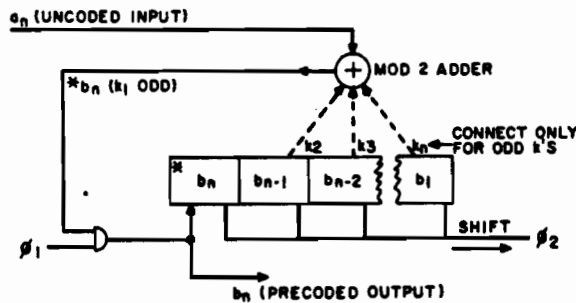


Fig. 1. Decoder.



*NOTE FOR k_1 EVEN, b_{n-1} REPLACES b_n , ETC

Fig. 2. Precoder.

$$a_n \equiv [k_1 b_n + k_2 b_{n-1} + \dots + k_n b_1] \text{ mod } 2. \quad (3)$$

Comparing this with (1), we find that

$$a_n \equiv c_n \text{ mod } 2, \quad (4)$$

i.e., an even-numbered value of the received sample implies a zero in the original sequence, an odd-numbered value a one. The precoding operation thus simplifies the interpretation of the received signal and eliminates error propagation; on the other hand, it limits the "decision distance" (difference in signal values to be distinguished by detector) to unity instead of k_1 —a matter of consequence only where k_1 exceeds unity, as in Class 3. The precoding implementation appears in Fig. 2.

PERFORMANCE

Performance is appraised by two criteria:

- 1) the speed tolerance, i.e., the percentage increase over the Nyquist rate which will just cause overlap between adjacent levels ("eye closure"); and
- 2) the required increase in signal-to-noise ratio over a reference binary system (operating at like bit rate) for a fixed error rate in the presence of white Gaussian noise (ignoring both error propagation and error detection capability).

Criterion 1 has been evaluated [9] by computer simulation using "worst possible" data sequences (see Table II). Criterion 2 is readily formulated in terms of the "decision distance" and the "noise bandwidth" of $[H(f)]^{1/2}$. The exponent 1/2 arises from a geometrically equal apportionment of $H(f)$ between transmitter

TABLE II

Class	k_1	k_2	k_3	k_4	Speed Tolerance k_1 (in percent)	S/N Increase Req. at $2F$ b/s
Ideal binary	1				0	0
1 ($n=2$)	1	1			43	2.1 dB
2 ($n=3$)	1	2	1		40	6.0 dB
3 ($n=3$)	2	1	-1		38	1.2 dB
4 ($n=3$)	1	0	-1		15	7.2 dB (precode)
5 ($n=3$)	-1	0	2	0	8	2.1 dB
($n=5$)						6.0 dB

and receiver. Consequently the ratio of the area under $H(f)$ relative to the area for the binary reference system gives a direct measure of the required increase in signal power. The area for each $H(f)$ given in Table I. The "decision distance" is unity in all cases except two for Class 3 when precoding is not used, since $k_1 = 2$, contributing a 6-dB improvement to the performance of this particular system [10]. The performance figures appear in Table II.

The speed tolerances listed in Table II may give some indication of other allowable departures from the prescribed transfer function. Further analysis is under way, including simulation of typical channel impairments superimposed on $H(f)$. For Classes 4 and 5, for example, this will test the practicability of a data baseband format achieving very efficient band utilization without dc transmission.

The virtues of the channels described are mainly that they afford relatively simple means of attaining $2F$ bps in bandwidth F with relatively gentle filter cutoffs, and that one can choose from a number of differently tapered spectral distributions which may be advantageous in specific applications. Instead of the basic $(2\pi Ft)/2\pi Ft$ impulse, any other functions satisfying Nyquist's first criterion [7] may be used as channel response components. An infinite number of such functions exist, their bandwidth ranging from F to $2F$.

ACKNOWLEDGMENT

The author acknowledges the contributions of F. K. Becker and H. O. Burton in clarifying some of the preceding concepts.

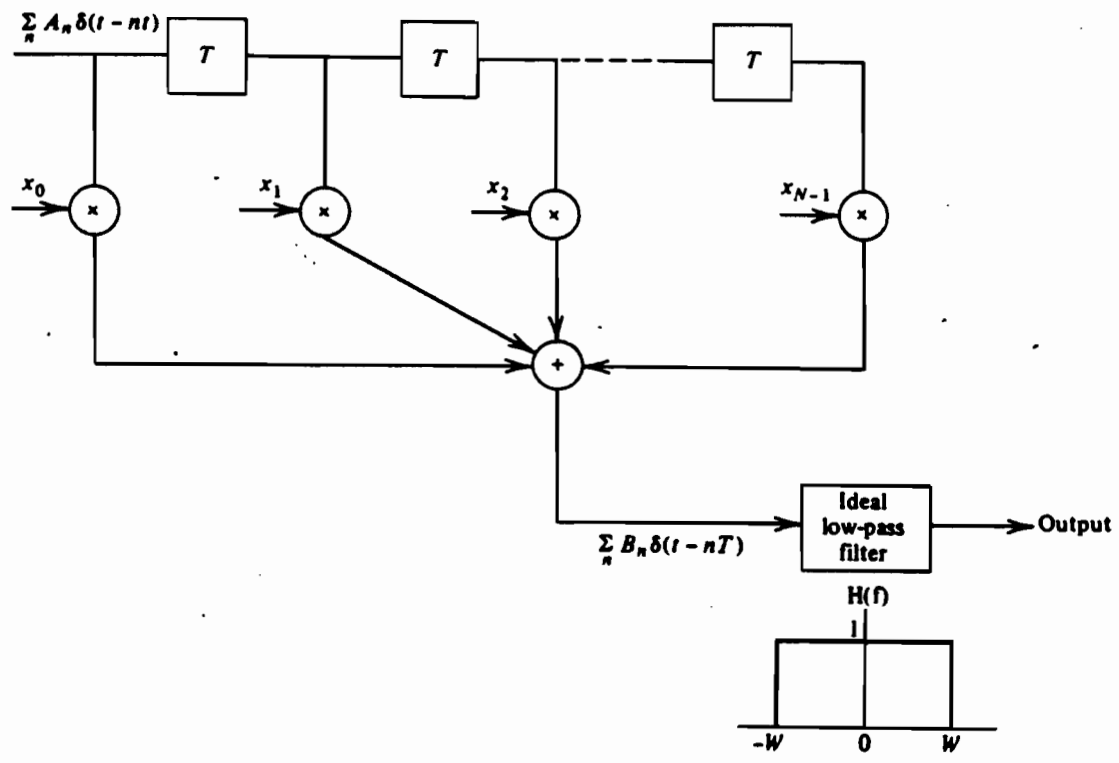
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- [2] A. P. Brogie, "A new transmission method for pulse-code modulation communication systems," *IRE Trans. on Communications Systems*, CS-8, pp. 155-160, September 1960.
- [3] O. E. Ringelhaan, "System for transmission of binary information at twice the normal rate," U. S. Patent 3 162 724, December 1964.
- [4] A. Lender, "The duobinary technique for high-speed data transmission," *IEEE Trans. on Communication and Electronics*, vol. 82, pp. 214-219, May 1963. See also A. Lender, "Correlative digital communication techniques," *IEEE Trans. on Communication Technology*, vol. COM-12, pp. 128-135, December 1964.
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- [8] L. A. McColl, "Signalling method and apparatus," U. S. Patent 2 056 200, October 1936.
- [9] W. R. Bennett and J. R. Davay, *Data Transmission*. New York: McGraw-Hill, 1965, ch. 7. See also, S. Habib, internal communication.
- [10] R. L. Townsend, "A class of partial response systems," internal communication.

TABLE I

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4 ($n=3$)	1	0	-1				$2 \sin \frac{f}{F}$	3
5 ($n=5$)	-1	0	2	0	-1		$4 \sin^2 \frac{f}{F}$	5

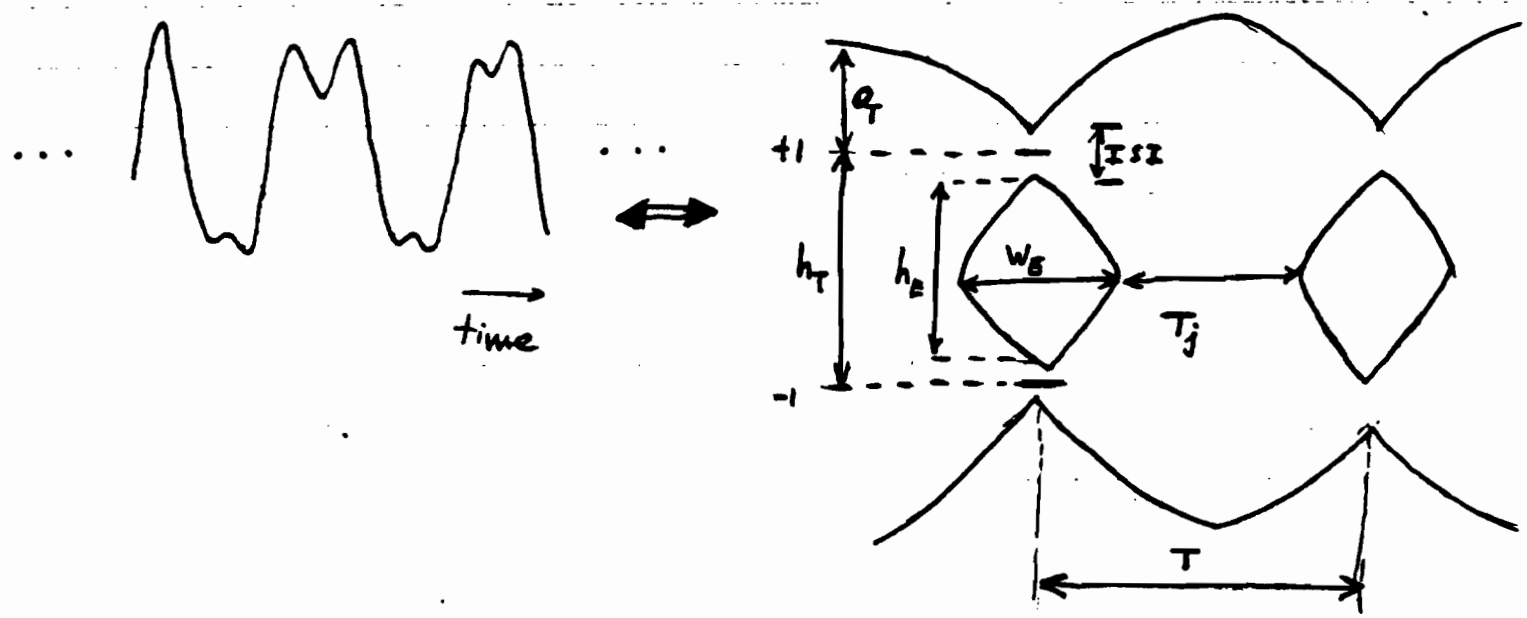
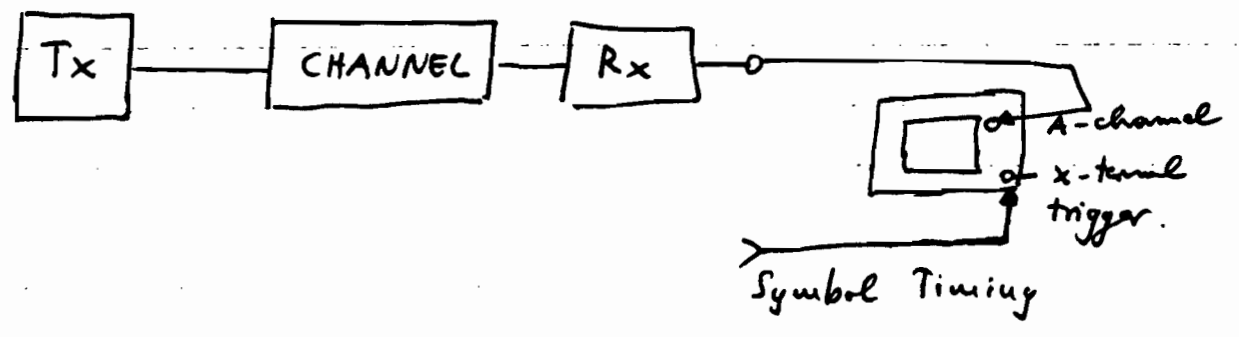
* Entry in figure is area under curve.



Implementation: Correlative Encoder.

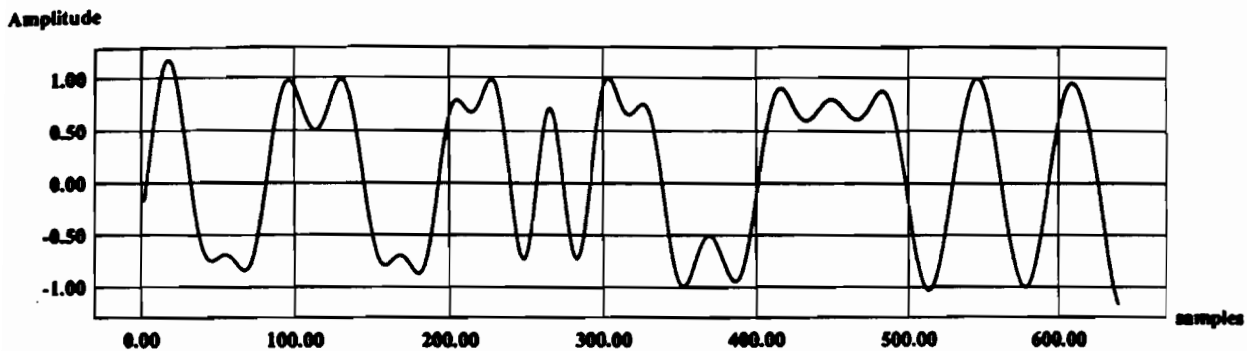
EYE DIAGRAMS.

- Imperfections → degradation
 - (Filtering, nonlinearities, carrier and symbol recovery)
- Graphical illustration of this degradation: Eye-diagram
 - Quick check of the performance of a modem in the field
 - Useful in the analysis and simulation.

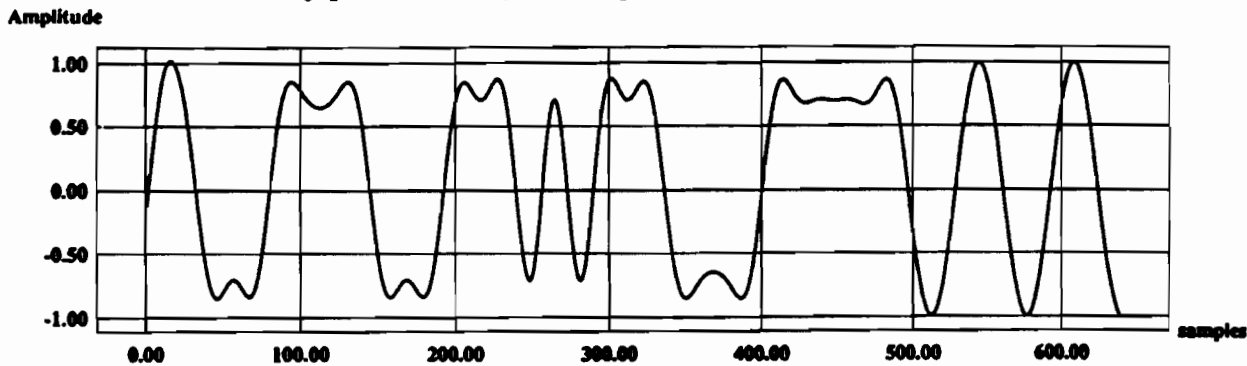


h_E = eye height ; w_E = eye width
 O_T = overshoot
 T_j = jitter

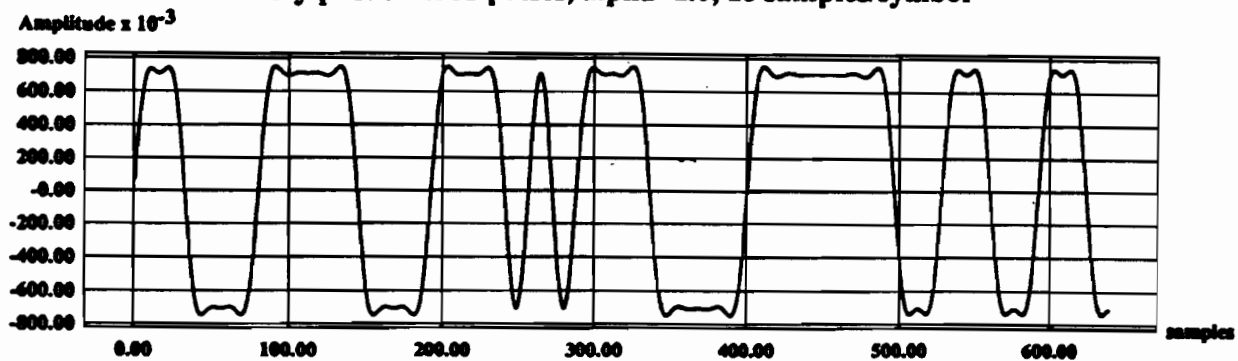
Nyquist filtered pulses, $\alpha=0.2$, 16 samples/symbol



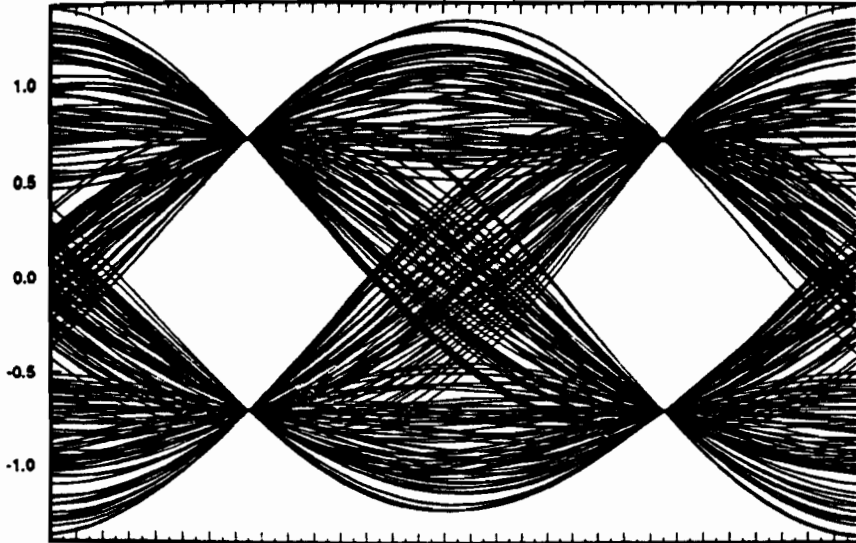
Nyquist filtered pulses, $\alpha=0.5$, 16 samples/symbol



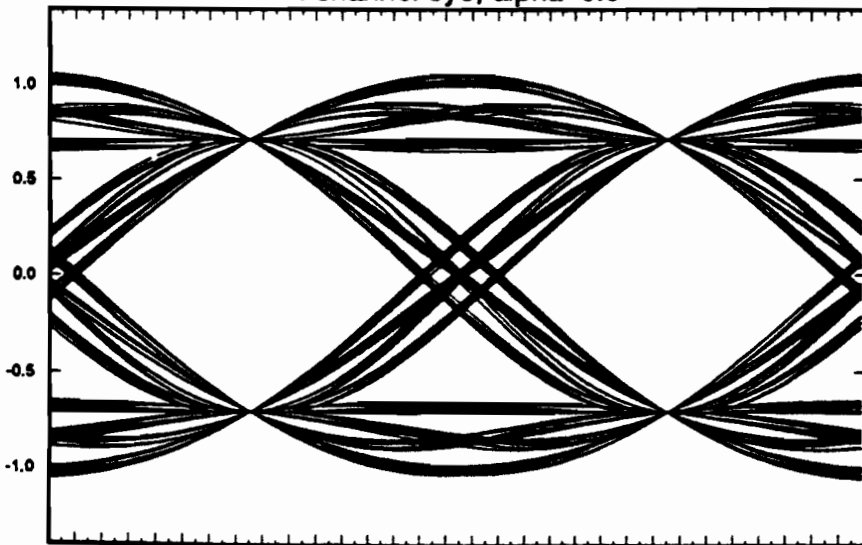
Nyquist filtered pulses, $\alpha=1.0$, 16 samples/symbol



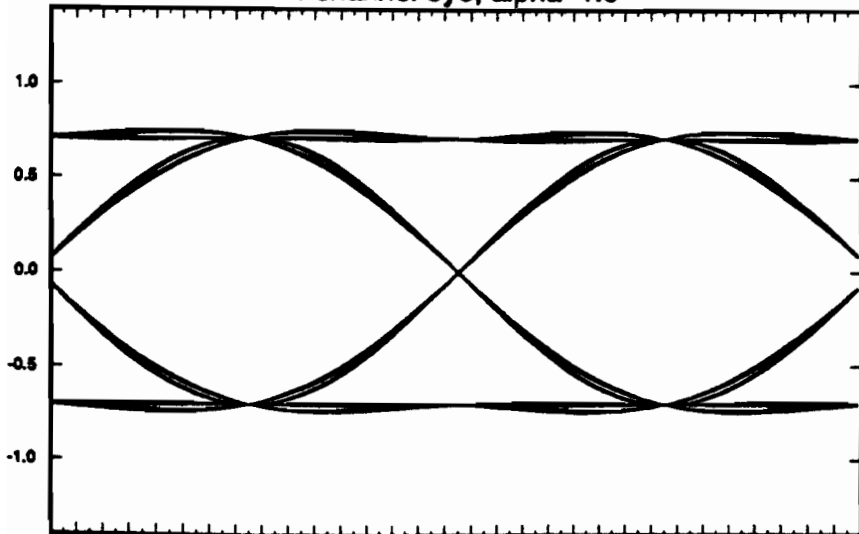
I-channel eye, alpha=0.2



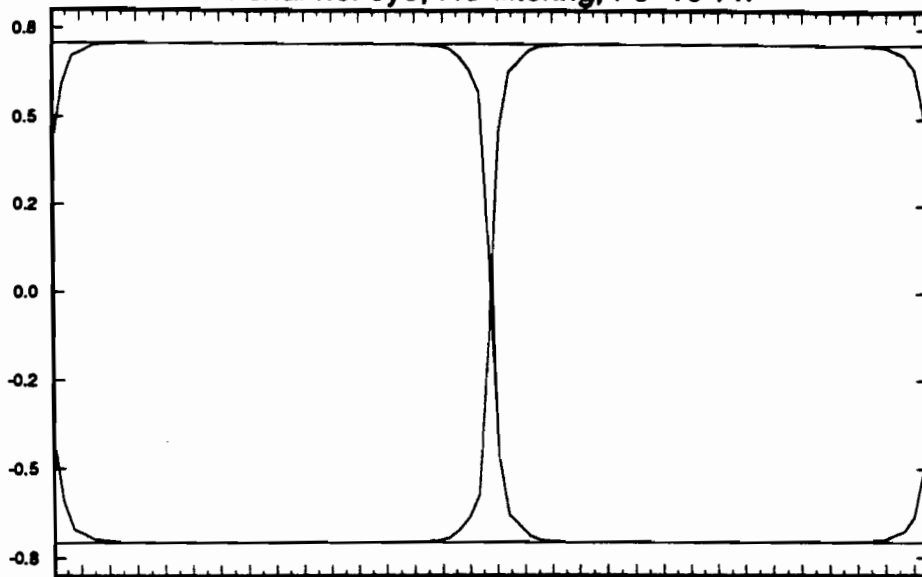
I-channel eye, alpha=0.5



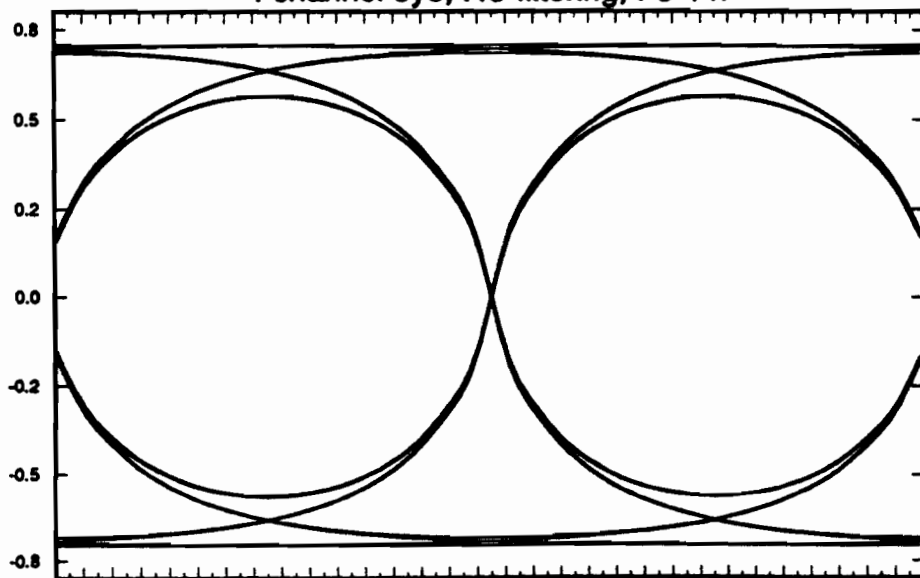
I-channel eye, alpha=1.0



I-channel eye, RC filtering, $F_c=10 \cdot F_n$



I-channel eye, RC filtering, $F_c=F_n$



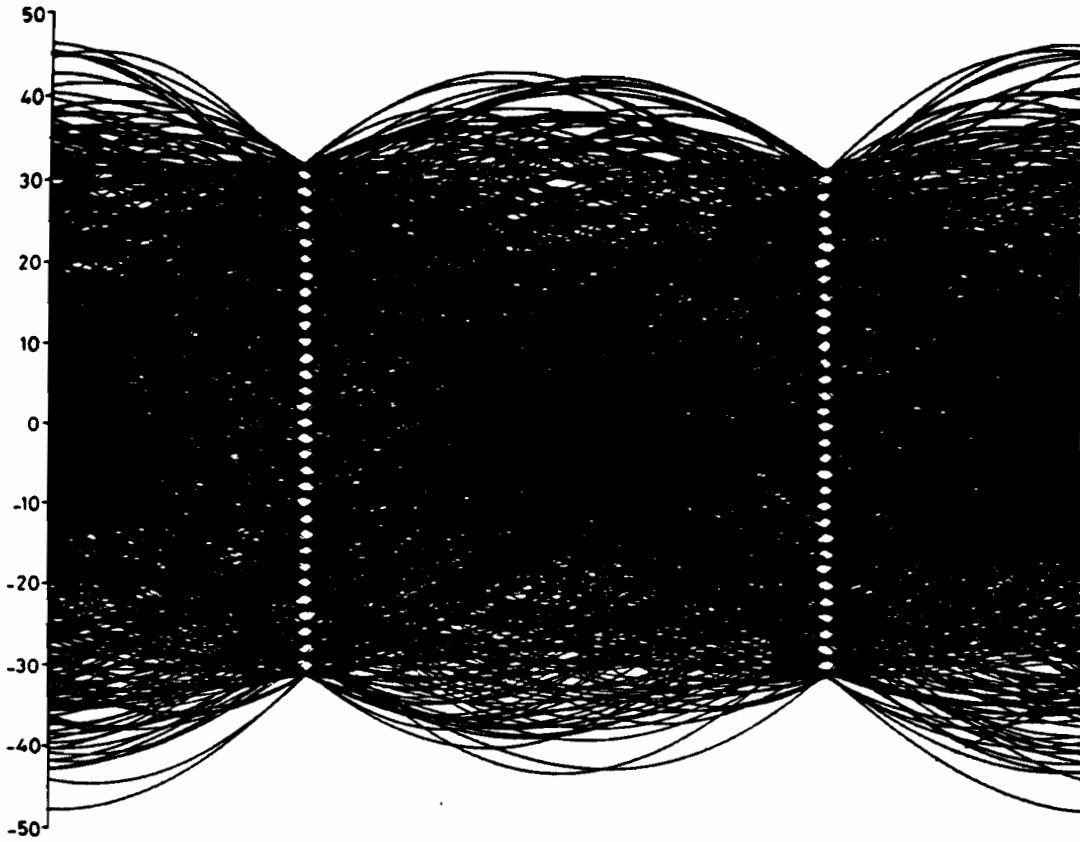


Fig. 3 Eye diagram of the in-phase channel of the 1024-QAM using $\alpha = 0.2$ raised cosine filters and a channel with parabolic amplitude distortion

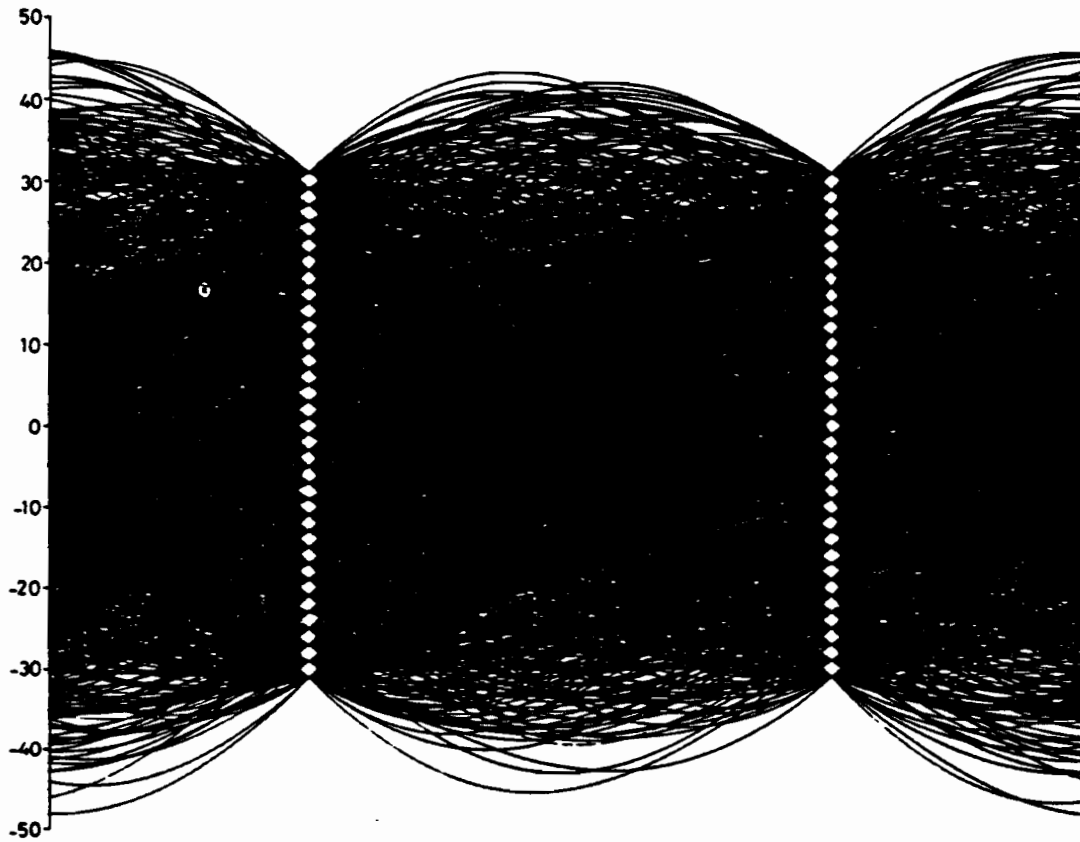
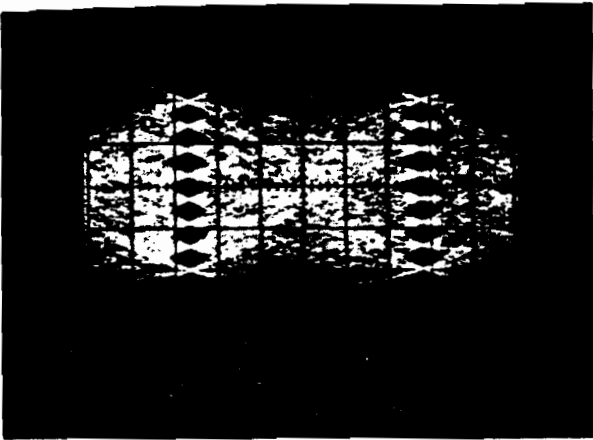


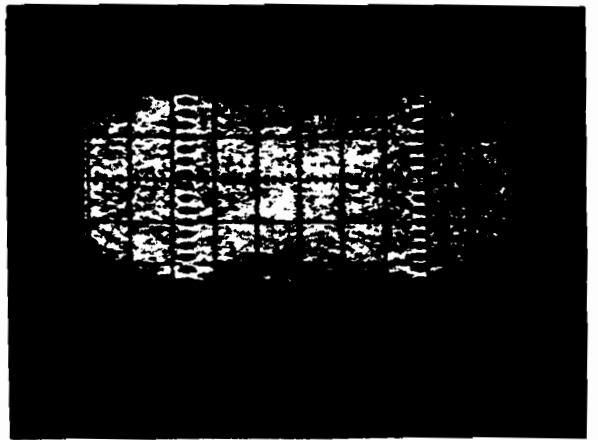
Fig. 2 Eye diagram of the in-phase channel of the 1024-QAM system using raised cosine filters with a roll-off factor of $\alpha = 0.2$

Ref. P. Mathiopoulos, H. Ohnishi and K. Fchar, "Study of 1024-QAM system performance in the presence of filtering imperfections," IEE Proc., Pt I, Apr 1990.

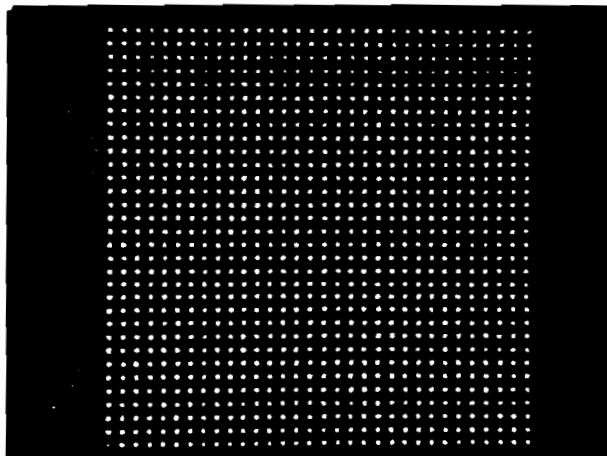
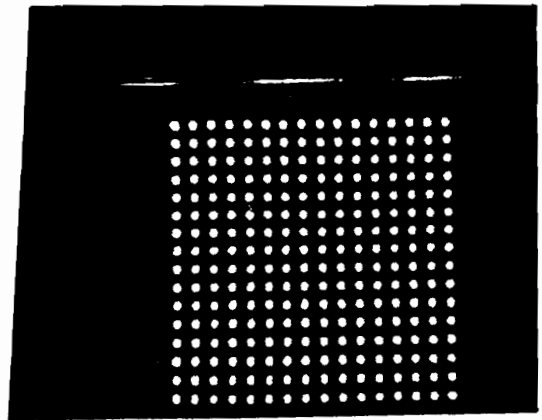
EXPERIMENTAL EYES & STATE-SPACE DIAGRAMS



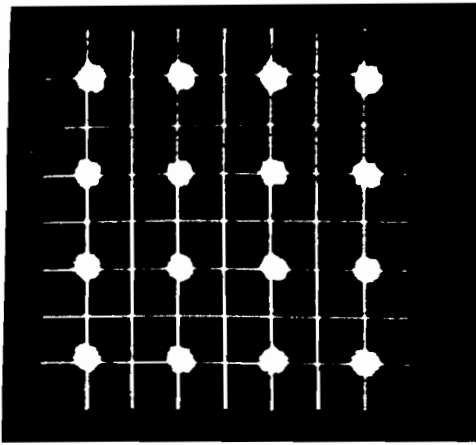
64-QAM



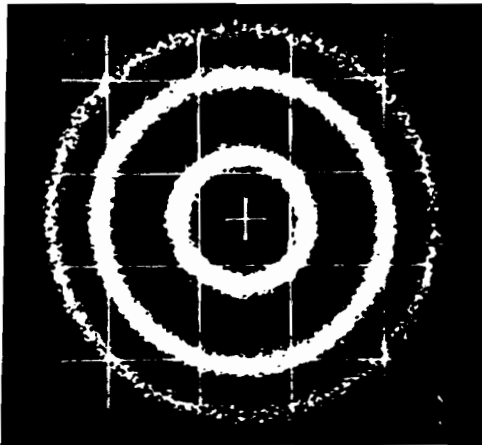
256-QAM



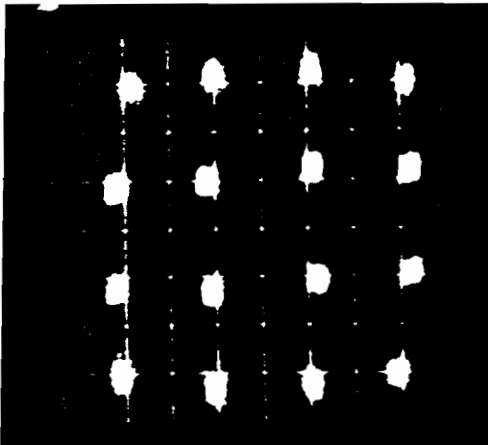
1024-QAM



(a)



(b)

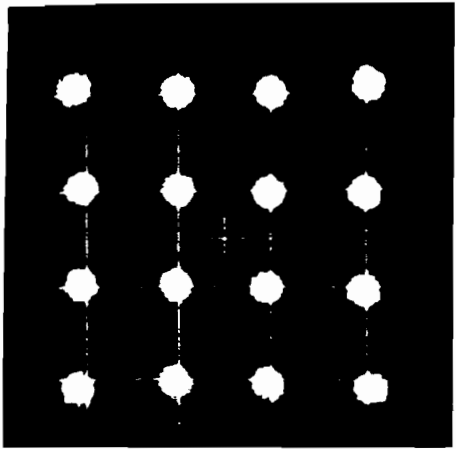


(c)

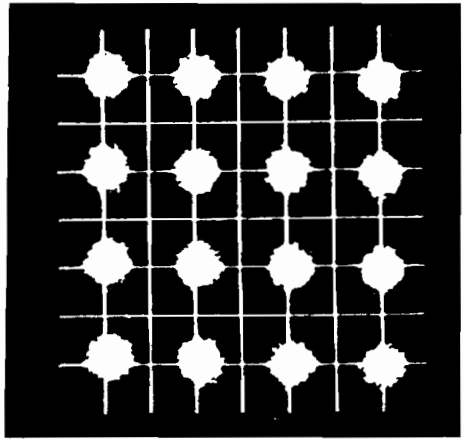
Figure 7.22 16-QAM constellation with geometric defects due to specific faults. (a) 16-QAM radio: Normal constellation plus eye-closure data. (b) Receiver out of lock. (c) 16-QAM radio: 3-dB transmitter power overdrive (AM-AM and AM-PM).

(continued)

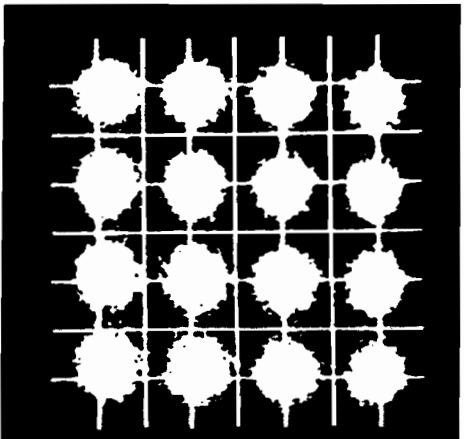
Ref. Fehar/Eng. of HP, Telecommunications, Measurement Analysis and Instrumentation, Prentice Hall, 1987.



(d)

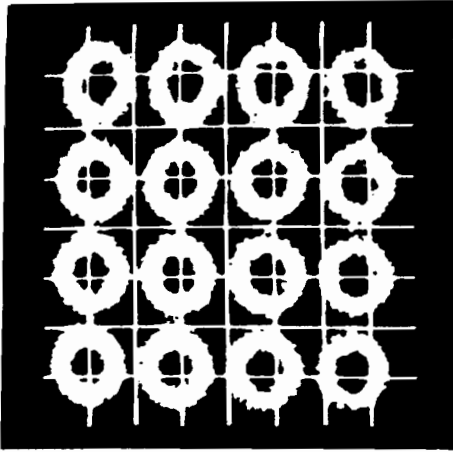


(e)

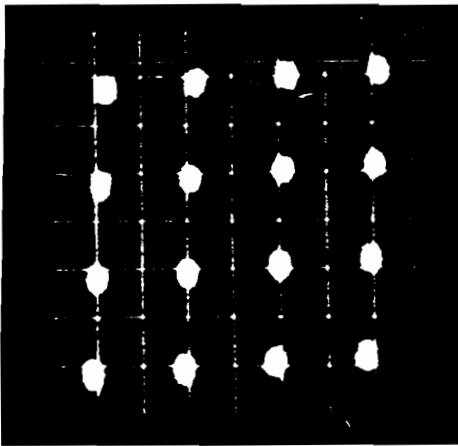


(f)

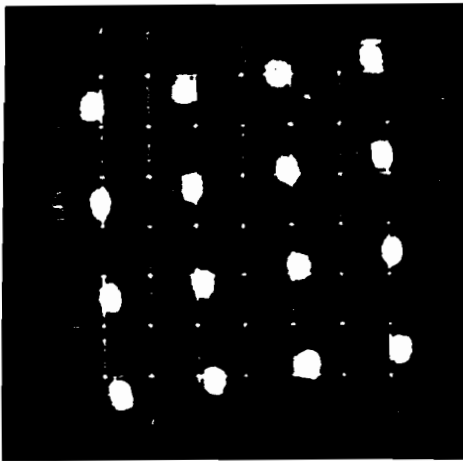
Figure 7.22 (continued) (d) Expansion due to underdrive of TWT. (e) 16-QAM radio: C/N ratio 20 dB (setup using 3708A). (f) 16-QAM radio: C/N ratio 15 dB (setup using 3708A).
(continued)



(g)



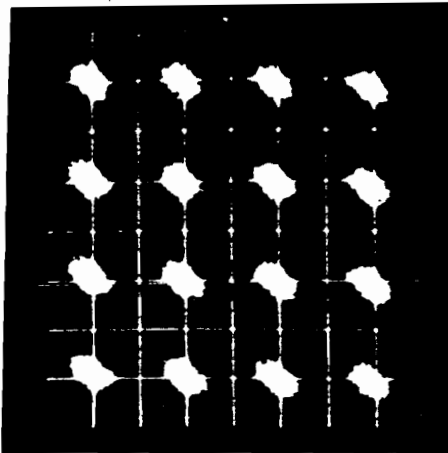
(h)



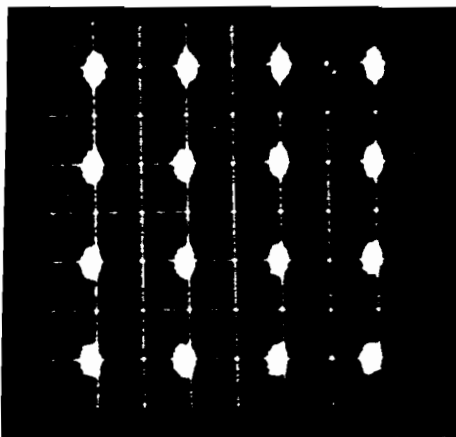
(i)

Figure 7.22 (continued) (g) 16-QAM radio: interferer tone present. (h) 16-QAM radio: recovered carrier / vs. Q phase nonorthogonal. (i) 16-QAM radio: carrier recovery loop lock misadjusted.

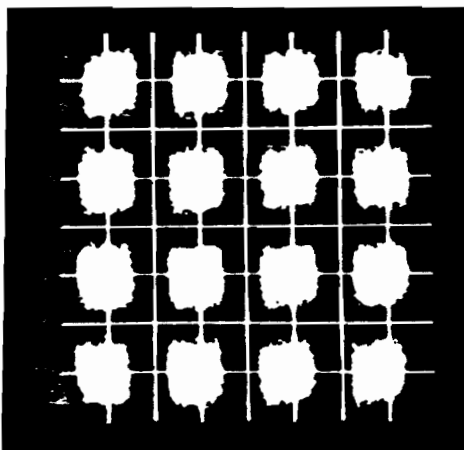
(continued)



(j)



(k)



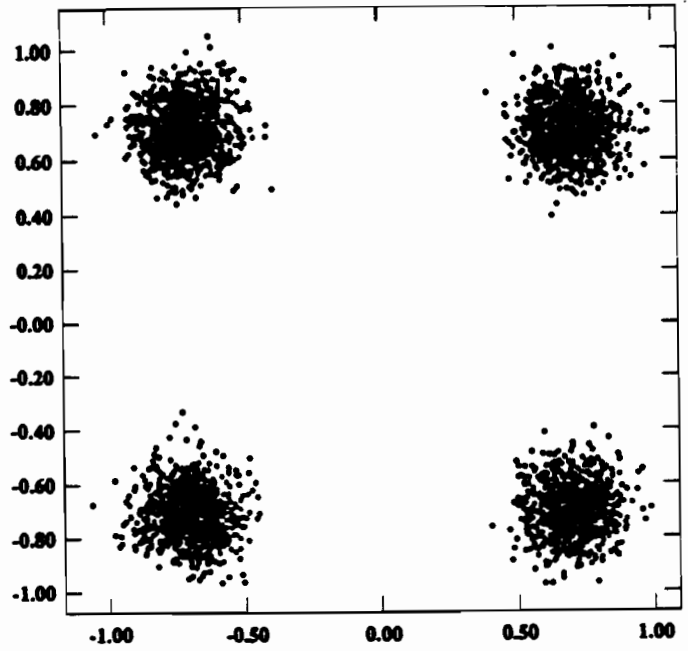
(l)

Figure 7.22 (continued) (j) 16-QAM radio: multipath fade (+6 dB slope, 50 MHz to 90 MHz). (k) 16-QAM radio: multipath fade (6 dB symmetrical notch, 70 MHz). (l) 16-QAM radio: multipath fade (10 dB symmetrical notch, 70 MHz).

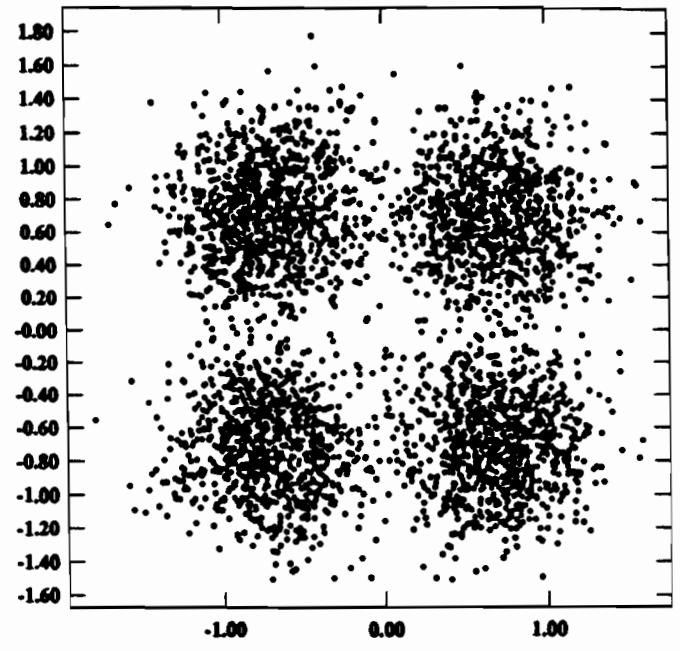
**TABLE 7.7 16-QAM Degradation
Allocation at BER = 10^{-8} (from [Bates,
7.7] with permission of the IEEE)**

16 QAM	
Quadrature carrier offset	
amplitude imbalance	9%
Nyquist filters	15%
Timing recovery	6%
Carrier recovery	21%
Power amplifier	15%
Radio	2%
Total eye closure	68%
Resulting eye opening	32%

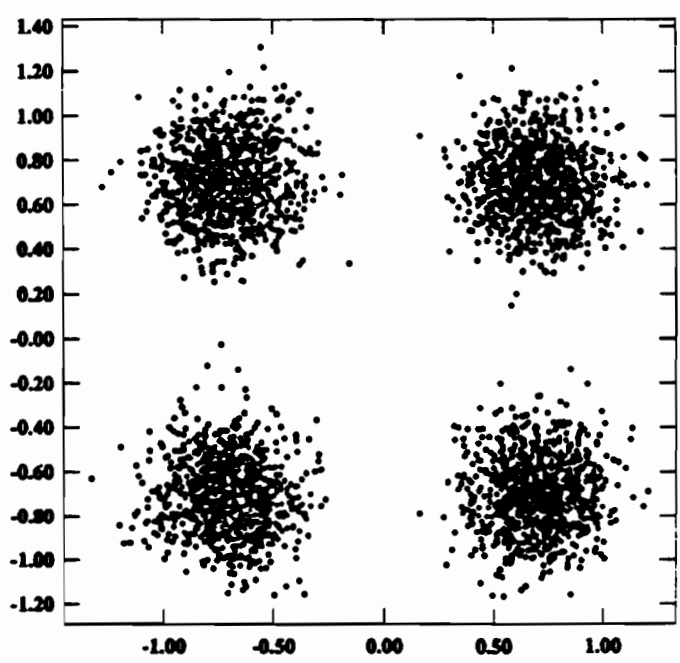
QPSK, 4K symbols, alpha=0.35, Eb/No=15dB



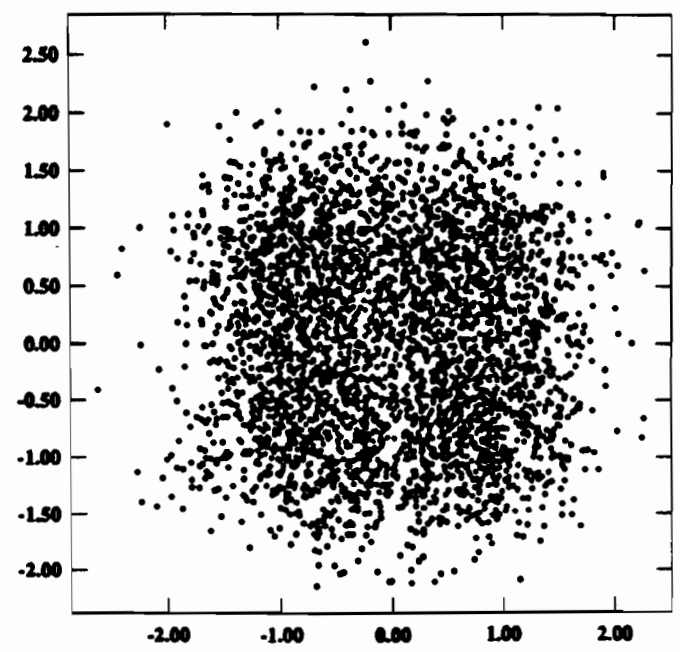
QPSK, 4K symbols, alpha=0.35, Eb/No=5dB



QPSK, 4K symbols, alpha=0.35, Eb/No=10dB

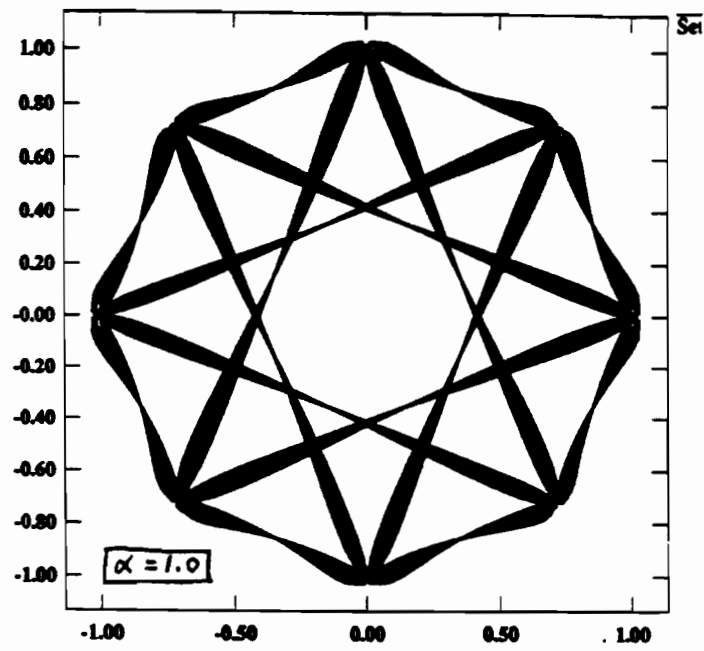
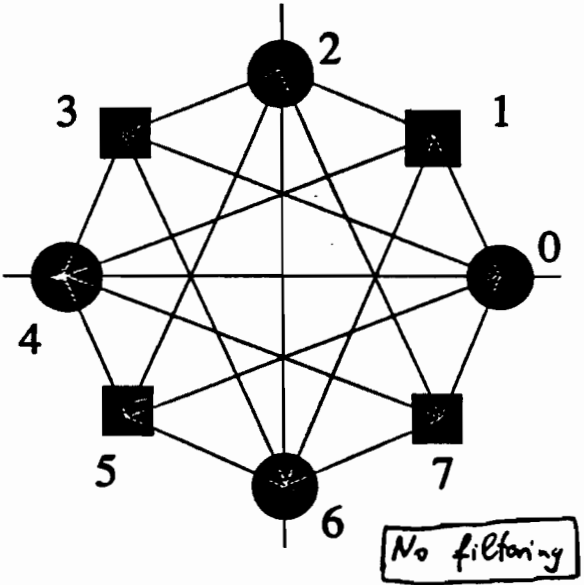


QPSK, 4K symbols, alpha=0.35, Eb/No=0dB



$\frac{\pi}{4}$ -QPSK : State-space diagrams

$\Delta\theta \in \{\frac{\pi}{4}, \frac{3\pi}{4}, \frac{5\pi}{4}, \frac{7\pi}{4}\}$



$\theta_i = \theta_{i-1} + \Delta\theta_i$

\uparrow
Tx-phase

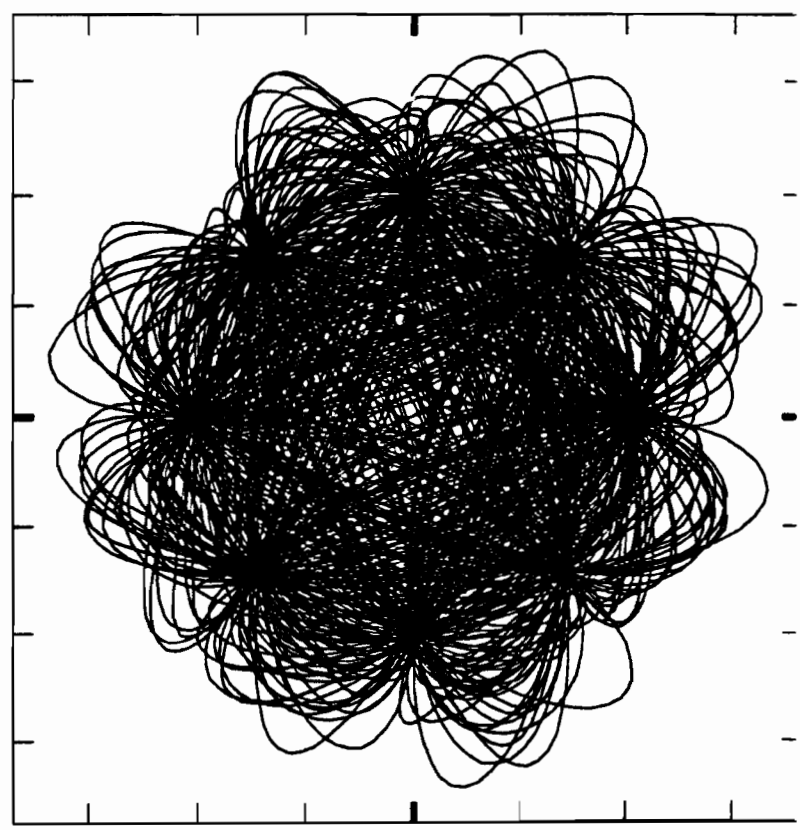
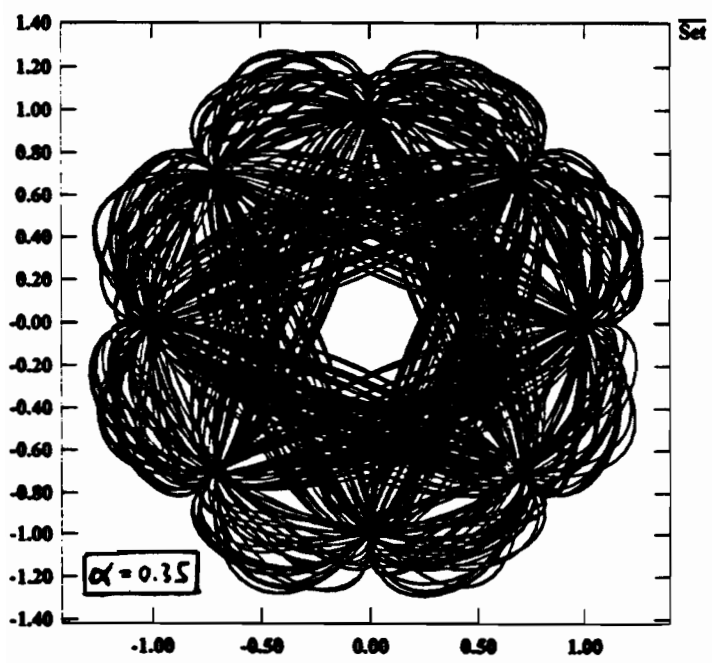
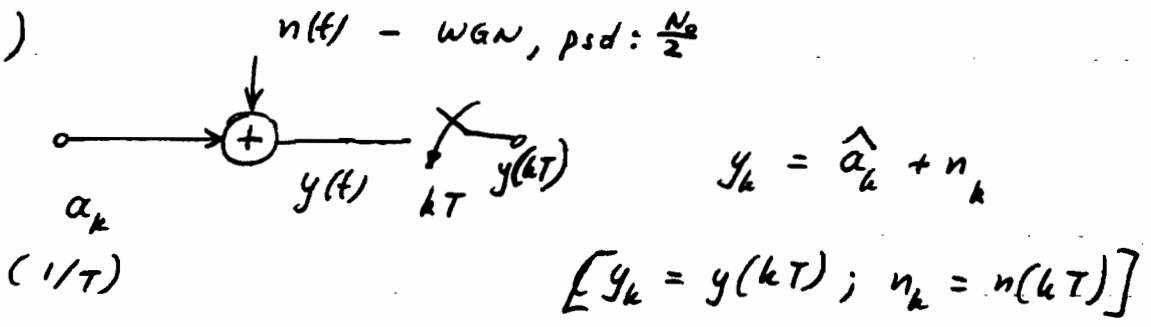


Figure 10 State-space diagram of a computer-simulated $\pi/4$ -QPSK system employing raised cosine filters with an excess bandwidth of 35% and which is operated in the presence of CCI (C/I = 55 dB) and Gaussian noise (C/N = 40 dB). The number of iterations is equal to 4.

ADDITIVE WHITE GAUSSIAN NOISE (AWGN):
FILTER CONSIDERATIONS

Example 1

(No filter)

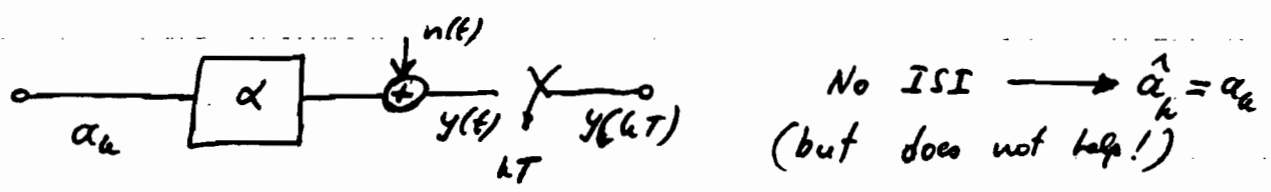


- ① No filtering $\Rightarrow \hat{a}_k = a_k$
- ② Power of $n_k = P_n = E\{|n_k|^2\} = \int_{-\infty}^{\infty} \frac{N_0}{2} df \Rightarrow \infty$

$\therefore a_k$ is totally hidden in noise !!

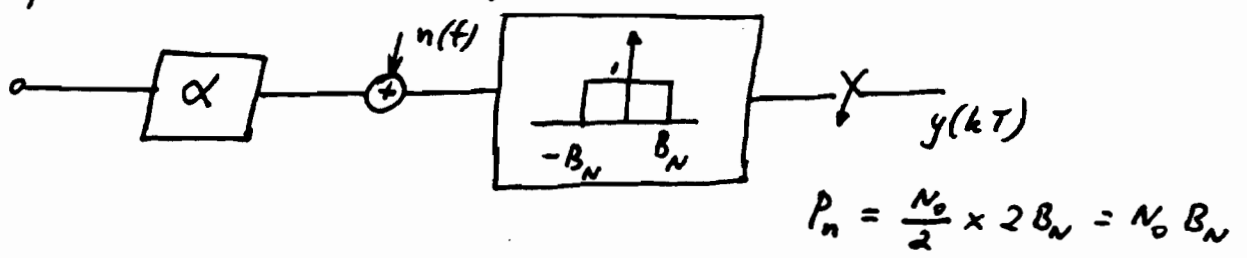
Example 2

(Nyquist filter at the Tx)



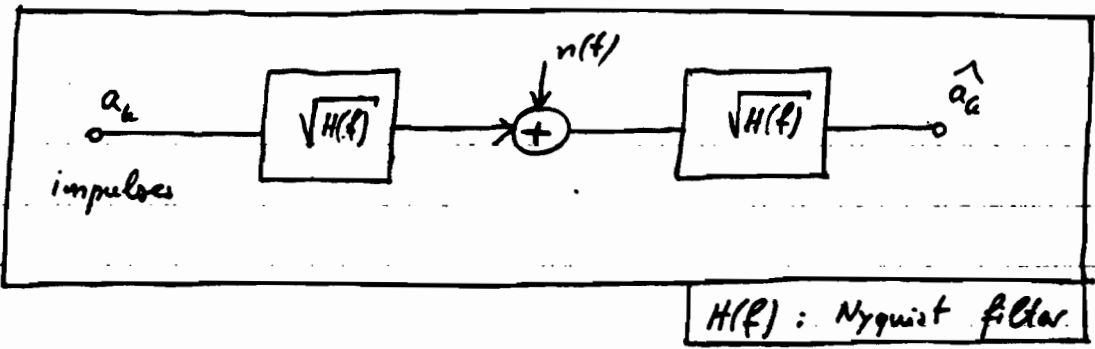
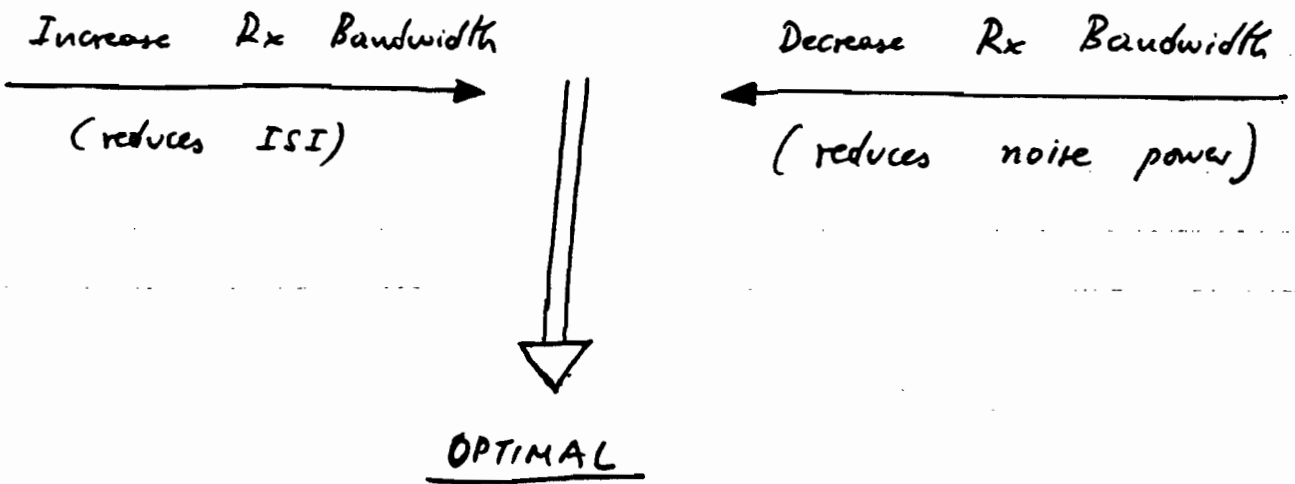
Example 3

(Rx filter to reduce noise)



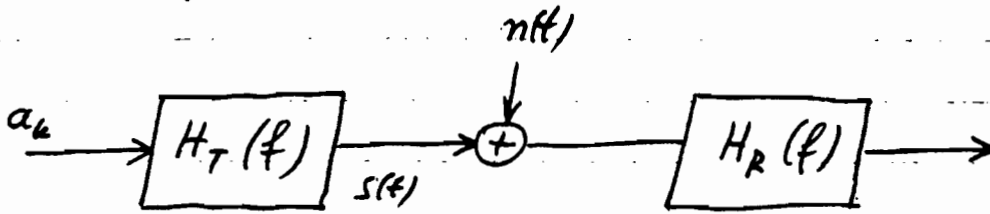
\therefore The noise power which corrupts the received sample depends upon B_N .

BUT Rx filter reduces the power of the noise at the same time it should not distort (in the Nyquist sense) i.e., should not create any ISI at the sampling instant.



(Optimal: in the sense that the signal-to-noise ratio is maximized.)

Proof



$$h_T(t) \leftrightarrow H_T(f)$$

$$h_R(t) \leftrightarrow H_R(f)$$

$$\frac{1}{T} = \text{symbol rate}$$

$$H(f) = H_T(f) H_R(f) = \begin{cases} \text{Nyquist} \\ \text{Filter} \end{cases}$$

$$s(t) = \sum_k \alpha_k h_T(t - kT)$$

$$P_s = \left\langle \lim_{K \rightarrow \infty} \frac{1}{2KT} \int_{-KT}^{KT} \left[\sum_{k=-K}^K \alpha_k h_T(t - kT) \right]^2 dt \right\rangle$$

$$= \lim_{K \rightarrow \infty} \frac{1}{2KT} \sum_n \sum_m \langle a_n a_m \rangle \int_{-KT}^{KT} h_T(t - nT) h_T(t - mT) dt$$

a_n, a_m are uncorrelated for $n \neq m$, i.e.,

$$\langle a_n a_m \rangle = \begin{cases} 0 & n \neq m \\ \frac{1}{a^2} & n = m \end{cases}$$

$$\therefore P_s = \lim_{K \rightarrow \infty} \frac{\overline{a^2}}{2KT} \sum_{k=-K}^K \int_{-kT}^{kT} h_T^2(t-kT) dt$$

$$= \frac{\overline{a^2}}{T} \int_{-\infty}^{\infty} h_T^2(t) dt$$

Parseval's theorem

$$\int_{-\infty}^{\infty} h^2(t) dt = \int_{-\infty}^{\infty} |H(f)|^2 df$$

$$\therefore P_s = \frac{\overline{a^2}}{T} \int_{-\infty}^{\infty} |H_T(f)|^2 df$$

Since $H(f) = H_T(f) H_R(f)$ has to satisfy Nyquist \Rightarrow the signal power at the output of the Rx is invariant under changes of $H_R(f)$.

\therefore The signal-to-noise ratio is maximized by minimizing the noise power at the output of the Rx filter.

$$P_N = \int_{-\infty}^{\infty} N(f) |H_R(f)|^2 df$$

under the conditions that

$$\textcircled{1} \quad P_S = \frac{a^2}{T} \int_{-\infty}^{\infty} |H_T(f)|^2 df$$

$$\textcircled{2} \quad P_R = P_S \int |H_R(f)|^2 df = \text{constant}$$

Using techniques of the variational analysis we can form the following functional V

$$V = \lambda P_S + P_N$$

λ = Lagrange multiplier.

$$V = \int_{-\infty}^{\infty} \left[\lambda \frac{\bar{a}^2}{T} \frac{|H(f)|^2}{|H_R(f)|^2} + N(f) |H_R(f)|^2 \right] df$$

Minimization of V is achieved on the stationary point of the above equation, i.e.,

$$H_R(f) = \frac{\sqrt{H(f)}}{N^{1/4}(f)}$$

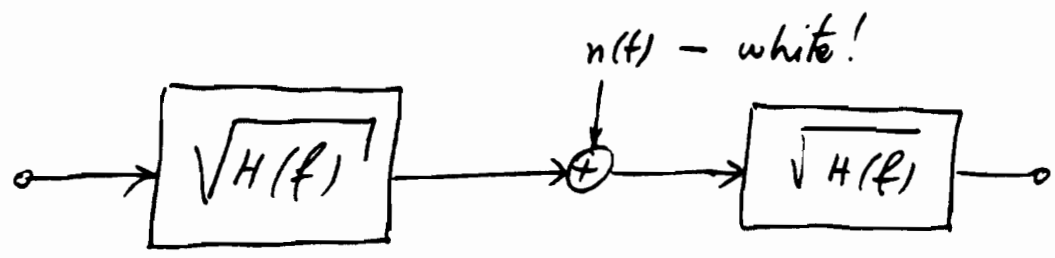
so that

$$H_T(f) = \frac{H(f) N^{1/4}(f)}{\sqrt{H(f)}} \quad (\text{real filters})$$

For flat noise (const. p) $\Rightarrow N^{1/4}(f) = 1$

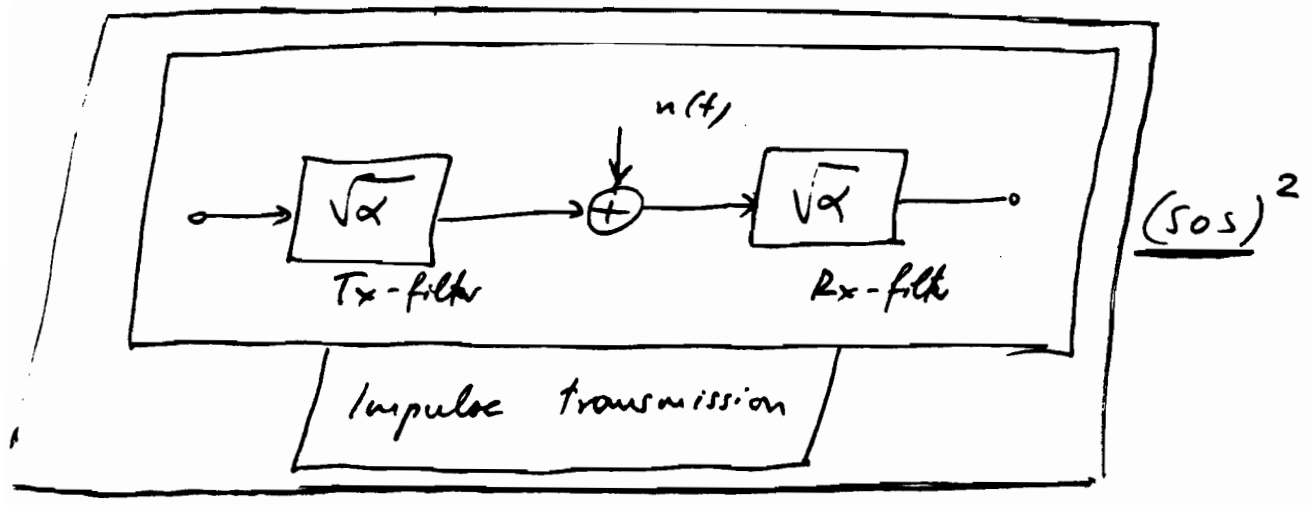
$$H_R(f) = H_T(f) = \sqrt{H(f)}$$

∞



$H(f)$: Nyquist filter

In practice we use raised cosine filter α .



For pulse transmission an extra $\frac{f}{\sin f} = H_E(f)$ filter is needed.

ORTHOGONAL SIGNALS, MATCHED FILTER & CORRELATION RECEIVERS AND ERROR PROBABILITY PERFORMANCE IN AN AWGN CHANNEL

• ORTHOGONAL SIGNALS

Consider N orthonormal functions $\phi_q(t)$, $q = 1, 2, \dots, N$ over the interval $(0, T]$

$$\int_0^T \phi_i(t) \phi_k(t) dt = \delta_{ik} = \begin{cases} 1 & \text{for } i=k \\ 0 & \text{for } i \neq k \end{cases}$$

These functions can be used to define a set of M finite energy signals

$$s_i(t) = \sum_{j=1}^N s_{ij} \phi_j(t) \quad \begin{matrix} i = 1, 2, \dots, M \\ 0 < t \leq T \end{matrix}$$

with $s_{ij} = \int_0^T s_i(t) \phi_j(t) dt$

∴ $s_i(t)$ is completely specified by $\bar{s}_i = (s_{i1}, s_{i2}, \dots, s_{iN})$ assuming $\{\phi_j(t)\}$ are known.

$$\begin{aligned} \text{Energy} = E_i &= \int_0^T s_i^2(t) dt = \int_0^T \sum_{j=1}^N \sum_{k=1}^N s_{ij} s_{ik} \phi_j(t) \phi_k(t) dt = \\ &= \sum_{j=1}^N \sum_{k=1}^N s_{ij} s_{ik} \delta_{jk} = \sum_{j=1}^N s_{ij}^2 \end{aligned}$$

Alternatively $E_i = \bar{s}_i \cdot \bar{s}_i = |\bar{s}_i|^2$

• $\hat{=}$ dot product ; $\bar{s}_i \cdot \bar{s}_j \hat{=} |\bar{s}_i| |\bar{s}_j| \cos \theta_{ij}$

Correlation coefficient between the i -th and the j -th signal

$$\rho_{ij} \hat{=} \frac{\int_0^T s_i(t) s_j(t) dt}{\sqrt{E_i E_j}} = \frac{\sum_{k=1}^N s_{ik} s_{jk}}{\sqrt{E_i E_j}} = \frac{\bar{s}_i \cdot \bar{s}_j}{\sqrt{E_i E_j}}$$

Alternatively: $\rho_{ij} = \cos \theta_{ij}$

θ_{ij} is the angle between the vectors \bar{s}_i and \bar{s}_j .

Euclidean distance

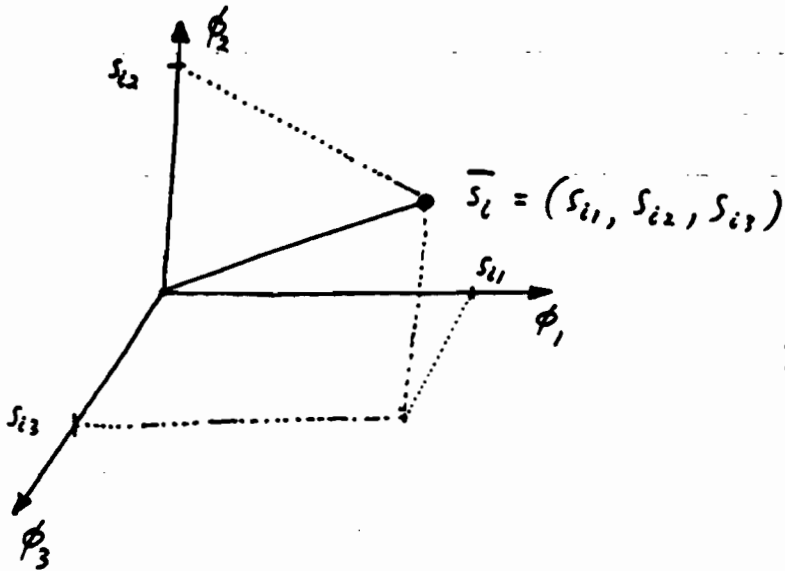
$$d_{ij}^2 \hat{=} \int_0^T [s_i(t) - s_j(t)]^2 dt = \sum_{k=1}^N s_{ik}^2 + \sum_{k=1}^N s_{jk}^2 - 2 \sum_{k=1}^N s_{ik} s_{jk}$$

$$= E_i + E_j - 2 \rho_{ij} \sqrt{E_i E_j}$$

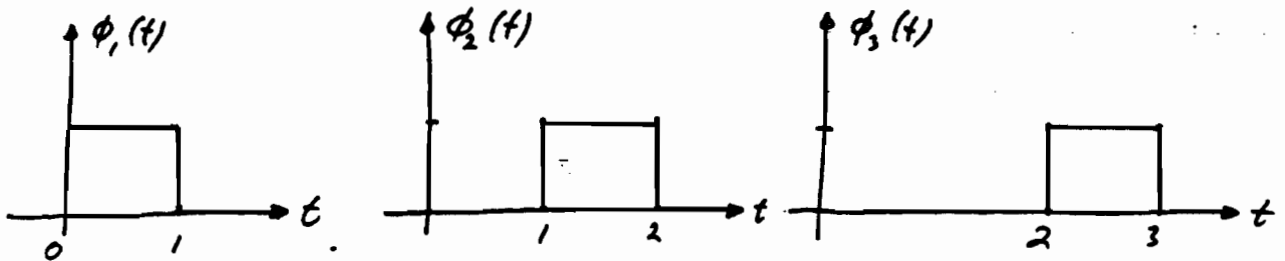
Examples

$$N = 3$$

(a)



(b)



$$N = 2 \rightarrow 2\text{-dimensions}$$

$$\phi_1(t) = \sqrt{\frac{2}{T}} \cos(2\pi f_c t)$$

$$\phi_2(t) = \sqrt{\frac{2}{T}} \sin(2\pi f_c t)$$

$$0 \leq t \leq T$$

2-dim. \rightarrow practically any signal representation
(PSK, QAM, QPSK...)

$$s_i(t) = s_{i1} \sqrt{\frac{2}{T}} \cos(2\pi f_c t) - s_{i2} \sqrt{\frac{2}{T}} \sin(2\pi f_c t)$$

$$= \sqrt{\frac{2E}{T}} \cos(2\pi f_c t + \theta_i) \quad E = s_{i1}^2 + s_{i2}^2$$

$$\theta_i = \arctan \frac{s_{i2}}{s_{i1}}$$

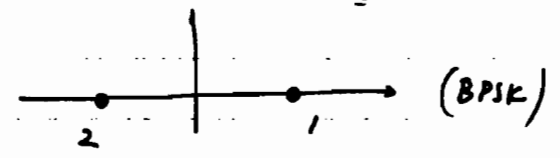
$$\rho_{ij} = \frac{1}{E} \int_0^T s_i(t) s_j(t) dt = \dots = \cos(\theta_i - \theta_j) \quad (\text{assuming that } f_c = k \frac{1}{T})$$

For M-ary PSK systems $\Rightarrow \theta_i = \frac{2\pi}{M} i \Rightarrow \rho_{ij} = \cos\left[\frac{2\pi(i-j)}{M}\right]$

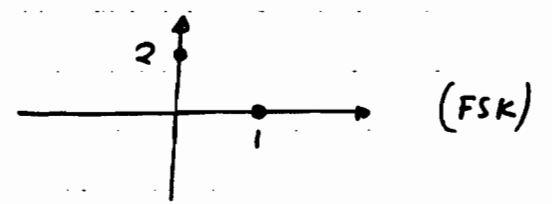
Examples

M=2

$\rho_{12} = -1$ (antipodal)



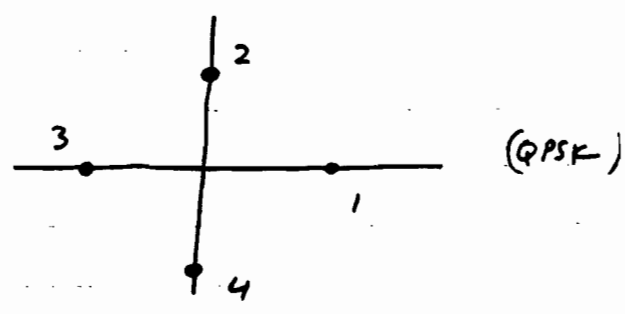
$\rho_{12} = 0$ (orthogonal)



M=4

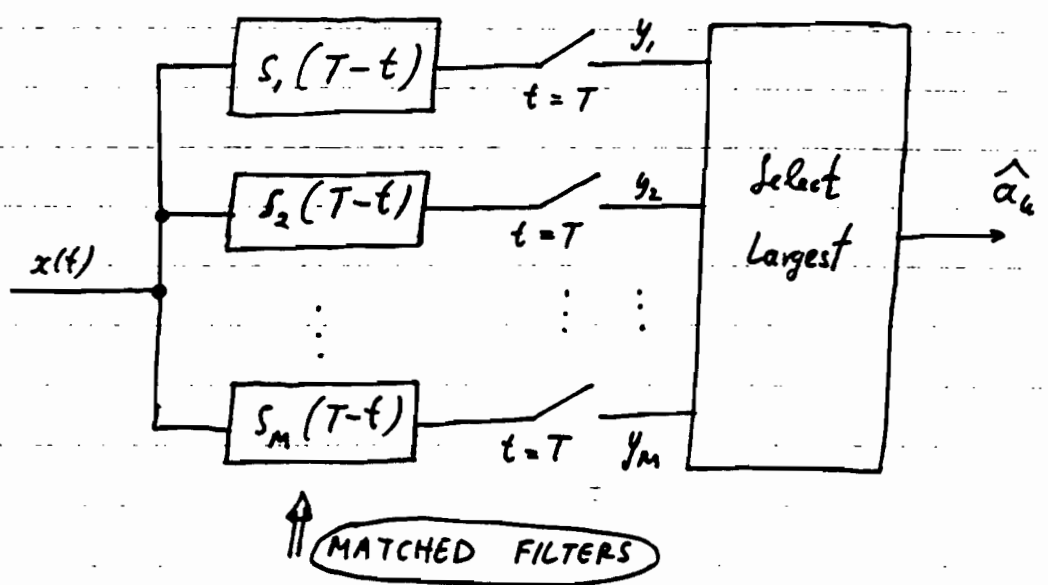
$$\rho_{12} = \rho_{23} = \rho_{34} = \rho_{41} = 0$$

$$\rho_{13} = \rho_{24} = -1$$

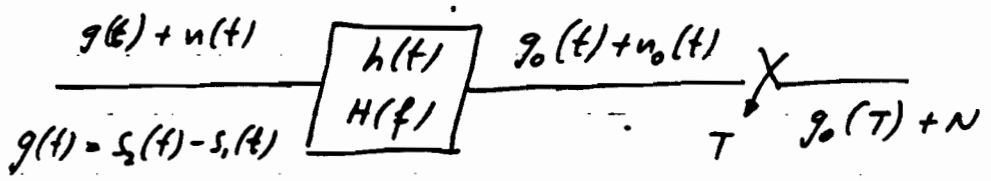


MATCHED FILTER RECEIVERS

- AWGN channel
- M-ary signals $\rightarrow s_1(t), s_2(t), \dots, s_M(t)$ (equiprobable)
- It can be shown that the following receiver structure minimizes the error probability



(For $M=2$ we have



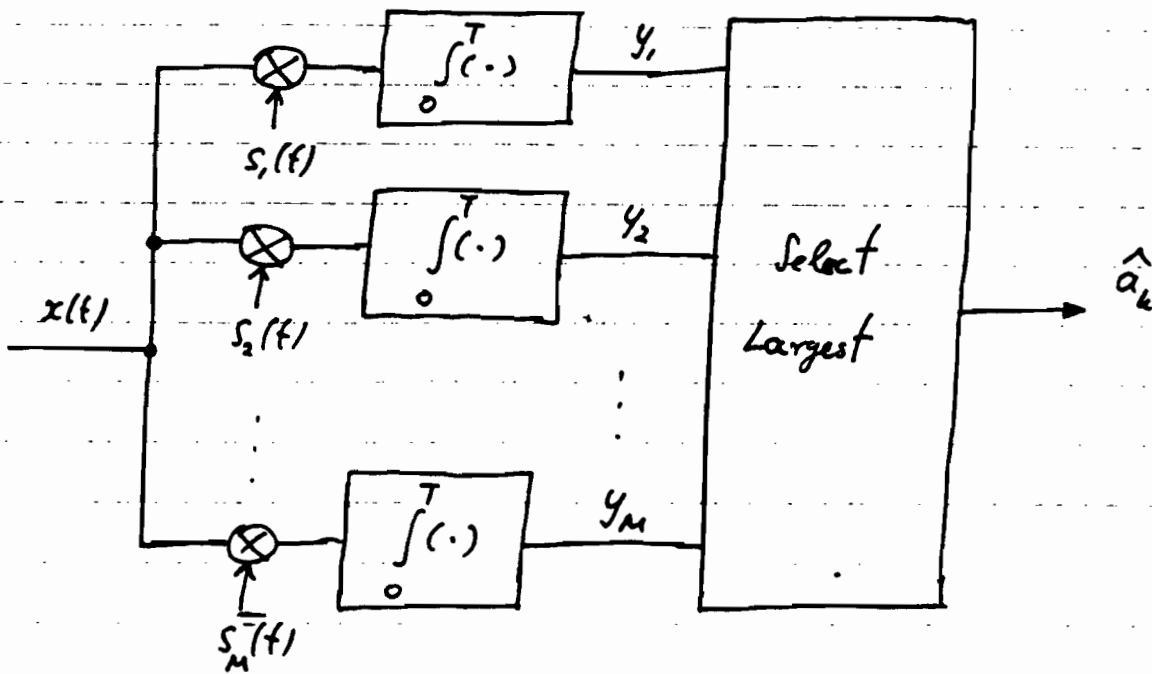
Find $h(t)$ which maximizes $J = \frac{s_{02}(T) - s_{01}(T)}{\sigma}$

Proof

(Hw)

• CORRELATION RECEIVERS

It is easy to show that an equivalent to the "Matched Filter" receiver is the following Correlation Receiver



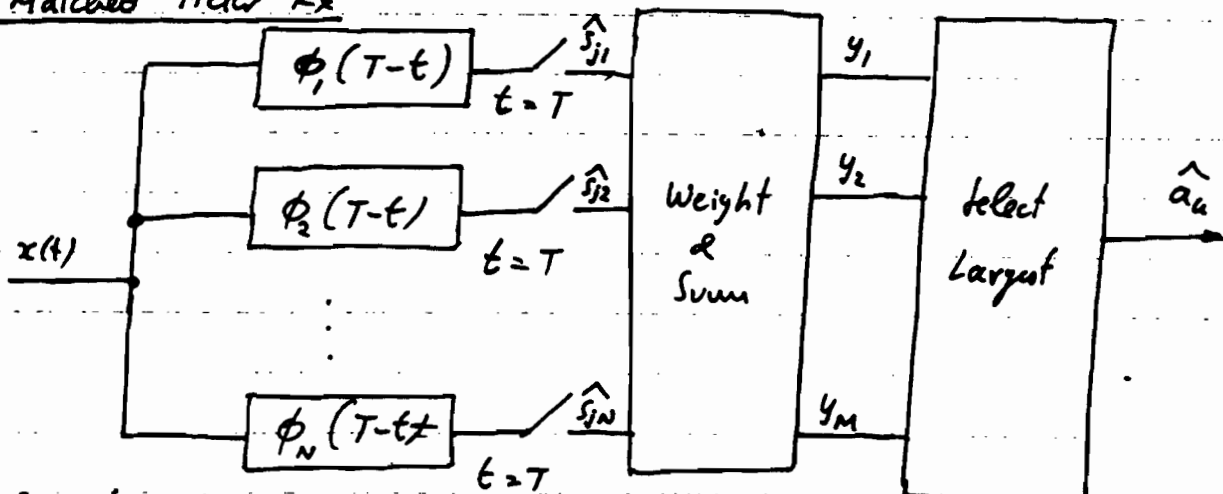
Although its the "same thing", it might have implementation advantages.

For almost any "practical" digital communication system the transmitted signals can be represented by two orthonormal functions e.g., $\phi_1(t)$ & $\phi_2(t)$.

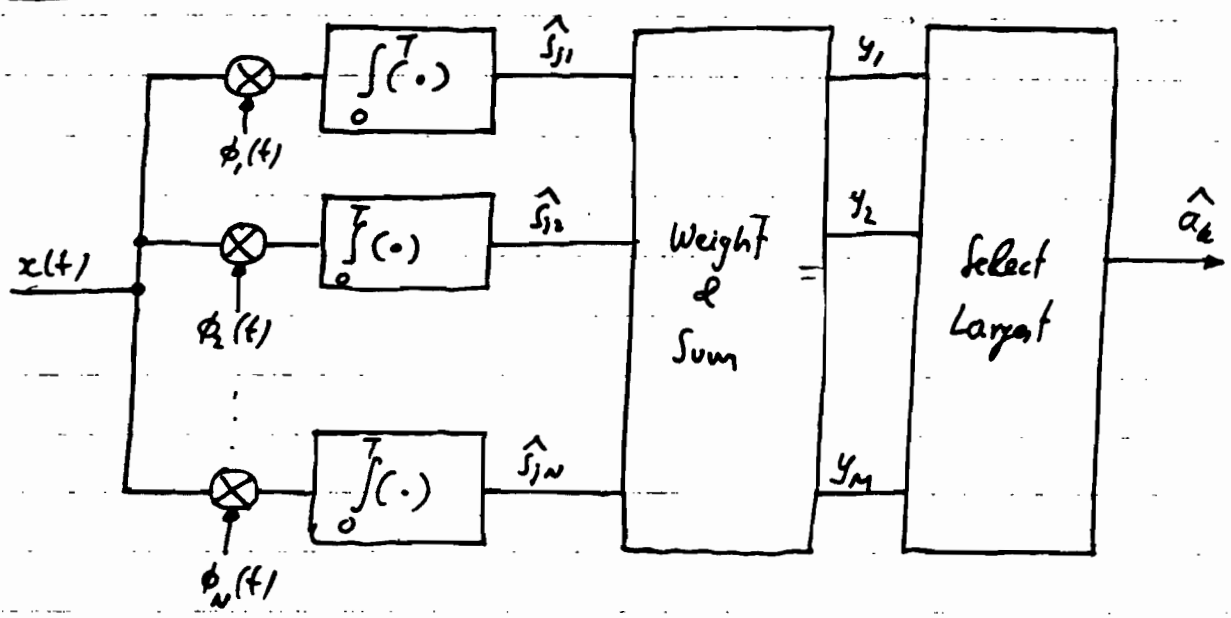
⇒ The previous receivers can be very much simplified by using these two orthonormal functions.

In more general terms: Assume N orthonormal functions.

• Matched Filter Rx



• Correlation Rx



$$s_i(t) = \sum_{j=1}^N s_{ij} \phi_j(t)$$

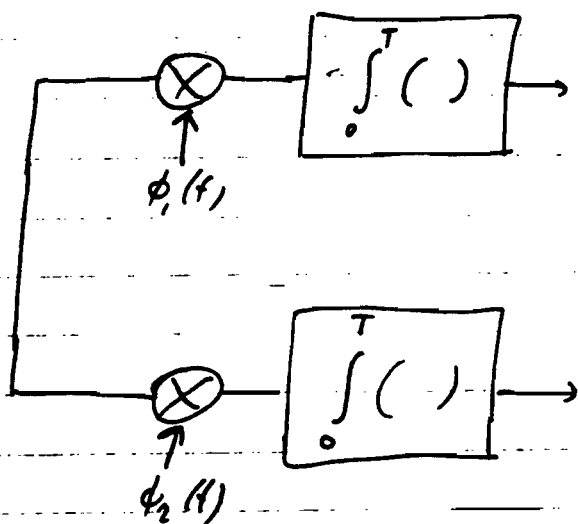
N correlation receivers which produce a set of N coefficients

$$\begin{aligned} \hat{s}_{jk} &\equiv \text{estimate of } s_{jk} \text{ when } s_j(t) \text{ is transmitted} \\ &= \int_0^T x(t) \phi_k(t) dt \end{aligned}$$

In order to obtain the signal samples y_i , the following mathematical operations are needed.

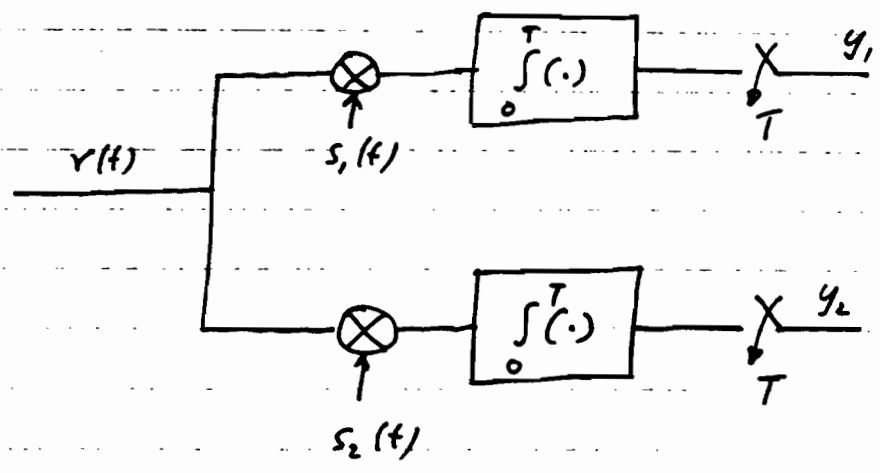
$$\begin{aligned} y_i &= \int_0^T x(t) s_i(t) dt = \int_0^T x(t) \sum_{k=1}^N s_{ik} \phi_k(t) dt \\ &= \sum_{k=1}^N s_{ik} \int_0^T x(t) \phi_k(t) dt = \sum_{k=1}^N s_{ik} \hat{s}_{jk} \quad i = 1, 2, \dots, M \end{aligned}$$

For QAM - only 2 functions \rightarrow \cos & \sin !!



PROBABILITY PERFORMANCE IN AN AWGN CHANNEL

Binary system : $s_i(t) \begin{cases} i = 1, 2 \\ 0 \leq t \leq T \\ \text{equiprobable} \end{cases} ; E = \int_0^T s_i^2(t) dt$



$$y_1 \begin{matrix} \geq \\ \leq \end{matrix} \begin{matrix} s_1(t) \\ s_2(t) \end{matrix}$$

$$\Downarrow$$

$$z = y_1 - y_2 \begin{matrix} \geq \\ \leq \end{matrix} \begin{matrix} s_1(t) \\ s_2(t) \end{matrix}$$

$$r(t) = s_i(t) + \underbrace{n(t)}_{\text{WGN}} \rightarrow R(z) = \frac{N_0}{2} \delta(z)$$

Assume that $s_1(t)$ is transmitted

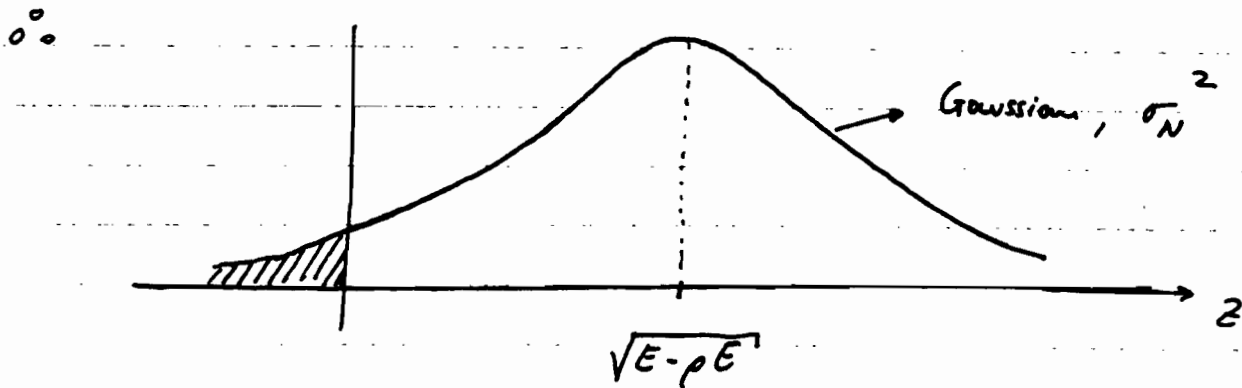
$$y_1 = \int_0^T r(t) s_1(t) dt = \int_0^T s_1^2(t) dt + \int_0^T s_1(t) n(t) dt$$

$$y_2 = \int_0^T r(t) s_2(t) dt = \int_0^T s_1(t) s_2(t) dt + \int_0^T s_2(t) n(t) dt$$

$$\begin{aligned} \therefore z = y_1 - y_2 &= \int_0^T s_1^2(t) dt - \int_0^T s_1(t) s_2(t) dt + \int_0^T n(t) [s_1(t) - s_2(t)] dt \\ &= \underline{E - \rho E + N} \end{aligned}$$

$$\rho \triangleq \text{correlation coefficient} = \frac{1}{E} \int_0^T s_1(t) s_2(t) dt$$

$N = \int_0^T n(t) [s_1(t) - s_2(t)] dt$ is a G.R.V. with zero mean and variance $\sigma_N^2 = E(N^2) = \dots = E(1-\rho) N_0$.



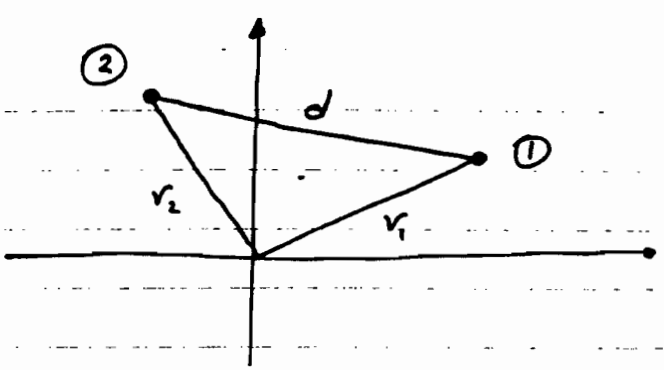
Prob. of an error = Prob that $z < 0$

∴

$$P_e = \int_{-\infty}^0 \frac{1}{\sqrt{2\pi} \sigma_N} \exp\left[-\frac{z - \sqrt{E - \rho E}}{2\sigma_N^2}\right] dz$$

$$P_e = \frac{1}{2} \operatorname{erfc}\left(\sqrt{\frac{E(1-\rho)}{2N_0}}\right)$$

A Geometrical Interpretation



signal ① has energy E_1
 - ② - - - E_2

$$r_1 = \sqrt{E_1}$$

$$r_2 = \sqrt{E_2}$$

$$d^2 = E_1 + E_2 - 2\rho\sqrt{E_1 E_2}$$

For equal energy signals, i.e., $E_1 = E_2 = E$

$$d^2 = 2E(1-\rho) \Rightarrow d = \sqrt{2} \sqrt{E(1-\rho)}$$

It can be easily shown that for equiprobable ①, ②

$$P_e = Q\left(\frac{d/2}{\sigma}\right) \quad [-\operatorname{erfc}(x) = 2Q(\sqrt{2}x)]$$

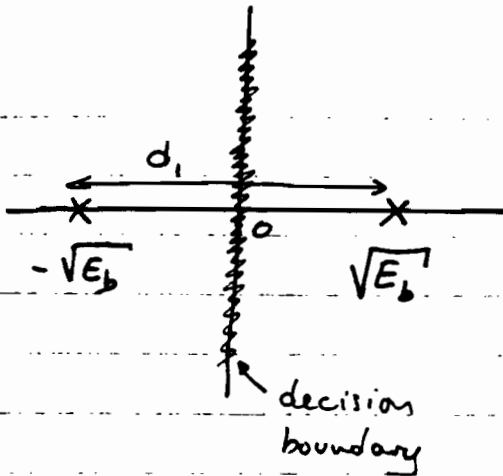
$\frac{d}{2}$ is half the ^{Euclidean} distance between ①, ② \rightarrow decision boundary

σ^2 is the noise variance = $\frac{N_0}{2}$

$$\therefore P_e = Q\left(\frac{d}{2\sigma}\right) = Q\left(\frac{\sqrt{2} \sqrt{E(1-\rho)}}{2 \sqrt{N_0/2}}\right) = Q\left(\sqrt{\frac{E(1-\rho)}{N_0}}\right)$$

Examples

(a)



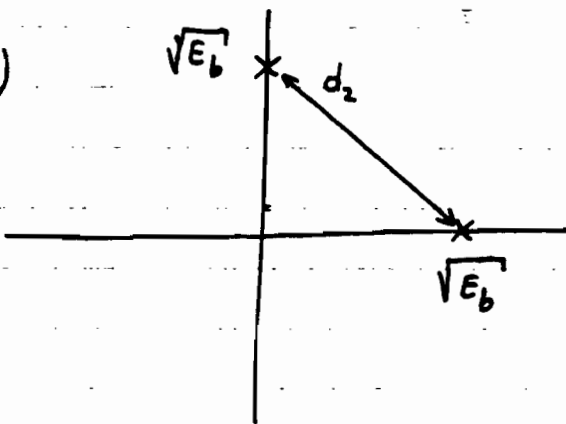
$\rho = -1 \Rightarrow$ antipodal
(BPSK)

$$\therefore P_e = \frac{1}{2} \operatorname{erfc}\left(\sqrt{\frac{E_b}{N_0}}\right)$$

In terms of the Q-function: $d_1 = 2\sqrt{E_b}$

$$\therefore P_e = Q\left(\frac{d_1/2}{\sigma}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

(b)



$\rho = 0 \Rightarrow$ orthogonal
(FSK)

$$\therefore P_e = \frac{1}{2} \operatorname{erfc}\left(\sqrt{\frac{E_b}{2N_0}}\right)$$

Q-function approach: $d_2 = \sqrt{2}\sqrt{E_b}$

$$\therefore P_e = Q\left(\frac{d_2/2}{\sigma}\right) = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$

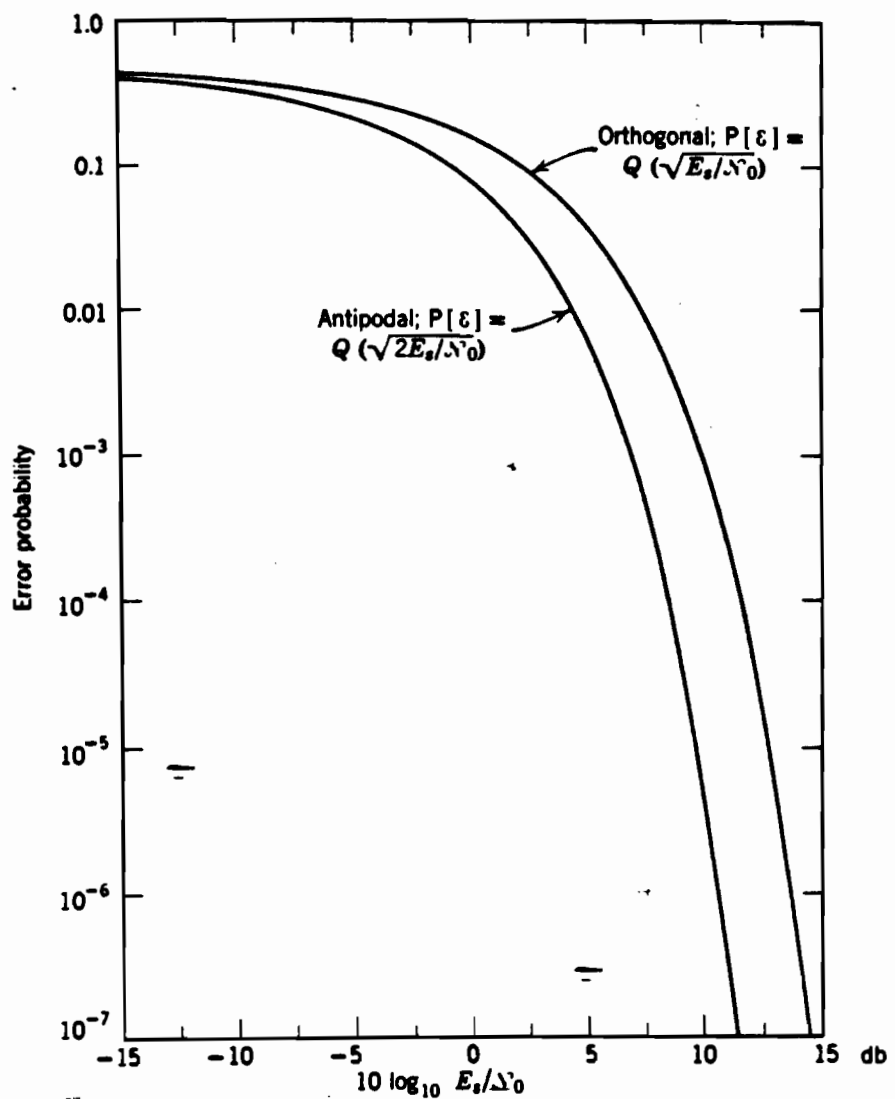


Figure 4.31 Probability of error for binary antipodal and binary orthogonal signaling with equally likely messages.

Ref. Wozencraft and Jacobs, Principles of Communication Engineering, J. Wiley, 1965

$$\frac{1}{\sqrt{2\pi\alpha}} e^{-\alpha^2/2} \left(1 - \frac{1}{\alpha^2}\right) < Q(\alpha) < \frac{1}{\sqrt{2\pi\alpha}} e^{-\alpha^2/2}; \quad \alpha > 0. \quad (2.121)$$

These two bounds are plotted together with $Q(\alpha)$ in Fig. 2.36.

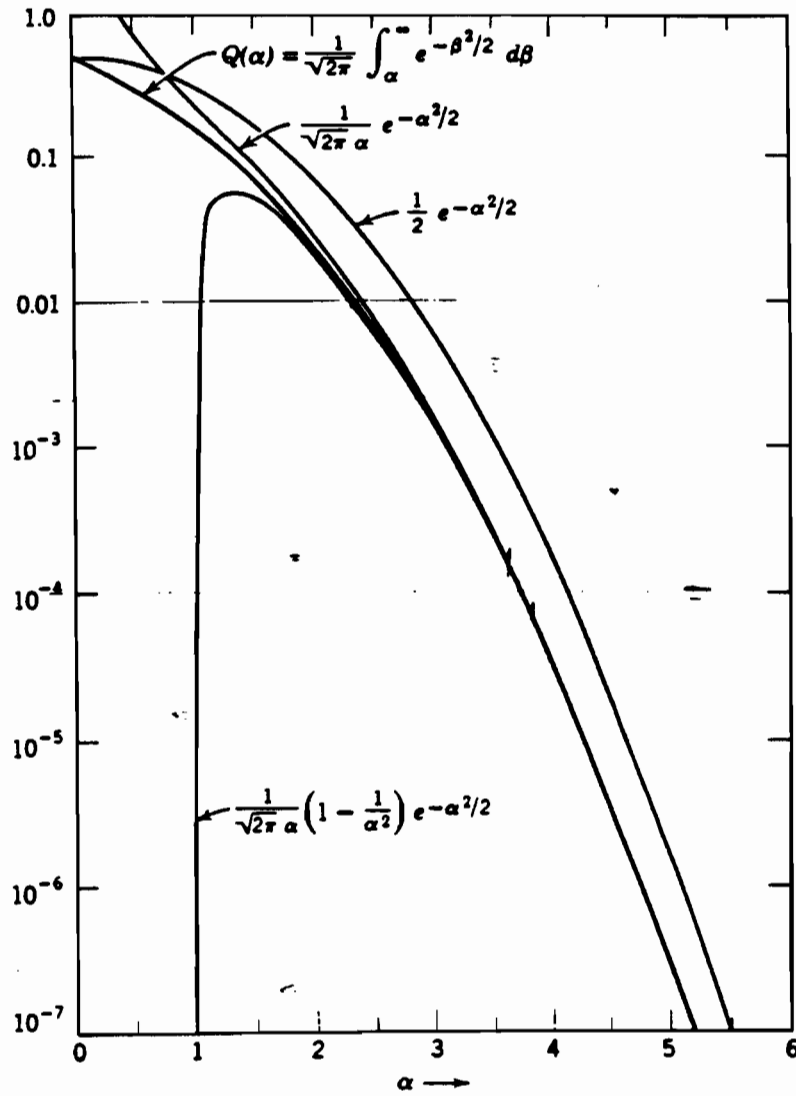
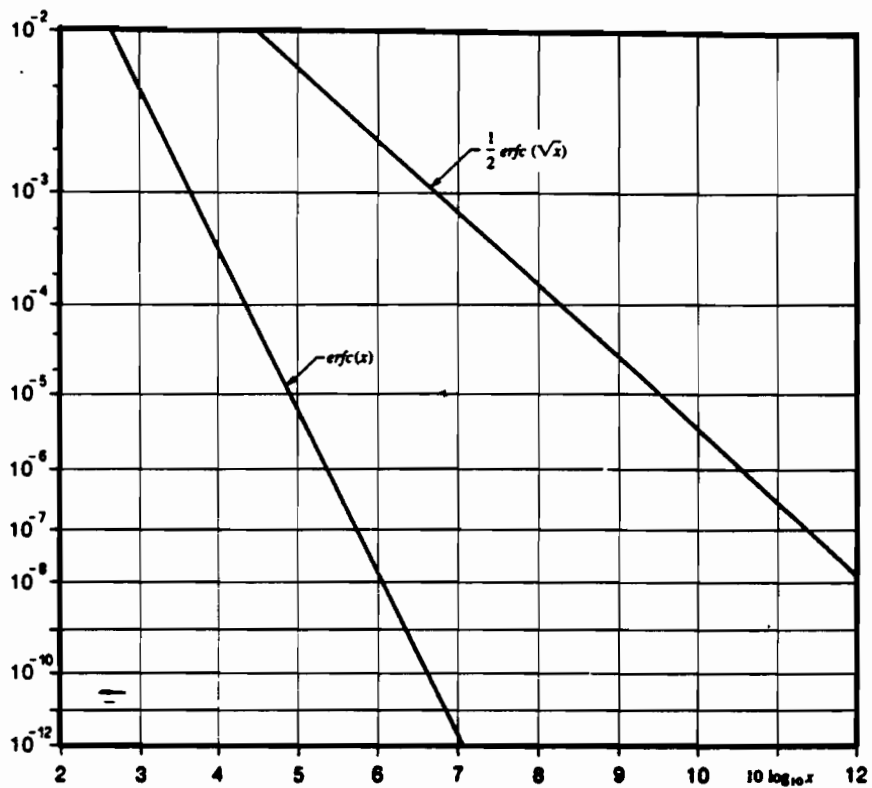
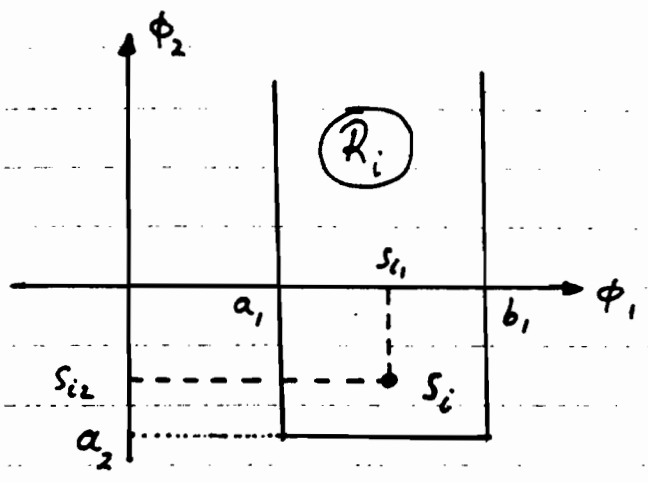


Figure 2.36 The function $Q(x)$ and three bounds.

Ref. W. + J.



Rectangular Signal Sets



Its always more efficient to find the Probability of having a correct decision P_c and then take $1 - P_c$ to find P_e .

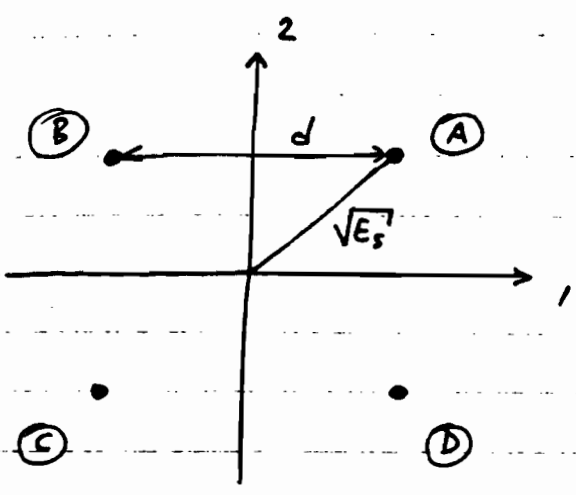
$$Pr(c/s_i) = Pr\{r \in R_i / s_i\} = Pr\{a_1 \leq r_1 \leq b_1, a_2 \leq r_2 < \infty / s_i\}$$

$$= Pr\{a_1 - s_{i1} \leq n_1 \leq b_1 - s_{i1}\} Pr\{a_2 - s_{i2} \leq n_2 < \infty\}$$

n_1, n_2 independent G.v.v., zero mean, with variance σ^2

$$\therefore Pr(c/s_i) = \int_{a_1 - s_{i1}}^{b_1 - s_{i1}} f_n(\alpha) d\alpha \int_{a_2 - s_{i2}}^{\infty} f_n(\alpha) d\alpha =$$

QPSK



All four points equiprobable, having the same energy E_s

Symmetry \Rightarrow only one point is needed to be examined

$$d^2 = 2E_s \Rightarrow d = \sqrt{2E_s}$$

$$\begin{aligned}
 Pr(\text{correct}/A) &= Pr\left\{ \frac{d}{2} < r_1 < \infty \text{ \& } -\infty < r_2 < -\frac{d}{2} \right\} \\
 &= Pr\left\{ n_1 < \frac{d}{2} \text{ \& } n_2 > -\frac{d}{2} \right\} = \underbrace{Pr\left\{ n_1 < \frac{d}{2} \right\}}_{1-p} \underbrace{Pr\left\{ n_2 > -\frac{d}{2} \right\}}_{1-p}
 \end{aligned}$$

$$p = Q\left(\frac{d/2}{\sigma}\right) = Q\left(\sqrt{\frac{E_s}{N_0}}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

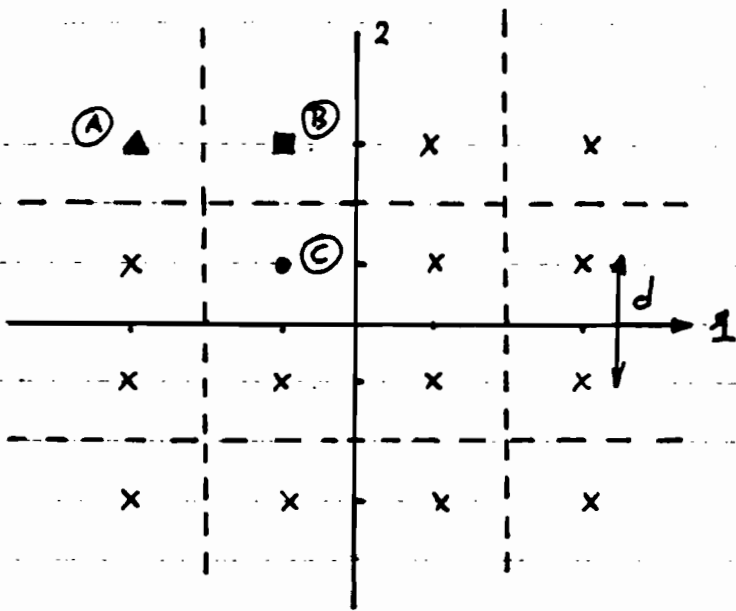
$$E_s = 2E_b$$

$$\left[\text{or } p = \frac{1}{2} \text{erfc}\left(\sqrt{\frac{E_b}{N_0}}\right) \right]$$

∴

$$P_e = 1 - Pr(\text{correct}/A) = 1 - (1-p)^2 = 2p - p^2 \approx 2p \quad (\text{for high } p > 10^{-2})$$

Performance of QPSK \equiv Performance of BPSK
 (2 b/s/Hz) 1 b/s/Hz

16-QAM

$$P_A = Pr(c/A)$$

$$P_B = Pr(c/B)$$

$$P_c = Pr(c/C)$$

$$P_A = (1-p)^2$$

$$P_B = (1-2p)(1-p)$$

$$P_c = (1-2p)^2$$

$$P_e = 1 - P_c = 1 - \frac{1}{M} \sum_{i=1}^M P(c/s_i)$$

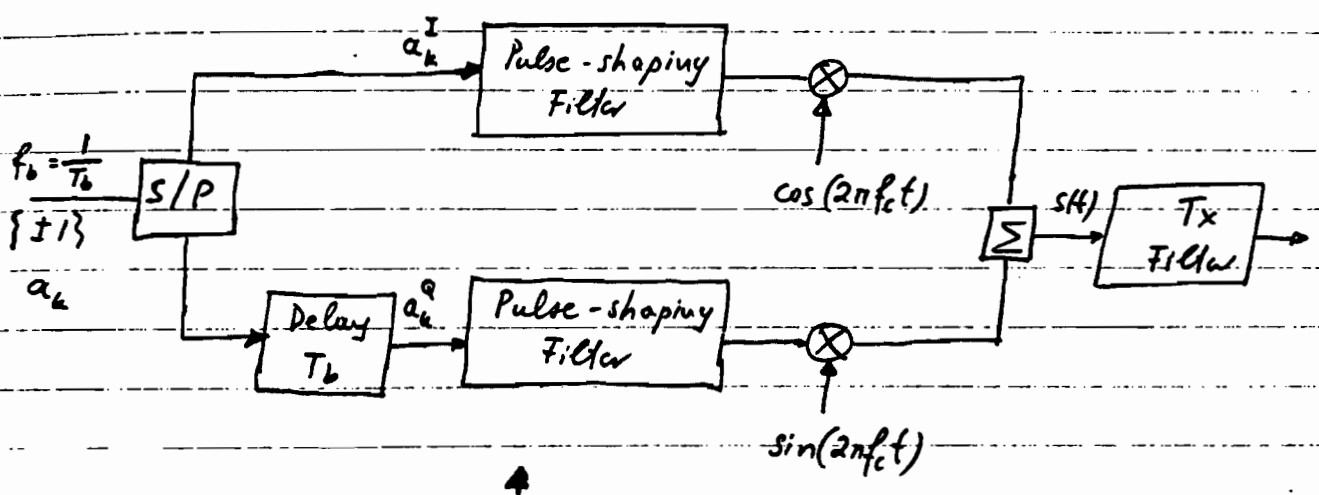
$$\therefore P_e = 1 - \frac{1}{16} [4P_A + 8P_B + 4P_c] = \dots = 1 - \frac{1}{4} (2-3p)^2$$

$$= 3p - \frac{3}{4} p^2$$

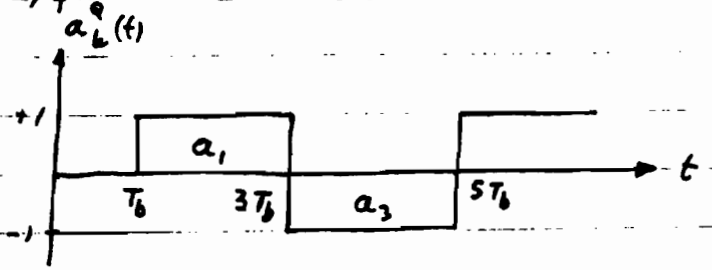
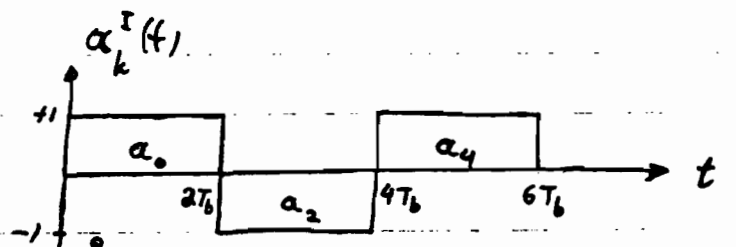
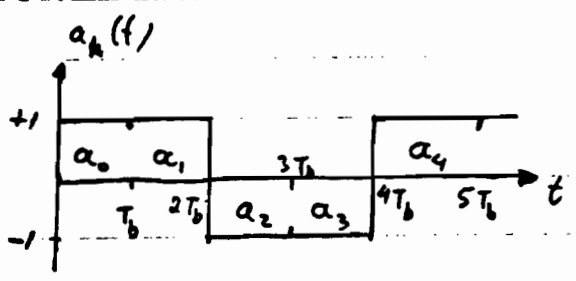
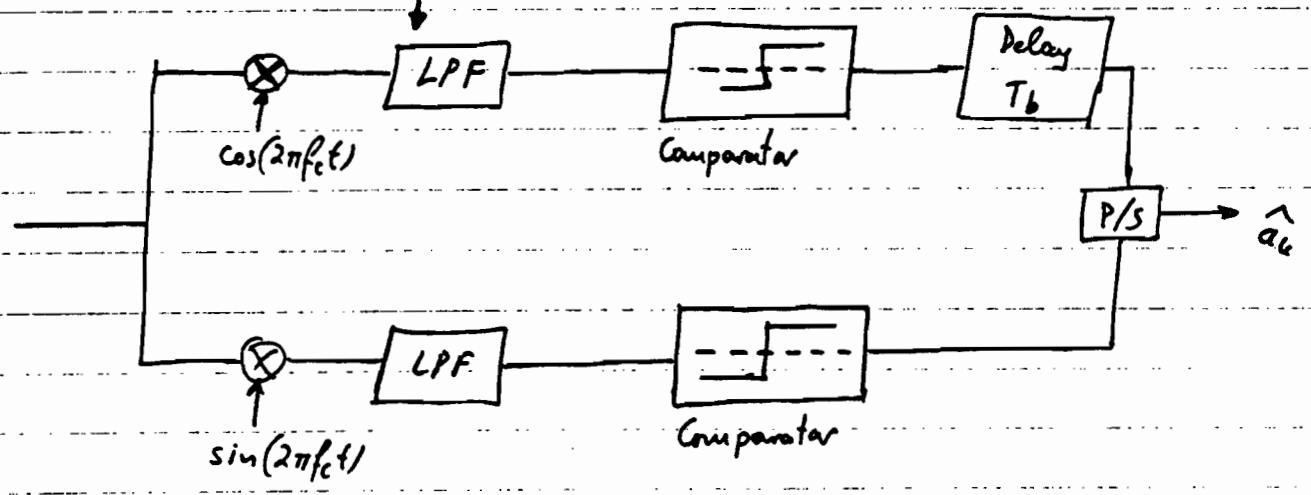
$$\text{where } p = Q\left(\frac{d/2}{\sigma}\right)$$

Offset - QPSK (OQPSK)

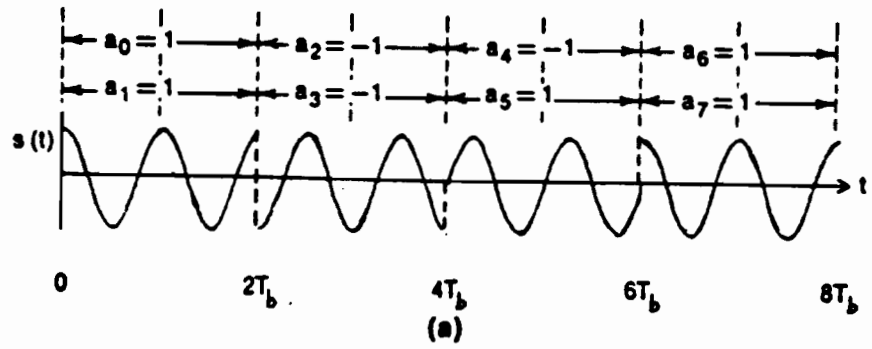
The I- and Q-channel are delayed by one bit duration



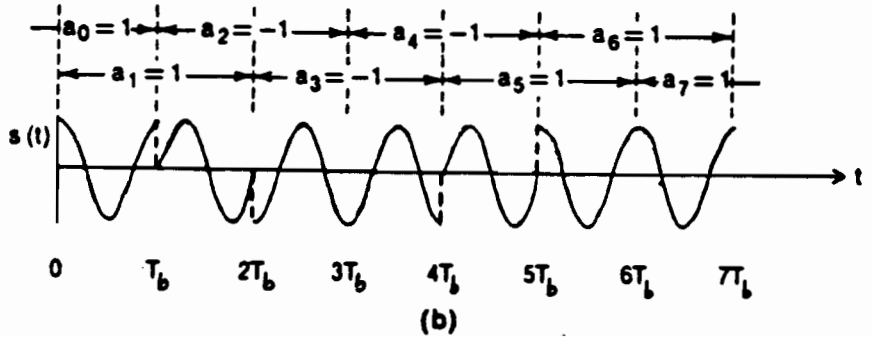
Transmitter
Receiver



Assuming no pulse-shaping filter, $s(t)$ is given by:



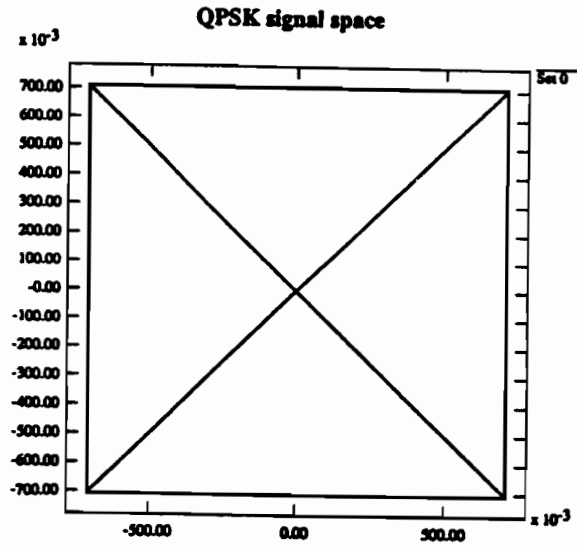
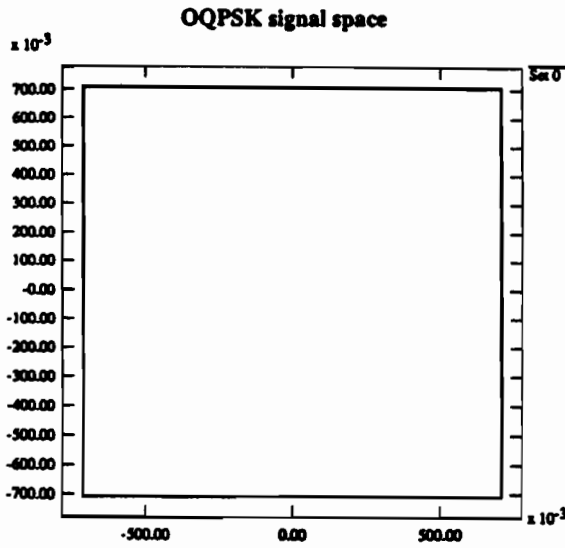
$$\Delta\phi \in \{0, \pm 90^\circ, 180^\circ\}$$



$$\Delta\phi \in \{0^\circ, \pm 90^\circ\}$$

(a) QPSK waveform; (b) OPQSK waveform.

State - Space Diagrams

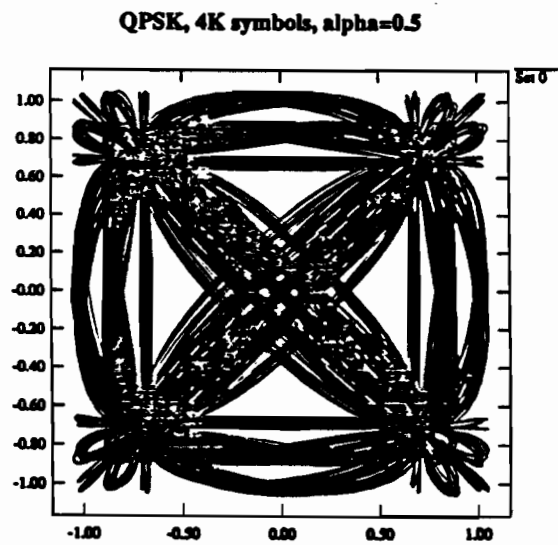
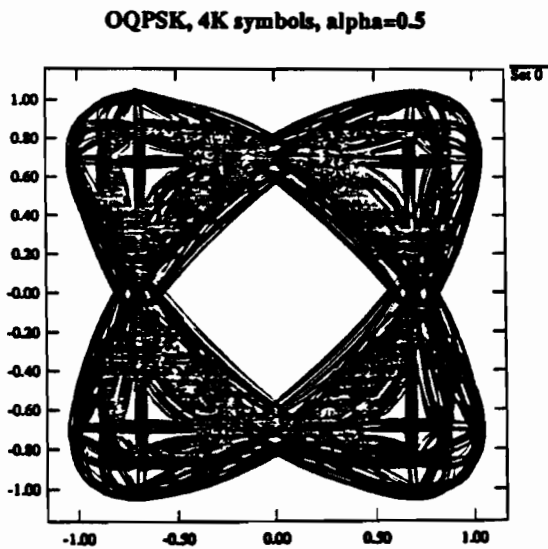


Comments

- Since the staggering does not change the orthogonality of the carriers \rightarrow in an AWGN channel OQPSK has the same performance as QPSK.
- The staggering does not change the power spectral density \rightarrow QPSK & OQPSK have the same spectrum (in a linear channel).
- For OQPSK there are no transitions through the origin. \Rightarrow Reduced envelope fluctuation \Rightarrow it has better spectral characteristics after bandlimiting & non-linear amplification.

Effects of bandlimiting

- State space diagrams



- In an AWGN channel the BER performance is "almost" the same
- The delay T_b has been shown to be optimal in terms of phase jitter immunity in an AWGN channel.
- When bandlimited OQPSK goes through a nonlinear amplifier (e.g., hardlimiter), the relatively small envelope fluctuations will be removed by the limiter. Furthermore, the absence of rapid phase shifts (\equiv high frequency content) means that this limiting will not regenerate high frequency components which originally were removed by the bandlimiting filter \Rightarrow out-of-band spectral regrowth (interference) will be small.

Typical system configuration (Tx):

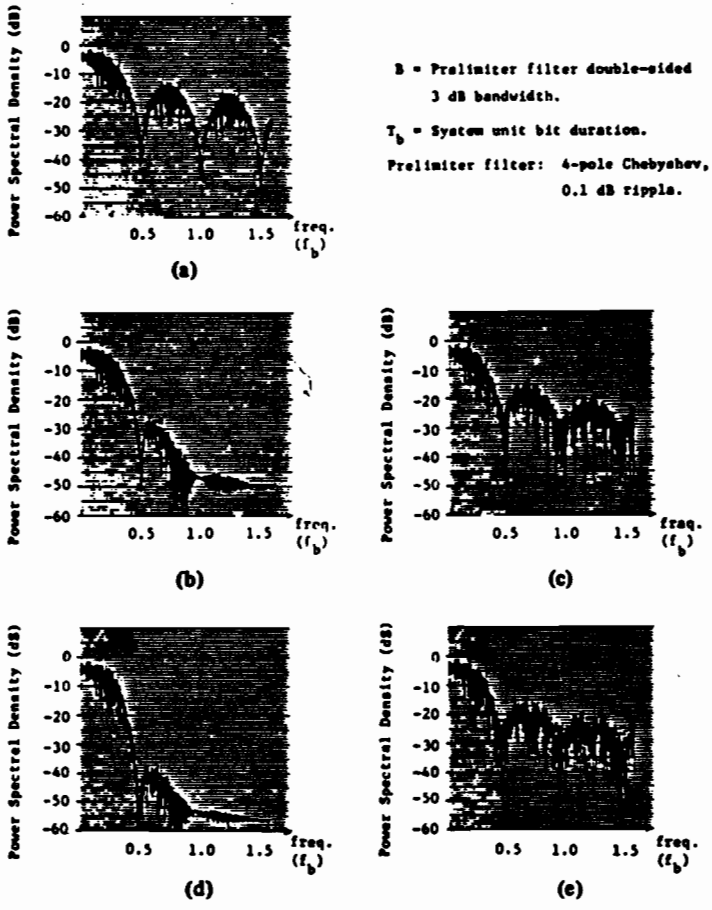
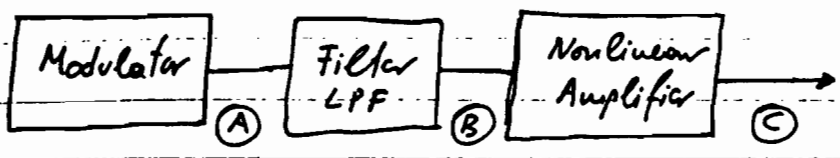


Fig. 2. Computed power spectral densities for QPSK. (a) Spectral density at pt. A, Fig. 1. Unfiltered, $BT_b = \infty$. (b) Spectral density at pt. B, Fig. 1. Filtered, $BT_b = 1$. (c) Spectral density at pt. C, Fig. 1. Filtered, $BT_b = 1$, then limited. (d) Spectral density at pt. B, Fig. 1. Filtered, $BT_b = 0.75$. (e) Spectral density at pt. C, Fig. 1. Filtered, $BT_b = 0.75$, then limited.

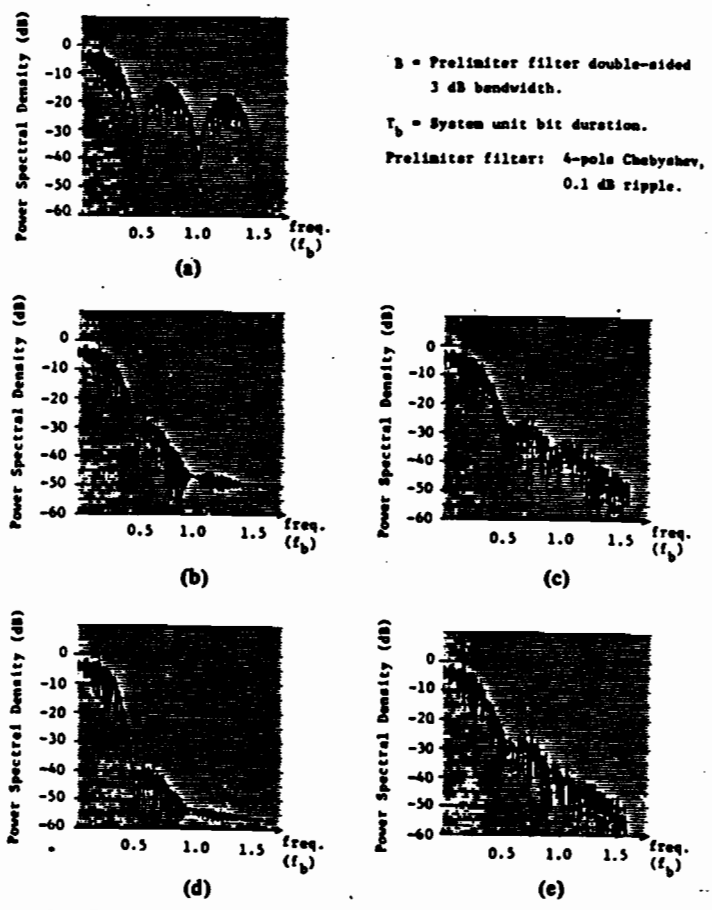


Fig. 3. Computed power spectral densities for offset QPSK. (a) Spectral density at pt. A, Fig. 1. Unfiltered, $BT_b = \infty$. (b) Spectral density at pt. B, Fig. 1. Filtered, $BT_b = 1$. (c) Spectral density at pt. C, Fig. 1. Filtered, $BT_b = 1$, then limited. (d) Spectral density at pt. B, Fig. 1. Filtered, $BT_b = 0.75$. (e) Spectral density at pt. C, Fig. 1. Filtered, $BT_b = 0.75$, then limited.

LPF: 4 pole Chebyshev filter
0.1 dB ripple

Conclusions

- OQPSK is a much better scheme than QPSK
 - Almost same BER performance
 - Significantly better spectrum
 - Same cost!

- These advantages are due to the fact that $\Delta\phi \in \{+90^\circ, -90^\circ, 0^\circ\}$ and not $\pm 180^\circ$!!

- However, it is not a constant envelope (amplitude) scheme, such as:

$$s(t) = \sqrt{\frac{2E}{T}} \cos[2\pi f_c t + \phi(t, \alpha)]$$

A large class of signals with constant envelope are the Continuous Phase Modulation (CPM) signals.

Perhaps the most simple signals in the CPM is the Minimum Shift Keying (MSK)

MINIMUM SHIFT KEYING (MSK)

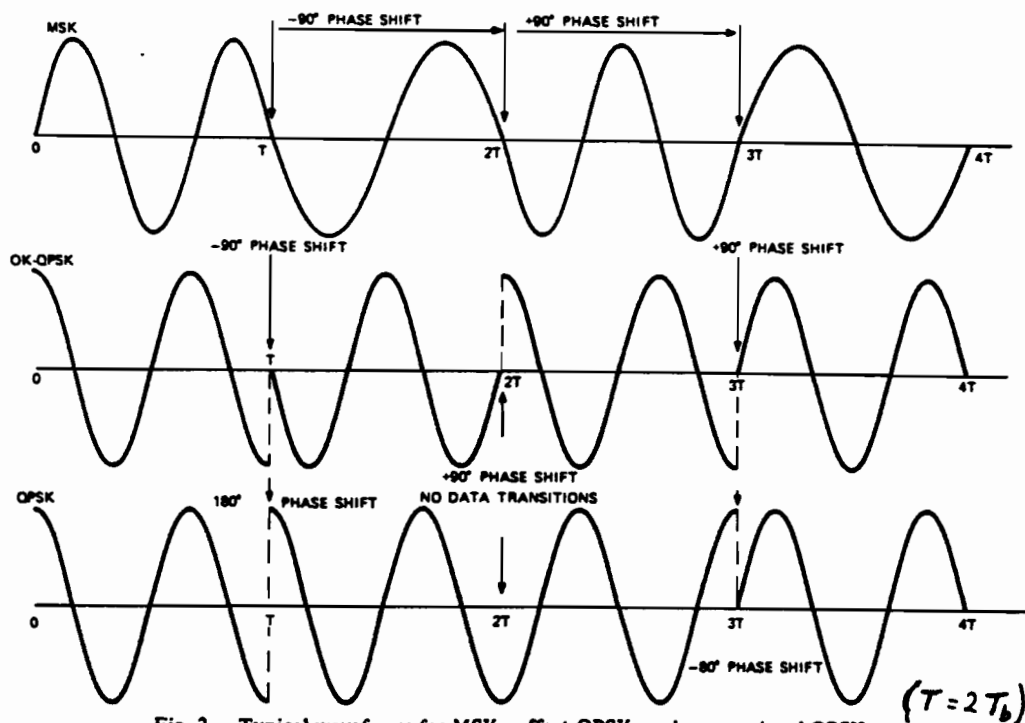


Fig. 3. Typical waveforms for MSK, offset QPSK, and conventional QPSK.

There are two different ways of "seeing" MSK:

- i) As a special case of OQPSK (with sinusoidal pulse shapes)
- ii) As a continuous phase FSK (with a frequency separation of $\frac{1}{2T_b}$, i.e., $\frac{1}{2}$ the bit rate)

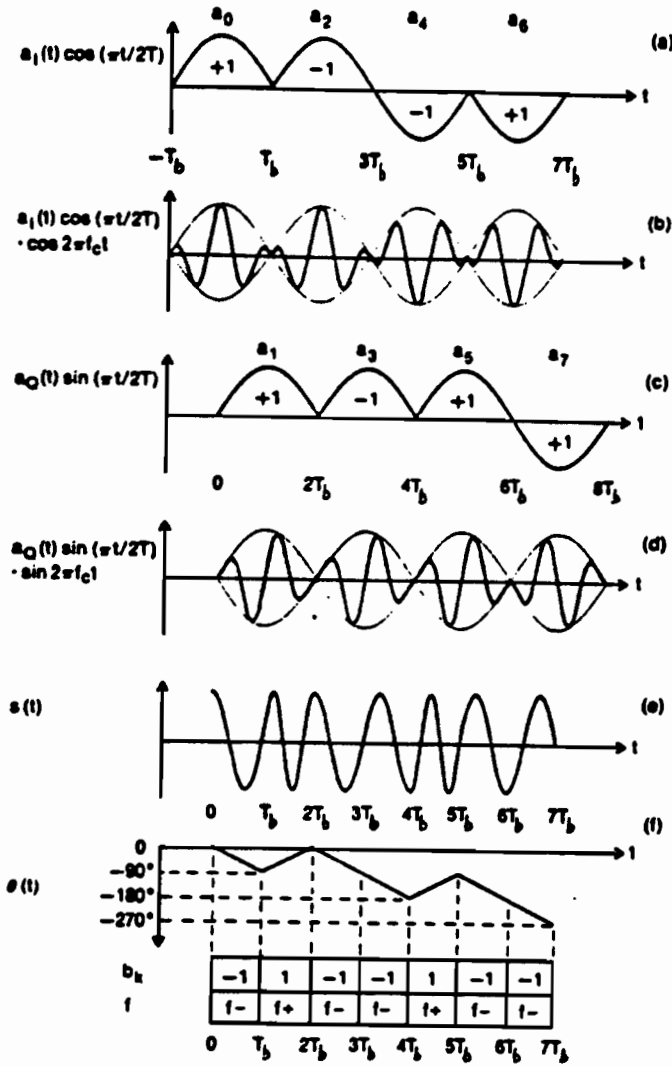


Fig. 5. MSK waveforms.

$$s(t) = a_I(t) \cos\left(\frac{\pi t}{2T_b}\right) \cos(2\pi f_c t) + a_Q(t) \sin\left(\frac{\pi t}{2T_b}\right) \sin(2\pi f_c t)$$

$$= \cos\left[2\pi f_c t + b_k(t) \frac{\pi t}{2T_b} + \phi_k\right]; \quad b_k(t) = -a_I(t)a_Q(t); \quad \phi_k = \begin{cases} 0; & a_I = 1 \\ \pi; & a_Q = -1 \end{cases}$$

$$s(t) = \cos\left[2\pi f_c t + \underbrace{\frac{1}{2} b_k(t) \frac{2\pi t}{2T_b}}_{\theta(t) \triangleq \text{excess phase}} + \phi_k\right]$$

$h \triangleq$ modulation index ($= \frac{1}{2}$ for MSK type of signals)

Comments

- $s(t)$ is a constant envelope signal
- There is a phase continuity in the carrier at the bit transition instants
- $s(t)$ can be also viewed as an FSK signal with signalling frequencies

$$f_+ = f_c + \frac{1}{4T_b} ; f_- = f_c - \frac{1}{4T_b}$$

∴ Frequency deviation $\Delta f = f_+ - f_- = \frac{1}{2T_b}$

Notice that this Δf is the minimum frequency which allows two FSK signals to be orthogonal to each other. In this respect, it is called Fast FSK (FFSK).

- Excess phase $\theta(t) = b_k(t) \frac{\pi t}{2T_b} = \pm \frac{\pi t}{2T_b}$

∴ Linear change (increase or decrease) during each bit period. T_b .

$b_k = +1 \Rightarrow$ increase carrier phase shift by $90^\circ (\frac{\pi}{2}) \leftrightarrow$ FSK; f_+
 $b_k = -1 \Rightarrow$ decrease carrier phase shift by $90^\circ (\frac{\pi}{2}) \leftrightarrow$ FSK; f_-

See also Figure of previous page.

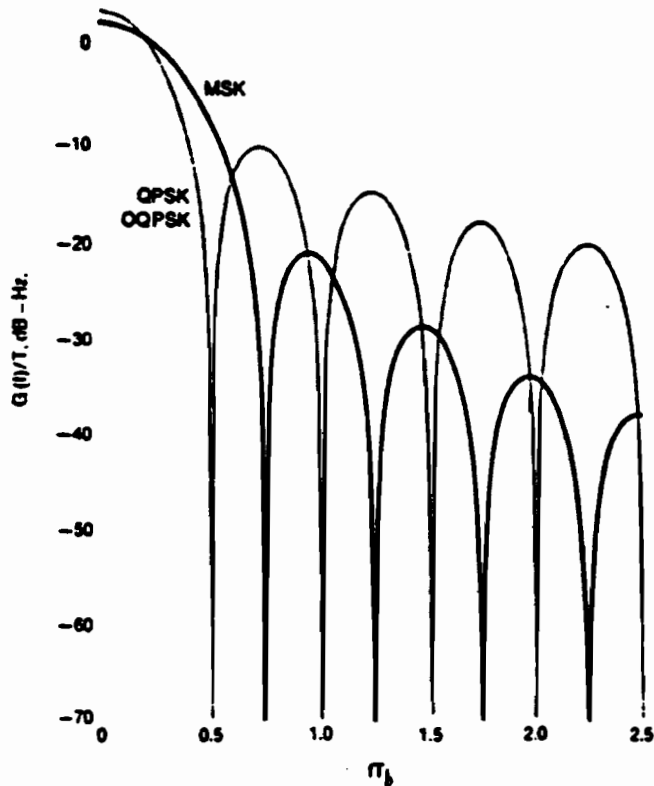


Fig. 6. Spectral density of QPSK, OQPSK, and MSK.

$BW \triangleq$ Bandwidth which contains 99% of the total power

$$MSK \rightarrow BW \approx \frac{1.2}{T_b} ; \quad QPSK \& OQPSK \rightarrow BW \approx \frac{8}{T_b}$$

However MSK has wider mainlobe \Rightarrow in narrowband channels MSK may not be the preferred mod. scheme.

With nonlinearities & ACI things are more complicated.

($\beta \triangleq$ channel spacing \times bit duration)

- MSK is better than QPSK for $\beta > 1.8$
 - - - - - - OQPSK - $\beta > 2.3$
 - OQPSK - - - QPSK - $\beta < 1.4$
- } i) Approximation
ii) Small diffn

Many other schemes which use pulse shaping to improve spectrum

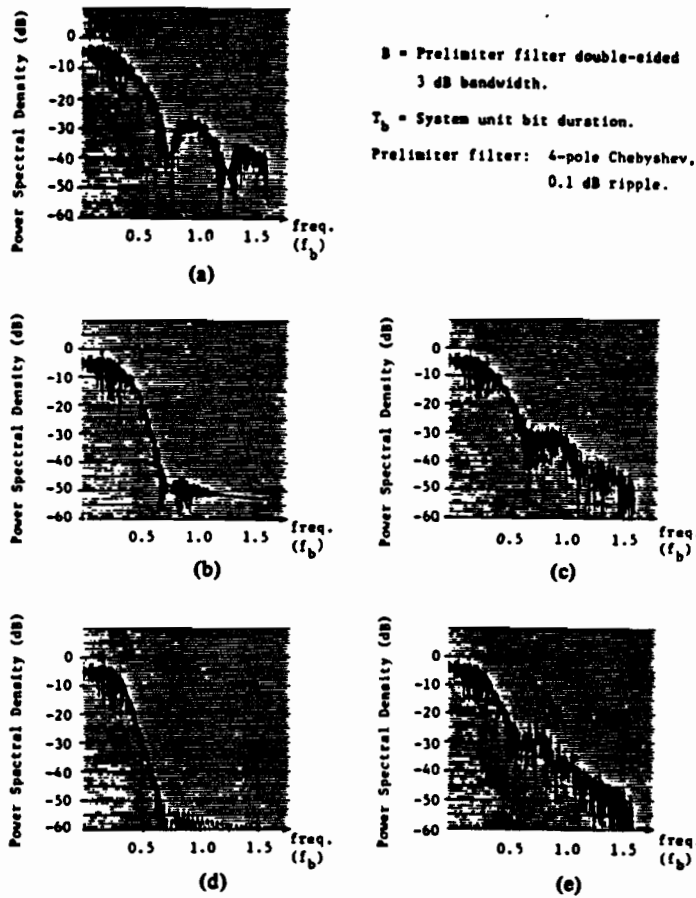


Fig. 4. Computed power spectral densities for MSK. (a) Spectral density at pt. A, Fig. 1. Unfiltered, $BT_b = \infty$. (b) Spectral density at pt. B, Fig. 1. Filtered, $BT_b = 1$. (c) Spectral density at pt. C, Fig. 1. Filtered, $BT_b = 1$, then limited. (d) Spectral density at pt. B, Fig. 1. Filtered, $BT_b = 0.75$. (e) Spectral density at pt. C, Fig. 1. Filtered, $BT_b = 0.75$, then limited.

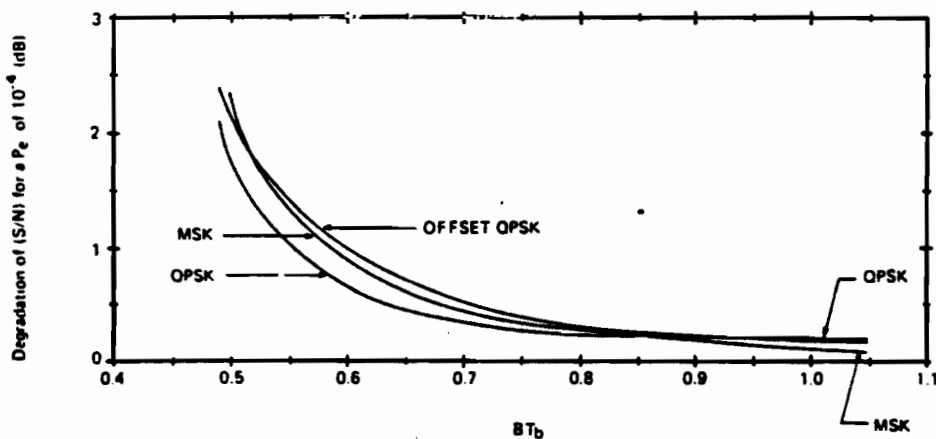
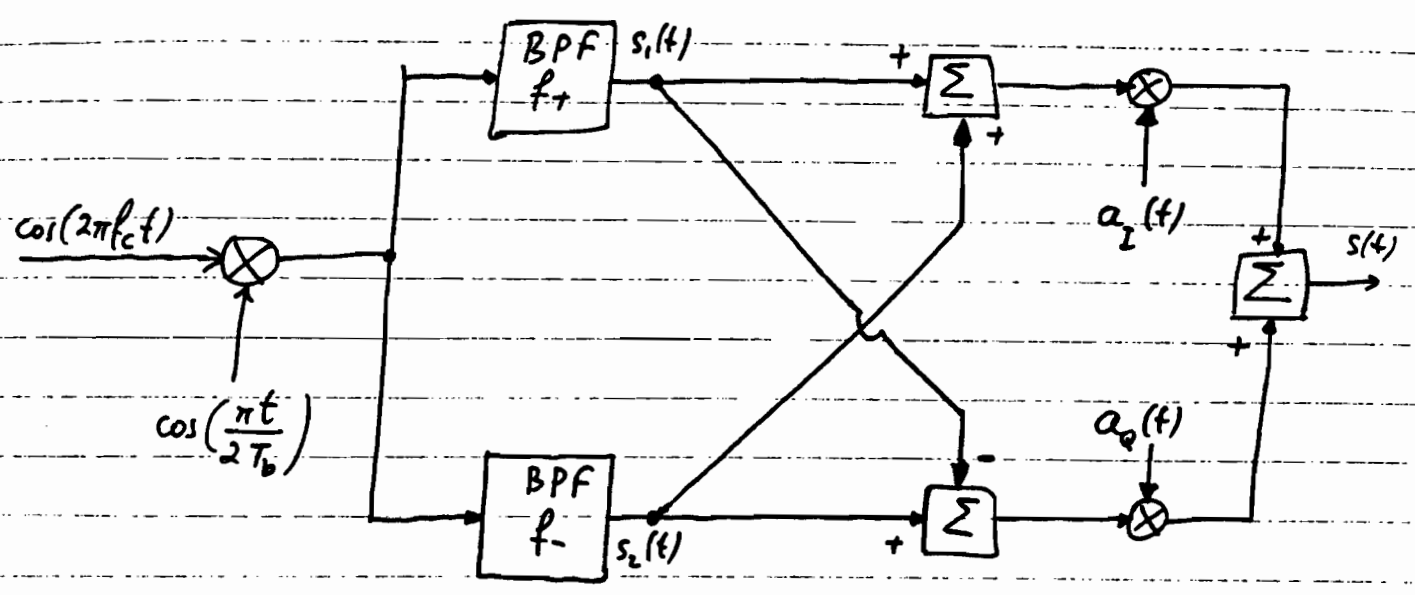


Fig. 9. Degradation, due to filtering, then limiting, of $(S/N)_{rcvr}$ input versus normalized prelimited filter bandwidth.

Ref. Morris & Fcker, "The effects of filtering and limiting on the performance of QPSK, OQPSK and MSK systems," IEEE Trans. Commun., vol. COM-28, pp. 1994-2009, Dec. 1980.

MSK transmitter implementation

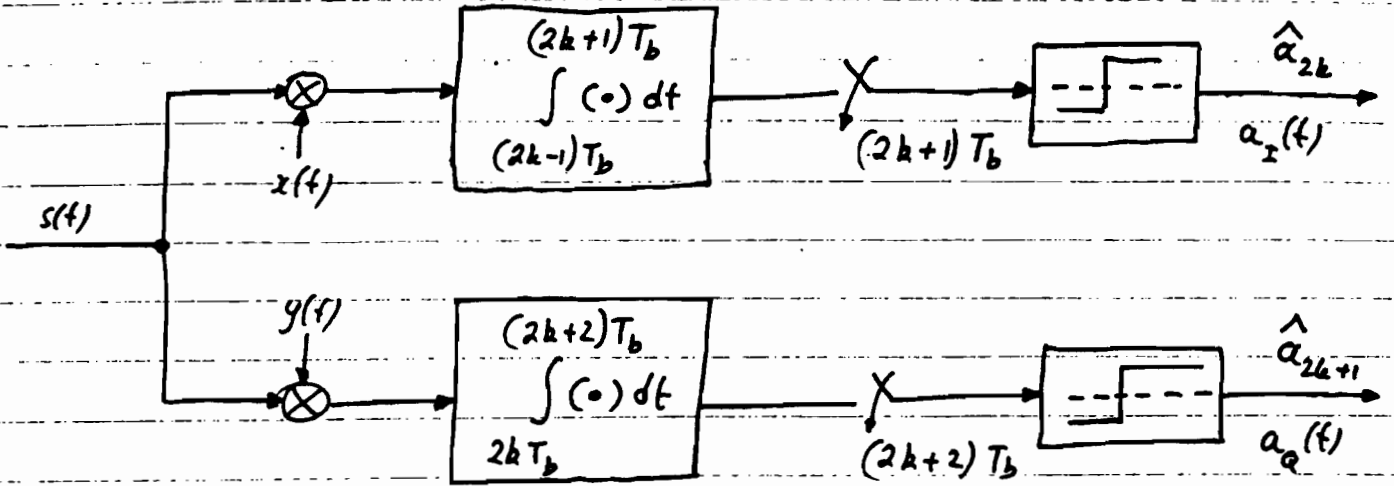


$$f_+ = f_c + \frac{1}{4T_b} \quad ; \quad f_- = f_c - \frac{1}{4T_b}$$

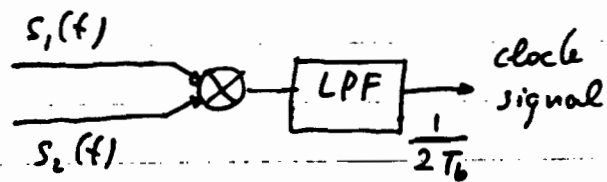
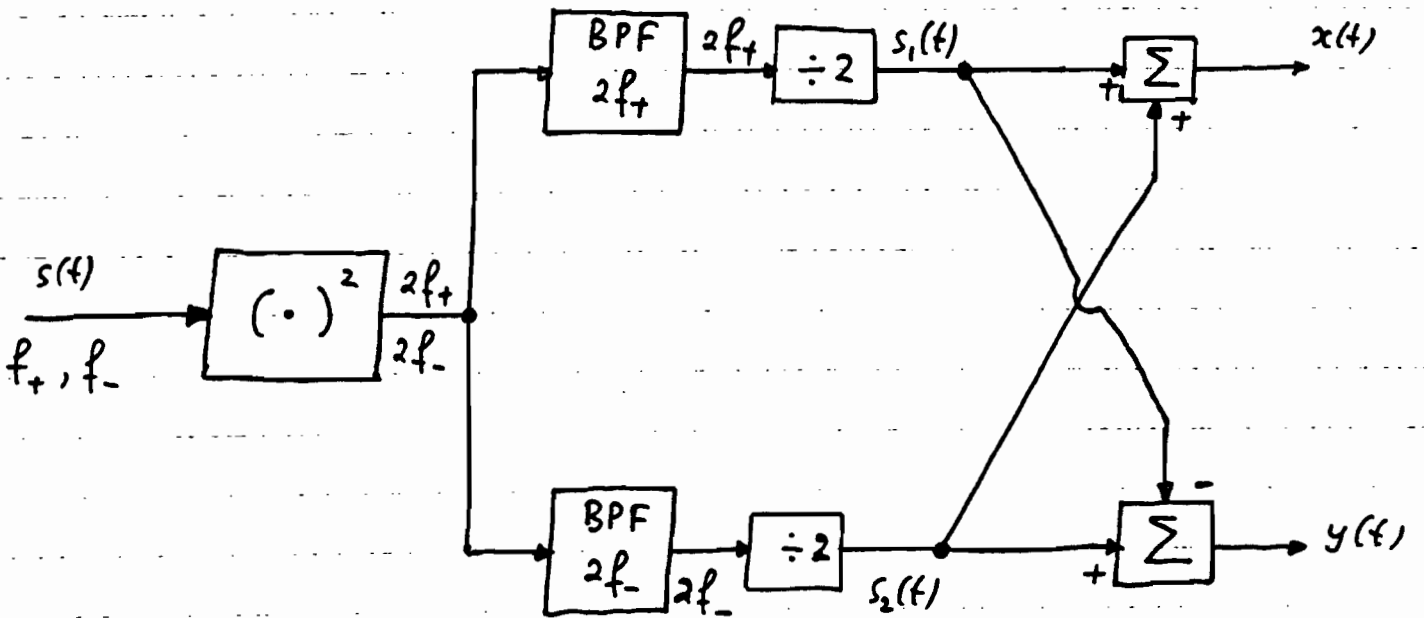
$$s_1(t) = \frac{1}{2} \cos\left(2\pi f_c t + \frac{\pi t}{2T_b}\right)$$

$$s_2(t) = -\frac{1}{2} \cos\left(2\pi f_c t - \frac{\pi t}{2T_b}\right)$$

MSK receiver implementation

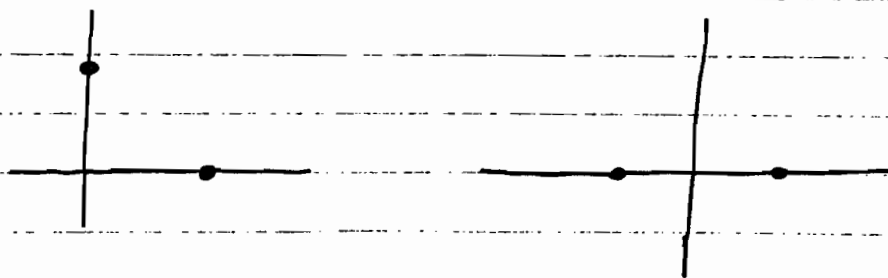


Synchronisation subsystems



Final comments

- For MSK an observation interval of $2T_b$ will result performance equal to QPSK, OQPSK & BPSK.
- If MSK is detected as an FSK signal, i.e., observation over $T_b \rightarrow 3\text{dB}$ degradation



- For other modulation indexes, larger observation intervals are necessary for optimal performance.

e.g. $h = \frac{2}{3} \rightarrow$ gain of 0.8dB for observation over $4T_b$.