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Προηγμένες Μέθοδοι Προγραμματισμού

ΠΜΣ (M135.CS1E, M135.CS23B, M135.IC1E, παλαιό: M117)

Concurrent programming (2)

Δρ. Κώστας Σαΐδης (saiko@di.uoa.gr)

Liveness

- The ability of the multi-threaded, concurrent application to execute in a timely manner.
- Three liveness problems: deadlock, starvation, livelock.

Deadlock

- A situation where two or more threads are waiting for each other, blocking indefinetely.
- There are four conditions that must hold simultaneously for a deadlock to occur:
    - Mutual exclusion: there is at least a non-shareable resource.
    - Hold and wait: a requesting thread holds a resource and waits for some other requested resource.
    - No-preemption allowed: a resource held by the lock cannot be ``taken back''.
    - Circular wait: threads form a chain where each one

Example 08 (BankAccount & BankTransfer)

How to avoid deadlocks

- Provide an ordering to your locks.
- Don't synchronize randomly.

Example 08 (contd.)

Starvation

- A thread is unable to gain access to shared resources and, thus, is unable to make progress.
- This happens when shared resources are made unavailable (locked) for long periods of time.

Livelock

- In cases where a thread (A) act in response to the actions of another thread (B) and the latter also responds to actions of some other thread (C), may lead to livelocked threads.
- As with deadlock, livelocked threads are unable to make progress; the threads are not blocked — they are simply too busy responding to each other to resume any actual work.

Guarded blocks, thread signaling and coordination

- The most common thread coordination idiom is the guarded block; a block begins by polling a condition that must be true before the block can proceed.

Example

```
public synchronized void guardedAction() {
  while(!safe) {
    try {
          wait();
        }
    catch(InterruptedException ie) {}
        ...
        // do the action
  }
}
```

The builtin Object's `wait()` and `notify()` methods

- These methods effectively allow threads to signal each other.
- They are available to all Java objects (they reside in the `java.lang.Object` class).
- A call to `obj.wait()` suspends the current thread. The thread is said to be "waiting on" the `obj`.
- Another thread calls `obj.notify()` or `obj.notifyAll()`. This "wakes up" the threads waiting on the `obj`.
- You need to have the lock/monitor of an object to be able to call wait/notify.

The `wait()` and `wait(long)` methods

- These methods, when called on an object, cause the current thread to wait until either another thread invokes `notify()` / `notifyAll()` on this object, or a specified amount of time has elapsed (timed waiting).
- The current thread releases the object's lock/monitor during a wait.
- When a thread is interrupted during a wait, the thread wakes up and the `wait()` method returns by throwing an `InterruptedException`.

The `notify()` and `notifyAll()` methods

- The first wakes up a single thread that is waiting on this object's monitor. If many threads are waiting on this object, one of them (arbitrary) is chosen to be awakened.
- The second wakes up all threads that are waiting on this object's monitor.

The current thread (the invoker of `notify()` or `notifyAll()`) should release the object's lock in order for the awakened threads to compete again to obtain the object's lock.

When to use guarded blocks with wait/notify signaling?

- In various scenarios that require thread communication
- Producer-Consumer (classic example):
    - A data structure,
    - where one or more threads write to it
    - and one or more thread read from it.

Example 09 (A logging producer/consumer)

- Multiple threads (producers) use a common logger to log application events.
- The logger writes events to standard output (consumer).

Can we avoid synchronization?

- Yes, when we don't share state!
- Making our objects immutable guarantees that no synchronization will be required (the state of our objects does not change).
- Immutability: since the state of immutable objects cannot change, the objects cannot be corrupted by thread interference or observed in an inconsistent state.